

Description

This invention relates to encoding apparatus and methods for error correction which may be applied to record, for example, audio pulse code modulation (PCM) signals onto a magnetic tape using rotary heads and, more particularly but not exclusively, to apparatus and methods for error correction provided in a digital encoder for recording high quality audio PCM data, for example as described in US patent no US-A-4 551 771.

Error correction encoding apparatus and methods have been previously proposed for use with information symbols arranged two-dimensionally in a matrix form in which encoding processes of error detection and error correction codes, for example Reed Solomon codes, are executed in each of the vertical and lateral directions of the information symbols. These codes are transmitted for each column in the vertical direction. On the reception side, the error correction is performed by using a first error detection code and a first error correction code and, at the same time, a pointer indicative of the presence or absence of errors is formed. The errors are then corrected by a second error detection code and a second error correction code with reference to this pointer.

In the case where the foregoing error correction encoded data is transmitted for each column, a synchronisation (sync) signal and sub-data such as a block address and the like are added thereby to form one block of data. For example, in US patent no US-A-4 630 272, there is shown a method whereby a sync signal and an address in which error detection can be independently performed by a cyclic redundancy check (CRC) code are added to each column of data to the the parity data of a first error correction code, thereby forming one block. In the above US patent no US-A-4 630 272, as shown in Figure 1A of the accompanying drawings, for the address, the error detection can be executed by the CRC code and for a data portion (PCM audio signals), encoding processes of a first error correction code (referred to as a C1 code) and a second error correction code (referred to as a C2 code) are performed. In the case of the encoding as shown in Figure 1A, however, since the C1 code is not applied to the address, the protection against errors may be insufficient.

To solve this problem, for example, as disclosed in US patent No. US-A-4 682 332 and as shown in Figure 1B of the accompanying drawings, an error correction encoding is proposed in which an encoding by a C1 code is also executed for the address.

When a header consists of only an address, the error correction encoding shown in Figure 1B is useful. However, if PCM audio signals (main data) are included in the header in addition to the address, the encoding by the C1 code is only executed for the main data and there is a problem in that the protection for errors may be insufficient, as discussed hereinafter. Encoding by the C2 code of the whole header, including the address, to eliminate this drawback causes an inconvenience in that the data area in which the addresses are recorded is lost by virtue of the existence of the C2 parity.

According to one aspect of the invention there is provided data encoding apparatus for error correction of blocks of data, the error correction being of the type in which encoding is performed such that information symbols are respectively included in different first and second error correction code series, and a block of data is formed by a header data portion containing a first number of symbols including at least a block address for discrimination of a data block, and a main data portion containing a second number of symbols of main data which does not include said address, said encoding apparatus comprising: means for forming said first error correction code series including first parity data for a set of symbols in which a plurality of said blocks are arranged with respect to all of the symbols included in said plurality of blocks; and being characterised by:

mixing means for mixing a part of the main data into a part of the header data portion of certain of said plurality of blocks such that the first number of symbols of the header data portion of certain blocks includes main data; and means for forming said second error correction code series including second parity data with regard to said part of the header data portion and all of said main data portion.

According to another aspect of the invention there is provided a data encoding method for error correction of blocks of data, the error correction being of the type in which encoding is performed such that information symbols are respectively included in different first and second error correction code series, and a block of data is formed by a header data portion containing a first number of symbols including at least a block address for discrimination of a data block, and a main data portion containing a second number of symbols of main data which does not include said address, said encoding method comprising the step of: forming said first error correction code series including first parity data for a set of symbols in which a plurality of said blocks are arranged with respect to all of the symbols included in said plurality of blocks; and being characterised by the steps of:

mixing a part of the main data into a part of the header data portion of certain of the plurality of blocks such that the first number of symbols of the header data portion of certain blocks includes main data; and forming said second error correction code series including second parity data with regard to said part of the header data portion and all of said main data portion.

5 An embodiment of the present invention, to be described in greater detail hereinafter, provides an encoding apparatus and method for error correction in which a whole header together with a data portion is C1 encoded and the encoding by the C2 code is performed for the main data included in the header, excluding addresses, thereby enabling error protection for the main data included in the header part to be made secure, so that main data can be recorded into the header part.

10 Preferably, one block is formed by: a header consisting of a block address, an identification (ID) signal, and data; and a data block consisting of only main data. Encoding by the first error correction code (C1 code) is performed for all of the symbols in a two-dimensional array of symbols in which a plurality of blocks are arranged. Therefore, all of the symbols can be error corrected by the C1 code. Encoding by the C2 code is performed for the main data included in the header and the main data in the data part.

15 Accordingly, the main data is securely protected from error, since the main data is encoded by the C1 code and the C2 code. This enables the header to include the main data, thereby increasing the degree of freedom in the frame arrangement of the code.

