

Controlling bandwidth in encoders and/or decoders

Description

5 **1. Prior art**

The present examples relate to encoders and decoders and methods for these apparatus, in particular for information signals, such as audio signals.

General audio codecs need to transmit music and speech signals in a very good quality. Such audio codecs are for instance used in Bluetooth where the audio
10 signals are transmitted from the mobile phone to a headset or headphone and vice versa.

Quantizing parts of a spectrum to zeros often leads to a perceptual degradation. Therefore, it is possible to replace zero-quantized spectral lines with noise using a noise filler tool operating in the frequency domain (FD).

15 Temporal noise shaping (TNS) uses open-loop linear prediction in the frequency domain (FD). This predictive encoding/decoding process over frequency effectively adapts the temporal structure of the quantization noise to that of the time signal, thereby efficiently using the signal to mask the effects of noise. In the MPEG2 Advanced Audio Coder (AAC) standard, TNS is currently implemented by
20 defining one filter for a given frequency band, and then switching to another filter for the adjacent frequency band when the signal structure in the adjacent band is different than the one in the previous band.

Especially, for speech signals, the audio content may be bandlimited, meaning the audio bandwidth contains only 4kHz (narrow band, NB), 8kHz (wide band, WB) or
25 16kHz (super wide band, SWB). Audio codecs need to detect the active audio bandwidth and control the coding tools accordingly. As the detection of the bandwidth is not 100% reliable, technical issues may arise.

Some audio coding tools, e.g. Temporal Noise Shaping (TNS) or noise filling (NF), may cause annoying artefacts when operating on bandlimited audio files, e.g., if

the tool is not aware about the active signal part. Assuming that the WB signal is coded at 32kHz, the tools might fill the upper spectrum (8-16kHz) with artificial noise.

Fig. 1 shows artificial noise generated by unguided tools: line 11 is the active
5 signal up to WB while the signal 12 is artificially generated by a parametric tool, e.g. by noise filling, which is not aware of the active audio bandwidth.

Therefore, the tools need to be restricted to operate only on the active frequency regions.

Some codecs like AAC are configured so as to send the information on active
10 spectrum per scale factor band. This information is also used to control the coding tools. This provides precise results but requires a significant amount of side information to be transmitted. As speech is usually just transmitted in NB, WB, SWB and FB, this limited set of possible active bandwidths is preferably used to limit the side information.

15 It is unavoidable that a bandwidth detector returns wrong results from time to time. For instance, a detector may see the fade out of a music signal and interprets this as a low bandwidth case. For codecs, which switch between the different bandwidth modes (NB, WB, SWB, FB) in a hard manner, e.g. 3GPP EVS codec [1], this results in a rectangular spectral hole. Hard manner means that the
20 complete coding operation is limited to the detected bandwidth. Such hard switch can result in audible artefacts. Figure 2 outlines the spectral hole 22 resulting from a wrong detection.

Fig. 2 shows a schematic outline of wrong bandwidth detection: all coding tools work on lower audio bandwidth, leading to rectangular spectral hole 22.

25 It is requested to overcome or reduce impairments such as those identified above.

1.1 References

[1] 3GPP EVS Codec,

http://www.3gpp.org/ftp//Specs/archive/26_series/26.445/26445-e10.zip, Section 5.1.6 "Bandwidth detection"

2. Summary of the invention

5 In accordance with examples, there is provided an encoder apparatus comprising:

- a plurality of frequency domain, FD, encoder tools for encoding an information signal, the information signal presenting a plurality of frames; and
- an encoder bandwidth detector and controller configured to select a
10 bandwidth for at least a subgroup of the plurality of FD encoder tools, the subgroup including less FD encoder tools than the plurality of FD encoder tools, on the basis of information signal characteristics so that at least one of the FD encoder tools of the subgroup has a different bandwidth with
15 respect to at least one of the FD encoder tools which are not in the subgroup.

Accordingly, it is possible to avoid spectral holes while maintaining in case of wrong detection of the bandwidth.

In accordance with examples, at least one FD encoder tool of the subgroup may be a temporal noise shaping, TNS, tool and/or a noise level estimator tool.

20 In accordance with examples, at least one FD encoder tool which is not in the subgroup is chosen among at least one of linear predictive coding, LPC, based spectral shaper, a spectral noise shaper, SNS, tool a spectral quantizer, and a residual coder.

25 In accordance with examples, the encoder bandwidth detector and controller is configured to select the bandwidth of the at least one FD encoder tool of the subgroup between at least a first bandwidth common to at least one of the FD encoder tools which are not in the subgroup and a second bandwidth different

from the bandwidth of the at least one of the FD encoder tools which are not in the subgroup.