The invention will now be described by way of example with reference to the accompanying drawings throughout which like parts are referred to by like references, and in which:

20 Figures 1A and 1B are schematic diagrams which are useful for explanation of a previously-proposed error correction code;

Figure 2 is a schematic diagram showing a block constitution of an embodiment of the present invention;

Figure 3 is a schematic diagram showing a frame constitution of the embodiment shown in Figure 2;

Figure 4 is a schematic diagram which is useful for explanation of the formation of a C2 code;

25 Figures 5A, 5B, 5C and 6 are schematic diagrams which show an example of interleaving of data;

Figures 7A, 7B-I, 7B-II, 7B-III, 7B-IV, 7C-I and 7C-II are schematic diagrams which are useful for explanation of the constitution of a header;

Figure 8 is a block diagram of an example of a recording circuit; and

30 Figures 9A and 9B are schematic diagrams of another example of a frame constitution embodying the invention.

An embodiment of the present invention will be described hereinbelow, in accordance with the following order:

(a) a block constitution and a frame constitution;

(b) an example of data interleave;

35 (c) an example of the constitution of a header;

(d) a recording circuit; and

(e) a modification.

40 (a) A block constitution and a frame constitution

Figure 2 shows the constitution of one block of data, for example, audio data which is recorded onto a magnetic tape. This block corresponds to the amount of PCM data recorded in one track of the magnetic tape. One block consists of 49 symbols. A block sync signal of one symbol is located at the head of one block. This is followed by a header of four symbols, and then by a data portion of 44 symbols. As will be explained later, the header may comprise an ID signal, data, or a C2 parity; a block address ADR; and an error detection code EDC such as simple parity, cyclic redundancy check (CRC), or the like. The data portion may comprise data (audio PCM signals) or a C2 parity; or either data or a C2 parity, and a C1 parity.

45 As shown in Figure 3, one frame is formed by arranging the above-mentioned blocks into 100 columns. Block-in-addresses of 0 to 47 are added to 48 symbols in the vertical direction, excluding the block sync signal, in a matrix-shaped frame construction, and block addresses 0 to 99 are added to 100 blocks in the lateral direction.

The audio PCM signals and C1 parity are included in 44 symbols x 80 blocks in the block addresses 20 to 99 and block-in addresses 4 to 47. When the sampling frequency is 48 kHz, the PCM audio signals of one field according to the NTSC system are set to:

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$$48,000 \times \frac{1}{59.94} = 801 \text{ words}$$

In the case of performing a linear digitisation of 16 bits, each word is divided into an upper eight bits and a lower eight bits and one symbol is set to eight bits. On the other hand, when one word consists of 12 bits, one symbol is set to six bits. Instead of 48 kHz, the sampling frequency can be set to 44.1 kHz or 32 kHz.

Referring to Figures 5A, 5B, 5C and 5D, symbols L0u and R0u (Figure 5D) on the upper side of words L0 and R0 in the heads of the audio PCM signals L0 to L800, R0 to R800 of the left and right channels are set to 1 in the block-in address and are arranged in the block addresses 97 and 99. On the other hand, symbols L01 and R01 on the lower side are set to 3 in the block-in address and are arranged in the block addresses 97 and 99. Among the remaining 800 words in each channel, the 800 symbols of the odd-number designated words are arranged in the block addresses 20 to 59 and the 800 symbols of the even-number designated words are arranged in the block addresses 60 to 99.

The parity (C2 parity) of the second error correction code (C2 code) is included in 30 blocks x 44 symbols in the block addresses 0 to 19 and block-in addresses 4 to 47. As shown in Figure 4, the C2 code is a (25, 20) Reed Solomon code which is formed for groups of 20 symbols, every four blocks of the symbols of which are arranged in the lateral direction. Since four series of this C2 code are formed with respect to one row, the C2 parities of $4 \times 5 = 20$ symbols are included in one row. Therefore, the encoding processes of the C1 code and C2 code are executed for all of the symbols of 44 symbols x 80 blocks in the block addresses 20 to 99 and block-in addresses 4 to 47.

The error detection code EDC for the header is included in the block-in address 0 and block addresses 0 to 99. The encoding process of the C2 code is not performed for the error detection code EDC.

An ID signal IDu or data L0u and R0u is included in the block-in address 1 and block addresses 20 to 99. The C2 parities of $5 \times 2 = 10$ symbols formed from this data are included in the block-in address 1 and block addresses 1, 3, 5, 7, ..., 17, 19. Five symbols are needed for each of the data L0u and R0u. An ID signal IDu is included in the block-in address 1 and block addresses 0, 2, 4, 6, ..., 16, 18, as shown in the upper part of Figure 5A.

The block address ADR is included in the block-in address 2 and block addresses 0 to 99. Encoding by the C2 code is not executed for the block address ADR.

The ID signal ID1 or data L01 and R01 is included in the block-in address 3 and block addresses 20 to 99. The C2 parities of $5 \times 2 = 10$ symbols formed from this data are included in the block-in address 3 and block addresses 1, 3, 5, 7, ..., 17, 19. The ID signal ID1 is included in the block-in address 3 and block addresses 0, 2, 4, 6, ..., 16, 18.

The encoding by the C1 code is executed for all one hundred of the blocks in a manner similar to the method described in US patent no US-A-4 630 272. The C1 code is a (48, 44) Reed Solomon code. The series of this C1 code is constructed so as to exist in two adjacent blocks. That is, one C1 series is formed by the even-number designated symbols 0, 2, 4, 6, ... of the block-in address of the series of symbols for two adjacent blocks (for example, blocks 20 and 21 shown in Figure 6). Another C1' series is formed by the odd-number designated symbols 1, 3, 5, 7, ... of the block-in address. The reason why the C1 series are formed so as to exist in two blocks is to prevent the two symbols in one C1 series from including errors due to errors generated at the boundary of two continuous symbols upon recording. The C1 parities (8 symbols) of two adjacent blocks are arranged in the block-in addresses 40 to 47 of the odd-number block addresses.

Upon recording, the C2 parity is formed by the data and ID signal. Next, the C1 parity is formed from these data. Upon reproduction, error detection and/or error correction is executed by means of the C1 code. A pointer is set for the symbols which cannot be error corrected. The error detection and error correction of the C2 code are executed by reference to this pointer. In addition, upon reproduction, error detection regarding the header is performed using the error detection code EDC.

The production of the C2 code will be explained again with reference to Figure 4. As shown in Figure 4, the block-in addresses 0, 1, ..., 47 are expressed as 1 and the block addresses 0, 1, ..., 99 are expressed as k. When $1=0$ and $1=2$, the encoding by the C2 code is not executed, since audio data does not exist. When $1=1$, only the C2 series including the data L0u and R0u (series marked by 0 and X) is formed. When $1=3$, only the C2 series including the data L01 and R01 (series marked by 0 and X) is formed. When $1=4$ to 47, the C2 series (series marked by 0, X, Δ and □) is formed for all of the data.

(b) An example of data interleave

Figures 5 and 6 show the interleave of data of 801 words (=1602 symbols)/one channel of one field according to the NTSC system in detail. Figures 5A and 5B show the data constitution in the block addresses 0 to 59. Figures 5B, 5C and 5D show the data constitution in the block addresses 60 to 99, and Figure 6 shows those in the block addresses 20 and 21 in detail.

As mentioned above, four symbols L0u, L0l, R0u and R0l of two words are positioned in the block-in addresses 1 and 3 in the block addresses 97 and 99. The odd-number designated words L1 to L799, R1 to R799 are arranged in the block addresses 20 to 59. The even-number designated words L2 to L800, R2 to R800 are arranged in the block addresses 60 to 99. By interleaving in the recording position of the odd-number designated words and the even-number designated words, the poor effect caused upon reproduction when consecutive words are error words can be reduced.

An example of the interleave of the symbols of the odd-number designated words will now be described. As shown in Figures 5A, 5B and 6, the data are sequentially arranged from the block-in addresses 4 and 6. In this case, the symbols L1u, R1u, L3u, R3u, ..., R19u on the upper side are successively arranged in the even-number designated block addresses 20, 22, 24, 26, ..., 58 in the block-in address 4. The symbols L1l, R1l, L3l, R3l, ..., R19l on the lower side are sequentially arranged in the even-number designated block addresses 20, 22, 24, 26, ..., 58 in the block-in address 6. The next odd-number designated symbols are arranged in the block-in addresses 5 and 7 in a manner similar to that described above. By repeating the data arrangement in this manner, the symbols R799u and R799l are located in the block-in addresses 37 and 39 in the block address 59.

In Figure 6, P00 to P13 denote C1 parities regarding two blocks in the block addresses, for example, 20 and 21. That is, in the two blocks in the block addresses 20 and 21, and parities P00, P01, P02 and P03 of a (48, 44) Reed Solomon code (C1 code) are formed from the 48 symbols located in the even-number designated block-in address 20. The parities P10, P11, P12 and P13 of the (48, 44) Reed Solomon code (C1' code) are formed from the 48 symbols located in the odd-number designated block-in address 21.

As shown in Figures 5C and 5D, the even-number designated words are arranged in a manner similar to the odd-number designated words. The symbols R800u and R800l of the last word in the R channel are arranged in the block-in addresses 37 and 39 in the block address 99. According to the interleave shown in Figures 5 and 6, the recording positions of adjacent words are interleaved by four blocks in each of the even-number designated word series and the odd-number designated word series in each channel. Successive recording of the upper side symbols and lower side symbols of one word is prevented. The influence of burst errors is reduced. In the actual write operation of the data into memory, odd-numbered words and even-numbered words are written into the tables shown in Figures 5A, 5B, 5C, and 5D, for example L0u, R0u, L1u, R1u, L2u, R2u, L3u, R3u, ...

(c) An example of the constitution of a header

Figure 7A shows a portion of the header in one of the blocks in which C2 parity is not included. The header comprises the ID signals IDu and IDl comprised of two symbols; the block address ADR comprised of one symbol; and a parity of the error detection code (EDC) comprised of three symbols, for instance, of a simple parity. Hereinafter, this parity code is referred to as "third parity data". This error detection code is also used to detect a block sync signal. Figure 7B shows the information of each of the three symbols IDu, IDl and ADR when one symbol consists of six bits, namely one word consists of twelve bits. On the other hand, Figure 7C shows the information of each of the three symbols IDu, IDl, and ADR when one symbol consists of eight bits, namely one word consists of sixteen bits.