In accordance with examples, the encoder bandwidth detector and controller is configured to select the bandwidth of the at least one of the plurality of FD encoder
5 tools on the basis of at least one energy estimate on the information signal.

In accordance with examples, the encoder bandwidth detector and controller is configured to compare at least one energy estimation associated to a bandwidth of the information signal to a respective threshold to control the bandwidth for the at least one of the plurality of FD encoder tools.

10 In accordance with examples, the at least one of the plurality of FD encoder tools of the subgroup comprises a TNS configured to autocorrelate a TNS input signal within the bandwidth chosen by the encoder bandwidth detector and controller.

In accordance with examples, the at least one of the FD encoder tools which are not in the subgroup is configured to operate at a full bandwidth.

15 Therefore, the bandwidth selection operates only for the tools of the subgroup (e.g., TNS, noise estimator tool).

In accordance with examples, the encoder bandwidth detector and controller is configured to select at least one bandwidth which is within the full bandwidth at which the at least one of the FD encoder tools which are not in the subgroup is
20 configured to operate.

In accordance with examples, the at least one of the remaining FD encoder tools of the plurality of FD encoder tools is configured to operate in open chain with respect to the bandwidth chosen by the encoder bandwidth detector and controller.

25 In accordance with examples, the encoder bandwidth detector and controller is configured to select a bandwidth among a finite number of bandwidths and/or among a set of pre-defined bandwidths.

Therefore, the choice is limited and there is not the necessity of encoding too complicated and/or long parameters. In examples, only one single parameter (e.g., encoded in 0-3 bits) may be used for the bitstream.

In accordance with examples, the encoder bandwidth detector and controller is
5 configured to perform a selection among at least one or a combination of: a 8 KHz, 16 KHz, 24 KHz, 32 KHz, and 48 KHz, and/or NB, WB, SSWB, SWB, FB, etc.

In accordance with examples, the encoder bandwidth detector and controller is configured to control the signalling of the bandwidth to a decoder.

Therefore, also the bandwidth of signals processed by some tools at the decoder
10 may be controlled (e.g., using the same bandwidth).

In accordance with examples, the encoder apparatus is configured to encode a control data field including an information regarding the chosen bandwidth.

In accordance with examples, the encoder apparatus is configured to define a control data field including:

- 15 - 0 data bits corresponding to NB bandwidth;
- 1 data bit corresponding to NB, WB bandwidth;
- 2 data bits corresponding to NB, WB, SSWB bandwidth;
- 2 data bits corresponding to NB, WB, SSWB, SWB bandwidth;
- 3 data bits corresponding to NB, WB, SSWB, SWB, FB bandwidth.

20 In accordance with examples, the encoder apparatus at least one energy estimation is performed by:

$$E_B(n) = \sum_{k=I_{fs}(n)}^{I_{fs}(n+1)-1} \frac{X(k)^2}{I_{fs}(n+1) - I_{fs}(n)} \quad \text{for } n = 0 \dots N_B - 1$$

where $X(k)$ are MDCT (or MDST...) coefficients, N_B is the number of bands and $I_{fs}(n)$ are the indices associated to the band.

In accordance with examples, the encoder apparatus comprises a TNS tool which may be configured to perform a filtering operation including the calculation of an autocorrelation function. One of the possible autocorrelation functions may be in the following form:

5 for each $k = 0..8$

$$r(k) = \begin{cases} r_0(k) & , \text{ if } \prod_{s=0}^2 e(s) = 0 \\ \sum_{s=0}^2 \frac{\sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1-k} X_s(n)X_s(n+k)}{e(s)} & , \text{ otherwise} \end{cases}$$

with

$$r_0(k) = \begin{cases} 1 & , \text{ if } k = 0 \\ 0 & , \text{ otherwise} \end{cases}$$

and

$$e(s) = \sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1} X_s(n)^2 \quad \text{for } s = 0..2$$

10 where $X(k)$ are MDCT coefficients, $\text{sub_start}(f, s)$ and $\text{sub_stop}(f, s)$ are associated to the particular bandwidth as detected by the encoder bandwidth detector and controller.

In accordance with examples, the encoder apparatus may comprise a noise estimator tool which may be configured to estimate a noise level. One of the procedures used for such an estimation may be in the form of

$$L_{NF} = \frac{\sum_{k=0}^{N_E-1} I_{NF}(k) \cdot \frac{|X_f(k)|}{gg}}{\sum_{k=0}^{N_E-1} I_{NF}(k)}$$

15 where gg refers to the global gain, $I_{NF}(k)$ to the identification of the spectral lines on which the noise level is to be estimated, and $X_f(k)$ is the signal (e.g., the MDCT or MDST or another FD spectrum after TNS).