As shown in Figure 7B-I, when the least significant bit of the block address ADR is '0', namely, in the even number designated blocks (for example, the area P in Figure 5C), the block address is indicated by a total of eight bits consisting of six bits of the block address ADR and the lower two bits of the ID. The upper four bits of the signal IDl are set to a frame address. The frame address indicates a track number. This frame address is used to discriminate the frame (namely the track) in the high-speed reproducing mode in which the rotary heads scan a plurality of tracks, or in the editing mode. The lower three bits of the signal IDu are set to a track address. The track address, namely the channel address, is used to discriminate a channel when one track is divided into six channels. Six kinds of compressed audio PCM signals are recorded in each of the six channels. The upper three bits of the IDu are used as an ID signal. This ID signal can be used as a head search signal to select a desired track or a time code to indicate the position of a recording signal on the tape. This ID signal can also be used to discriminate the recording mode of the recording signals, for example, stereo/bilingual, prohibition of copy, or the like.

As shown in Figure 7B-II, when the least significant bit of the block address ADR is '1', that is, in the odd-number designated block addresses 97, 99, including main data (for example, the area Q in Figure 5D), the six bits of the ADR are used as a block address, which is an insufficient number of bits to express the full block address. However, the correct block address can be recovered by interpolation using the block addresses of the blocks before and after this block. In this case the area of the IDu is used for the symbols L0u and R0u in this embodiment on the upper side of the data. The area of the IDl is used for the symbols L0l and R0l in the embodiment on the lower side of the data.

Figure 7B-III shows the contents of data in the odd-numbered blocks 21, 23, 25, ..., 95 (for example, the area Y). The ID signal is the so-called sub-code data. This ID signal is used to discriminate the sampling frequency, the number of words, etc., in addition to the above-mentioned use. When three bits are not sufficient to express this control information, six bits included in the odd-numbered blocks are used.

Figure 7B-IV shows the contents of data in the odd-numbered blocks 1, 3, 5, ..., 19 (for example, the area Z) and C2 parity is inserted in the ID area.

Figure 7C shows the constitution of a header when one symbol consists of eight bits. In the block in which the least significant bit of the block address is '0', an ADR or 8 bits is set to a block address, the IDl is set to a track address (3 bits) and a frame address (5 bits), and the IDu is assigned to the ID signal (Figure 7C-I). In the block in which the least significant bit of the block address is '1', the 8 bits of the ADR are used as a block address and the IDu and IDl are assigned to the ID signal, data or C2 parity, respectively (Figure 7C-II).

Even in any of the cases where one symbol consists of 6 bits or where one symbol consists of 8 bits, no data is included in the blocks 20, 22, 24, ..., 112 in which the least significant bit of the block address is '0', so that, as mentioned above, the symbols in the block in which the least significant bit of the block address in the header is '0' are not C2 encoded. C2 encoding need not be performed, since no data is included in the blocks. In this way the block address, frame address, and track address are prevented from being lost by the C2 parity.

(d) Recording circuit

Figure 8 shows a constitution of a recording circuit according to an embodiment of the invention. In this recording circuit, analog audio signals are supplied to an input terminal 1. These analog audio signals are converted into PCM signals by an analog-to-digital (A/D) converter 2. Output data from the A/D converter 2 is supplied to an adder 3. An address signal and an ID signal from an address/ID generating circuit 4 are supplied to the adder 3 where they are added to the PCM audio signals.

The output signal of the adder 3 is used as a data input to random access memories (RAMs) 5 and 6. Each of the RAMs 5 and 6 has a capacity sufficient to store the symbols of one frame. An address generating circuit 7 and a timing generating circuit 8 are provided in conjunction with the RAMs 5 and 6 and are controlled so as to write in and read out data into and from the RAMs 5 and 6 on a symbol unit basis. The reason why the two RAMs 5 and 6 are provided is that, during a period when data is being written into one of the RAMs, data is read out of the other RAM and error correction encoded.

The audio PCM signals read out from one of the RAMs 5 and 6 are supplied to an encoder 10 of the C1 and C2 codes, and the parities of the C1 and C2 codes are formed. These parities are written into either one of the RAMs 5 and 6. A parity generating circuit 9 is provided and the parity of the error detection code for three symbols included in the header is formed. After completion of the error correction encoding process, the digital signals consisting of the parity symbol, block address, ID signal, and data are read out of the RAMs 5 or 6, block by block, and supplied to a parallel/serial converter 11 where they are converted into serial data.

The output data from the parallel/serial converter 11 is supplied to a channel encoder 12 and subjected to a channel encoding process such as 8 to 10 conversion or the like. The output of the channel encoder 12 is supplied to an adding circuit 13 where a block sync signal is added from a sync generator 14. The recording signals output from the adder 13 are supplied to rotary heads 17A and 17B through a recording amplifier 15 and a rotary transformer 16, and recorded onto a magnetic tape 19 by the rotary heads 17A and 17B. The rotary heads 17A and 17B are disposed with an interval of an angle of 180° therebetween. The magnetic tape 19 is run past the heads 17A and 17B at an angle range of 180° or more so as to obtain an overlapping interval of about 30° . The time-base compressed PCM audio signals are recorded in this overlapping interval.

(e) Modification

The invention is not limited to audio signals associated with the video signal of an NTSC system but can also be applied to the case where audio signals associated with the video signal of a CCIR system are recorded. In the CCIR system, since the field frequency is 50 Hz, the data of one field in the case of a sampling frequency of 48 kHz consists of 960 words, L0 to L959, R0 to R959.

Figures 9A and 9B show an example of a frame constitution in the case where the invention is applied to the CCIR system. Forty-eight symbols, excluding the block sync signal, are arranged in the vertical direction and 116 blocks are arranged in the lateral direction. The symbols of the head words L0 and R0 are included in a part of the header. The encoding by the C2 code is performed for this data. The interleave of the even-number designated words and the odd-number designated words, the encoding by the C1 code, and the encoding by the C2 code are similar to those in the case of the NTSC system.

In other embodiments of the invention, error correction codes other than the Reed Solomon code can be used.

Accordingly, the error correction code can be doubly encoded with regard to a part of a header, more specifically, main data included in the header. Therefore, the protection of the data can be made more secure. This means that it is possible to include data in the header. As a result, the degree of freedom in constructing the frame can be increased. On the other hand, address errors, such as in the block address, and the like can be detected, since one error correction code series (C1 code) is generated from symbols of all blocks. Therefore, address read errors can be detected with greater certainty.

Claims

1. Data encoding apparatus for error correction of blocks of data, the error correction being of the type in which encoding is performed such that information symbols are respectively included in different first and second error correction code series (C1, C2), and a block of data is formed by a header data portion containing a first number of symbols including at least a block address for discrimination of a data block, and a main data portion containing a second number of symbols of main data which does not include said address, said encoding apparatus comprising:
means (10) for forming said first error correction code series (C1) including first parity data for a set of symbols in which a plurality of said blocks are arranged with respect to all of the symbols included in said plurality of blocks; and being characterised by:
mixing means (10) for mixing a part of the main data into a part of the header data portion of certain of said plurality of blocks such that the first number of symbols of the header data portion of certain blocks includes main data; and
means (10) for forming said second error correction code series (C2) including second parity data with regard to said part of the header data portion and all of said main data portion.
2. Apparatus according to claim 1, wherein said second parity data is generated with regard to the data series of a block-in address including a part of said main data.
3. Apparatus according to claim 2, wherein sub-code ID data is additionally mixed into said header data portion.
4. Apparatus according to claim 3, wherein the sub-code has a frame address indicative of a track number in which data is recorded.
5. Apparatus according to claim 4, wherein the sub-code has a channel address indicative of the channel number in the track in which data is recorded.
6. Apparatus according to claim 5, wherein said second parity data is gathered in a predetermined area defined by said block address.
7. Apparatus according to claim 6, wherein said sub-code ID data and data of said second parity data are arranged alternately with regard to blocks.
8. Apparatus according to claim 3, wherein third parity data which can be error detected independently of the main data in the blocks is added to said header data portion.

**(e) Modification**

The invention is not limited to audio signals associated with the video signal of an NTSC system but can also be applied to the case where audio signals associated with the video signal of a CCIR system are recorded. In the CCIR system, since the field frequency is 50 Hz, the data of one field in the case of a sampling frequency of 48 kHz consists of 960 words, L0 to L959, R0 to R959.

Figures 9A and 9B show an example of a frame constitution in the case where the invention is applied to the CCIR system. Forty-eight symbols, excluding the block sync signal, are arranged in the vertical direction and 116 blocks are arranged in the lateral direction. The symbols of the head words L0 and R0 are included in a part of the header. The encoding by the C2 code is performed for this data. The interleave of the even-number designated words and the odd-number designated words, the encoding by the C1 code, and the encoding by the C2 code are similar to those in the case of the NTSC system.

In other embodiments of the invention, error correction codes other than the Reed Solomon code can be used.

Accordingly, the error correction code can be doubly encoded with regard to a part of a header, more specifically, main data included in the header. Therefore, the protection of the data can be made more secure. This means that it is possible to include data in the header. As a result, the degree of freedom in constructing the frame can be increased. On the other hand, address errors, such as in the block address, and the like can be detected, since one error correction code series (C1 code) is generated from symbols of all blocks. Therefore, address read errors can be detected with greater certainty.

Claims

1. Data encoding apparatus for error correction of blocks of data, the error correction being of the type in which encoding is performed such that information symbols are respectively included in different first and second error correction code series (C1, C2), and a block of data is formed by a header data portion containing a first number of symbols including at least a block address for discrimination of a data block, and a main data portion containing a second number of symbols of main data which does not include said address, said encoding apparatus comprising:
 means (10) for forming said first error correction code series (C1) including first parity data for a set of symbols in which a plurality of said blocks are arranged with respect to all of the symbols included in said plurality of blocks; and being characterised by:
 mixing means (10) for mixing a part of the main data into a part of the header data portion of certain of said plurality of blocks such that the first number of symbols of the header data portion of certain blocks includes main data; and
 means (10) for forming said second error correction code series (C2) including second parity data with regard to said part of the header data portion and all of said main data portion.
2. Apparatus according to claim 1, wherein said second parity data is generated with regard to the data series of a block-in address including a part of said main data.
3. Apparatus according to claim 2, wherein sub-code ID data is additionally mixed into said header data portion.
4. Apparatus according to claim 3, wherein the sub-code has a frame address indicative of a track number in which data is recorded.
5. Apparatus according to claim 4, wherein the sub-code has a channel address indicative of the channel number in the track in which data is recorded.
6. Apparatus according to claim 5, wherein said second parity data is gathered in a predetermined area defined by said block address.
7. Apparatus according to claim 6, wherein said sub-code ID data and data of said second parity data are arranged alternately with regard to blocks.
8. Apparatus according to claim 3, wherein third parity data which can be error detected independently of the main data in the blocks is added to said header data portion.