In examples, $I_{NF}(k)$ may be obtained with:

$$I_{NF}(k) = \begin{cases} 1 & \text{if } 24 \leq k < bw_{stop} \text{ and } X_q(i) == 0 \text{ for all } i = k - 3.. \min(bw_{stop}, k + 3) \\ 0 & \text{otherwise} \end{cases}$$

where bw_{stop} depends on the bandwidth detected by the encoder bandwidth detector and controller.

In accordance with examples, there may be provided a decoder apparatus comprising a plurality of FD decoder tools for decoding an information signal
5 encoded in a bitstream, wherein:

the FD decoder tools are subdivided:

- in a subgroup comprising at least one FD decoder tool;
- in remaining FD decoder tools comprising at least one FD decoder tool;

wherein the decoder apparatus is configured so that the at least one of the
10 plurality of decoder tools of the subgroup performs signal processing a different bandwidth with respect to at least one of the remaining FD decoder tools of the plurality of decoder tools.

In accordance with examples, the decoder apparatus may comprise a bandwidth controller configured to choose the bandwidth on the basis of the bandwidth
15 information.

In accordance with examples, the decoder apparatus may be such that the subgroup comprises at least one of a decoder noise estimator tool and/or a temporal noise shape, TNS, decoder.

In accordance with examples, the at least one of the remaining FD decoder tools
20 is at least one of a linear predictive coding, LPC, decoder tool, spectral noise shaper decoder, SNS, tool, a decoder global gain tool, an MDCT or MDST shaping tool.

In accordance with examples, the decoder apparatus may be configured to control the bandwidth of the at least one of the plurality of decoder tools in the subgroup
25 between:

at least a first bandwidth common to at least one of the remaining FD decoder tools; and

at least a second bandwidth different from the first bandwidth.

In accordance with examples, the at least one of the FD remaining decoder tools
5 is configured to operate at a full bandwidth.

In accordance with examples, the at least one of the remaining FD decoder tools is configured to operate in open chain with respect to the bandwidth (e.g., chosen by the bandwidth controller).

In accordance with examples, the bandwidth controller is configured to choose a
10 bandwidth among a finite number of bandwidths and/or among a set of pre-defined bandwidths.

In accordance with examples, the bandwidth controller is configured to perform a choice among at least one or a combination of: a 8 KHz, 16 KHz, 24 KHz, 32 KHz, and 48 KHz and/or NB, WB, SSWB, SWB, FB.

15 In accordance with examples, the decoder may be further comprising a noise filling tool (46) configured to apply a noise level using indices. A technique for obtaining the indices may provide, for example:

$$I_{NF}(k) = \begin{cases} 1 & \text{if } 24 \leq k < bw_{stop} \text{ and } \widehat{X}_q(i) == 0 \text{ for all } i = k - 3.. \min(bw_{stop}, k + 3) \\ 0 & \text{otherwise} \end{cases}$$

where bw_{stop} is obtained on the basis of bandwidth information in the bitstream.

In accordance with examples, the decoder apparatus may comprise a TNS
20 decoder tool configured to perform at least some of the following operations:

$$s^0(\text{start_freq}(0) - 1) = s^1(\text{start_freq}(0) - 1) = \dots = s^7(\text{start_freq}(0) - 1) = 0$$

for $f = 0$ to $\text{num_tns_filters}-1$ do

for $n = \text{start_freq}(f)$ to $\text{stop_freq}(f) - 1$ do

$$t^K(n) = \widehat{X}_f(n)$$

for $k = 7$ to 0 do

$$\begin{aligned}
 t^k(n) &= t^{k+1}(n) - rc_q(k)s^k(n-1) \\
 s^{k+1}(n) &= rc_q(k)t^k(n) + s^k(n-1) \\
 \widehat{X}_s(n) &= s^0(n) = t^0(n)
 \end{aligned}$$

where $\widehat{X}_s(n)$ is the output of the TNS decoder and $\widehat{X}_f(n)$ is the input of the TNS
 5 decoder, num_tns_filters, start_freq, stop_freq are obtained on the basis of
 bandwidth information in the bitstream.

Coding tools like TNS or noise filling can create unwanted artificial noise in the
 silent sections of band limited signals. Therefore, bandwidth detectors are usually
 in-cooperated to control the bandwidth all coding tools should work on. As
 10 bandwidth detection might lead to uncertain results, such wrong detection might
 lead to audible artefacts such as sudden limitation of audio bandwidth.

To overcome the problem, in some examples some tools, e.g., the quantizer, are
 not controlled by the bandwidth detector. In case of miss-detection, the quantizer
 can code the upper spectrum – even tough in low quality – to compensate the
 15 problem.