9. Apparatus according to claim 8, wherein said third parity data is generated from the address data in said header data portion and said ID data for a sub-code.
10. Apparatus according to claim 8, wherein said third parity data is generated from a part of said main data in said header data portion and said address data.
11. A data encoding method for error correction of blocks of data, the error correction being of the type in which encoding is performed such that information symbols are respectively included in different first and second error correction code series (C1, C2), and a block of data is formed by a header data portion containing a first number of symbols including at least a block address for discrimination of a data block, and a main data portion containing a second number of symbols of main data which does not include said address, said encoding method comprising the step of:
forming said first error correction code series (C1) including first parity data for a set of symbols in which a plurality of said blocks are arranged with respect to all of the symbols included in said plurality of blocks; and being characterised by the steps of:
mixing a part of the main data into a part of the header data portion of certain of said plurality of blocks such that the first number of symbols of the header data portion of certain blocks includes main data; and
forming said second error correction code series (C2) including second parity data with regard to said part of the header data portion and all of said main data portion.
12. A method according to claim 11, wherein said second parity data is generated with regard to the data series of a block-in address including a part of said main data.
13. A method according to claim 12, comprising the step of additionally mixing sub-code ID data into said header data.
14. A method according to claim 13, wherein the sub-code has a frame address indicative of a track number in which data is recorded.
15. A method according to claim 14, wherein the sub-code has a channel address indicative of the channel number in the track in which data is recorded.
16. A method according to claim 13, comprising the step of gathering said second parity data in a predetermined area defined by said block address.
17. A method according to claim 16, comprising the step of arranging said sub-code ID data and data of said second parity data alternately with regard to the blocks.
18. A method according to claim 13, comprising the step of adding to said header data a third parity data which can be error detected independently of the main data in the blocks.
19. A method according to claim 18, comprising the step of generating said third parity data from the address data in said header data and said ID data for the sub-code.
20. A method according to claim 18, comprising the step of generating said third parity data from a part of said main data in said header data and said address data.

Patentansprüche

1. Codierungsgerät zur Fehlerkorrektur von Datenblöcken, mit einer Fehlerkorrektur von dem Typus, in dem die Codierung derart durchgeführt wird, daß Informationssymbole entsprechend in verschiedenen ersten und zweiten Fehlerkorrekturcodierungsserien (C1, C2) enthalten sind und ein Datenblock durch ein eine erste Anzahl von Symbolen, die zumindest eine Blockadresse zur Unterscheidung eines Datenblockes umfassen, enthaltendes Kopfdatenteil gebildet wird und ein eine zweite Anzahl von Symbolen von Hauptdaten enthaltendes Hauptdatenteil, das nicht die besagte Adresse enthält, wobei das Codiergerät folgende Teile enthält:
Mittel (10) zur Ausbildung besagter Fehlerkorrekturcodierungsserien (C1), die erste Paritätsdaten für



einen Satz von Symbolen beinhalten in denen eine Vielzahl von besagten Blöcken mit Bezug zu allen in besagter Vielzahl von Blöcken enthaltenen Symbolen angeordnet ist; und **gekennzeichnet** ist durch:

Mischmittel (10) zum Mischen eines Teiles der Hauptdaten in einen Teil des Kopfdatenteiles von bestimmten Blöcken aus der Vielzahl von Blöcken, so daß die erste Anzahl von Symbolen des Kopfdatenteiles von bestimmten Blöcken Hauptdaten umfaßt; und

Mittel (10) zur Ausbildung besagter zweiter Fehlerkorrekturcodierungsserien (C2), die zweite Paritätsdaten in Bezug auf besagtes Teil des Kopfdatenteiles und alle der besagten Hauptdatenteile beinhaltet.

2. Codierungsgerät nach Anspruch 1, wobei die besagten zweiten Paritätsdaten im Hinblick auf die Datenserien einer in einem Block enthaltenen Adresse, die einen Teil der besagten Hauptdaten umfassen, erzeugt werden.
3. Codierungsgerät nach Anspruch 2, wobei Identifizierungsdaten (ID) in Form eines Sub-Codes zusätzlich in das Kopfdatenteil gemischt werden.
4. Codierungsgerät nach Anspruch 3, wobei der Sub-Code eine Rahmenadresse aufweist, die die Nummer einer Spur bezeichnet, in der Daten aufgezeichnet sind.
5. Codierungsgerät nach Anspruch 4, wobei der Sub-Code eine Kanaladresse aufweist, die die Nummer eines in der Spur in der die Daten aufgezeichnet sind, vorhandenen Kanals bezeichnet.
6. Codierungsgerät nach Anspruch 5, wobei die besagten zweiten Paritätsdaten in einer durch besagte Blockadresse vorbestimmten Zone gesammelt werden.
7. Codierungsgerät nach Anspruch 6, wobei die besagten Identifikationsdaten (ID) in Form eines Sub-Codes und Daten der besagten zweiten Paritätsdaten in Bezug auf die Blöcke abwechselnd angeordnet sind.
8. Codierungsgerät nach Anspruch 3, wobei dritte Paritätsdaten, bei denen die Fehlererfassung unabhängig von Hauptdaten in den Blöcken durchführbar ist, zu dem besagten Kopfdatenteil addiert werden.
9. Codierungsgerät nach Anspruch 8, wobei besagte dritte Paritätsdaten generiert werden von den Adreßdaten im besagten Kopfdatenteil und besagten Identifikationsdaten (ID) für einen Sub-Code.
10. Codierungsgerät nach Anspruch 8, wobei besagte dritte Paritätsdaten generiert werden von einem Teil der besagten Hauptdaten im besagten Kopfdatenteil und besagter Adreßdaten.
11. Verfahren zur Datencodierung zur Fehlerkorrektur an Datenblöcken, mit einer Fehlerkorrektur von dem Typus, bei dem die Codierung derart durchgeführt wird, daß Informationssymbole entsprechend in verschiedenen ersten und zweiten Fehlerkorrekturcodierungsserien (C1, C2) enthalten sind und ein Datenblock durch ein eine erste Anzahl von Symbolen, die zumindest eine Blockadresse zur Unterscheidung eines Datenblockes umfassen, enthaltendes Kopfdatenteil gebildet wird und ein eine zweite Anzahl von Symbolen von Hauptdaten enthaltendes Hauptdatenteil, das nicht die besagte Adresse enthält, wobei das Verfahren zur Codierung folgenden Schritt aufweist:
die Bildung besagter erster Fehlerkorrekturcodierungsserien (C1), die erste Paritätsdaten für einen Satz von Symbolen enthalten, in denen eine Vielzahl besagter Blöcke unter Bezug auf sämtliche in besagter Vielzahl von Blöcken enthaltene Symbole angeordnet ist; und **gekennzeichnet** ist durch folgende Schritte:
Mischen eines Teiles der Hauptdaten in ein Teil des Kopfdatenteiles von bestimmten Blöcken der besagten Vielzahl von Blöcken, so daß die erste Anzahl von Symbolen des Kopfdatenteiles der bestimmten Blöcke Hauptdaten enthält; und
Ausbildung besagter zweiter Fehlerkorrekturcodierungsserien (C2), die zweite Paritätsdaten im Hinblick auf das besagte Teil des Kopfdatenteiles und alle der besagten Hauptdatenteile umfaßt.
12. Verfahren zur Codierung nach Anspruch 11, wobei besagte zweite Paritätsdaten im Hinblick auf die Datenserien einer in einem Block enthaltenen Adresse, die einen Teil der besagten Hauptdaten umfassen, erzeugt werden.



13. Verfahren zur Codierung nach Anspruch 12, wobei der Schritt beinhaltet ist, daß Identifikationsdaten (ID) in Form eines Sub-Codes in die Kopfdaten gemischt werden.

14. Verfahren zur Codierung nach Anspruch 13, wobei der Sub-Code eine Rahmenadresse aufweist, die die Nummer einer Spur bezeichnet, in der Daten aufgezeichnet sind.

15. Verfahren zur Codierung nach Anspruch 14, wobei der Sub-Code eine Kanaladresse aufweist, die die Nummer eines in der Spur, in der die Daten aufgezeichnet sind, vorhandenen Kanales bezeichnet.

16. Verfahren zur Codierung nach Anspruch 13, wobei der Schritt beinhaltet ist, daß besagte zweite Paritätsdaten in einer durch besagte Blockadresse vorbestimmten Zone gesammelt werden.

17. Verfahren zur Codierung nach Anspruch 16, wobei der Schritt enthalten ist, daß die besagten Identifikationsdaten (ID) in Form eines Sub-Codes und Daten der besagten zweiten Paritätsdaten in bezug auf die Blöcke abwechselnd angeordnet sind.

18. Verfahren zur Codierung nach Anspruch 13, wobei der Schritt enthalten ist, daß dritte Paritätsdaten, bei denen die Fehlererfassung unabhängig von Hauptdaten in den Blöcken durchführbar ist, zu den besagten Hauptdaten addiert werden.

19. Verfahren zur Codierung nach Anspruch 18, wobei der Schritt enthalten ist, daß besagte dritte Paritätsdaten generiert werden von den Adreßdaten in besagten Kopfdaten und besagten Identifikationsdaten (ID) für den Sub-Code.

20. Verfahren zur Codierung nach Anspruch 18, wobei der Schritt enthalten ist, daß besagte dritte Paritätsdaten generiert werden von einem Teil der besagten Hauptdaten in besagten Kopfdaten und besagter Adreßdaten.

Revendications

1. Appareil de codage de données pour correction d'erreur dans des blocs de données, la correction d'erreur étant du type dans laquelle le codage est réalisé de telle sorte que des symboles d'information soient respectivement inclus dans des première et seconde séries différentes de codes de correction d'erreur (C1, C2) et qu'un bloc de données soit formé par une section de données d'en-tête contenant un premier nombre de symboles incluant au moins une adresse de bloc pour discriminer un bloc de données et par une partie de données principales contenant un second nombre de symboles de données principales qui n'incluent pas ladite adresse, ledit appareil de codage comprenant :

un moyen (10) pour former ladite première série de codes de correction d'erreur (C1) incluant des premières données de parité pour un jeu de symboles, dans lequel une pluralité desdits blocs sont agencés par rapport à tous les symboles inclus dans ladite pluralité de blocs,

caractérisé par :

un moyen de mélange (10) pour mélanger une partie des données principales dans une partie de la partie de données d'en-tête de certains blocs de ladite pluralité de blocs de telle sorte que le premier nombre de symboles de la partie de données d'en-tête de certains blocs incluent des données principales ; et

un moyen (10) pour former ladite seconde série de codes de correction d'erreur (C2) incluant des secondes données de parité relativement à ladite partie de la partie de données d'en-tête et à la totalité de ladite partie de données principales.

2. Appareil selon la revendication 1, dans lequel lesdites secondes données de parité sont générées par rapport à la série de données d'une adresse d'entrée de bloc incluant une partie desdites données principales.

3. Appareil selon la revendication 2, dans lequel des données d'identification (ID) de sous-code sont en plus mélangées dans ladite partie de données d'en-tête.

4. Appareil selon la revendication 3, dans lequel le sous-code comporte une adresse de trame indicative d'un numéro de piste dans laquelle des données sont enregistrées.



5. Appareil selon la revendication 4, dans lequel le sous-code comporte une adresse de canal indicative du numéro de canal de la piste dans lequel des données sont enregistrées.
6. Appareil selon la revendication 5, dans lequel lesdites secondes données de parité sont groupées dans une zone prédéterminée définie par ladite adresse de bloc.
7. Appareil selon la revendication 6, dans lequel lesdites données ID de sous-code et des données desdites secondes données de parité sont agencées en alternance par rapport à des blocs.
8. Appareil selon la revendication 3, dans lequel des troisièmes données de parité qui peuvent se voir détecter une erreur indépendamment des données principales contenues dans les blocs sont ajoutées à ladite partie de données d'en-tête.
9. Appareil selon la revendication 8, dans lequel lesdites troisièmes données de parité sont générées à partir d'une partie des données d'adresse contenue dans ladite partie de données d'en-tête et à partir desdites données ID pour un sous-code.
10. Appareil selon la revendication 8, dans lequel lesdites troisièmes données de parité sont générées à partir d'une partie desdites données principales contenue dans ladite partie de données d'en-tête et à partir desdites données d'adresse.
11. Procédé de codage de données pour correction d'erreur dans des blocs de données, la correction d'erreur étant du type dans lequel un codage est réalisé de telle sorte que des symboles d'information soient respectivement inclus dans différentes séries de première et seconde séries de codes de correction d'erreur (C1, C2) et un bloc de données étant formé par une partie de données d'en-tête contenant un premier nombre de symboles incluant au moins une adresse de bloc pour discriminer un bloc de données et par une partie de données principales contenant un second nombre de symboles de données principales qui n'incluent pas ladite adresse, ledit procédé de codage comprenant l'étape de :
 - formation de ladite première série de codes de correction d'erreur (C1) incluant des premières données de parité pour un jeu de symboles dans lequel une pluralité desdits blocs sont agencés par rapport à tous les symboles inclus dans ladite pluralité de blocs, caractérisé par les étapes de :
 - mélange d'une partie des données principales dans une partie de la partie de données d'en-tête de certains blocs de ladite pluralité de blocs de telle sorte que le premier nombre de symboles de la partie de données d'en-tête de certains blocs incluent des données principales ; et
 - formation de ladite seconde série de codes de correction d'erreur (C2) incluant des secondes données de parité par rapport à ladite partie de la partie de données d'en-tête et à la totalité de ladite partie de données principales.
12. Procédé selon la revendication 11, dans lequel lesdites secondes données de parité sont générées par rapport à la série de données d'une adresse de bloc incluant une partie desdites données principales.
13. Procédé selon la revendication 12, comprenant l'étape de mélange en plus de données d'identification (ID) de sous-code dans lesdites données d'en-tête.
14. Procédé selon la revendication 13, dans lequel le sous-code comporte une adresse de trame indicative d'un numéro de piste dans laquelle des données sont enregistrées.
15. Procédé selon la revendication 14, dans lequel le sous-code comporte une adresse de canal indicative du numéro de canal dans la piste dans lequel des données sont enregistrées.
16. Procédé selon la revendication 13, comprenant l'étape de groupement desdites secondes données de parité dans une zone prédéterminée définie par ladite adresse de bloc.
17. Procédé selon la revendication 16, comprenant l'étape d'agencement desdites données ID de sous-code et de données desdites secondes données de parité en alternance par rapport aux blocs.



18. Procédé selon la revendication 13, comprenant l'étape d'ajout auxdites données d'en-tête de troisièmes données de parité qui peuvent se voir détecter une erreur indépendamment des données principales contenues dans les blocs.

5 19. Procédé selon la revendication 18, comprenant l'étape de génération desdites troisièmes données de parité à partir des données d'adresse contenues dans lesdites données d'en-tête et à partir desdites données ID pour le sous-code.

10 20. Procédé selon la revendication 18, comprenant l'étape de génération desdites troisièmes données de parité à partir d'une partie desdites données principales contenues dans lesdites données d'en-tête et à partir desdites données d'adresse.

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