3. Description of the drawings

Fig. 1 shows artificial noise generated by unguided tools;

Fig. 2 shows a schematic outline of wrong bandwidth detection;

Figs. 3 and 3a show encoder apparatus according to examples;

20 Figs. 4 and 4a show decoder apparatus according to examples;

Fig. 5 shows a scheme in case of wrong bandwidth detection;

Fig. 6a-6c show methods according to examples;

Fig. 7 shows a method for TNS at the encoder according to an example;

Figs. 8a and 8b show apparatus according to examples.

25 4. Discussion on the invention

The invention described in this document permits to avoid the occurrence of spectral holes even when the bandwidth detector returns a wrong result. In particular, soft band switching for audiocoding applications may be obtained.

5 A key aspect is that parametric coding tools, e.g. TNS and NF, may be strictly controlled by the bandwidth detector and controller 39 while the remaining coding, i.e. LPC based spectral shaper or spectral noise shaper, SNS, spectral quantizer and residual coder, still work on the full audio bandwidth up to the Nyquist frequency.

10 Figs. 3 and 3a outline examples of an encoder apparatus 30 and 30a where a bandwidth (BW) detector and controller 39 estimates the current audio bandwidth in the frame based on energies derived from an MDCT or MDST spectrum (or other FD spectrum).

On the decoder side (Figs. 4 and 4a), the guiding bandwidth information for TNS and NF is extracted from the bitstream and the tools are controlled accordingly.

15 As a result, artificially generated noise in non-active spectral regions is avoided due to the bandwidth parameter used to control the TNS and NF coding tools (unguided tools). The tool just work on the active audio part and therefore do not generate any artificial noise.

20 On the other side, the audible effect of wrong detections (false bandwidth detection) can be reduced significantly as the remaining coding tools, e.g. spectral quantizer, LPC shaper or SNS (spectral noise shaper) and residual coder, still work up to the Nyquist frequency. In case of wrong detections, these tools can code the upper frequency – at least with some more distortions compared to a regular coding – and therefore avoid the more severe impression that the audio
25 bandwidth suddenly drops.

Fig. 5 shows a new scheme in case of wrong bandwidth detection: spectral hole is quantized sparsely but avoids an audible bandwidth drop.

In case the region outlined in the figure above contains mostly zero values, the arithmetic coder does not need to code those as the information on the last non-

zero spectral tuple is transmitted as side information for the arithmetic coder. This means there is no overhead required for the arithmetic coder.

The required side information for the transmitted bandwidth is also minimized. Due to the robust switching behavior, a signaling of the typically used communication
5 audio bandwidths, i.e. NB, WB, SSWB and SWB, is appropriate.

This technique also allows to build less complex bandwidth detectors which do not use frame dependencies and long history memories to get stable decisions, see the EVS codec [1] Section 5.1.6. This means, the new technique allows the bandwidth detector and controller 39 to react very fast on any audio bandwidth
10 change.

Accordingly, a bandwidth information is used to only control specific tools of a codec (e.g., audio codec) while keeping the remaining tools in another operation mode (e.g., full bandwidth).

5. Examples

15 5.1 The bandwidth detection and control of the tools

An information signal (e.g., an audio signal) may be described in the time domain, TD, as a succession of samples (e.g., $x(n)$) acquired at different discrete time instants (n). The TD representation may be made of a plurality of frames, each associated to a plurality of samples (e.g., 2048 samples per frame). In the
20 frequency domain, FD, a frame may be represented as a succession of bins (e.g., $X(k)$), each associated to a particular frequency (each frequency being associated to an index k).

Figs. 3 and 3a show encoder apparatus 30 and 30a, respectively, each comprising an encoder bandwidth detector and controller 39 which is capable of selecting a
25 bandwidth for some tools (a subgroup at the encoder) of the encoder apparatus 30 or 30a, so that other tools operate at different bandwidth. The encoder bandwidth detector and controller 39 is also capable of selecting the bandwidth for at least some of the tools of a decoder (a subgroup at the decoder). 39a refers to the

bandwidth selection information provided by the encoder bandwidth detector and controller 39 to the tools of the subgroup (e.g., 33, 36) and/or to a decoder.

Each of the encoder apparatus 30 and 30a may comprise a low delay modified discrete cosine transform, MDCT, tool 31 or low delay modified discrete sine
5 transform, MDST, tool 31 (or a tool based on another transformation, such as a lapped transformation) which may convert an information signal (e.g., an audio signal) from a time domain, TD, representation to a frequency domain, FD, representation (e.g., to obtain MDCT, MDST, or, more in general, FD coefficients).

The encoder apparatus 30 may comprise a linear predictive coding, LPC, tool 32
10 for performing an LPC analysis in the FD.

The encoder apparatus 30a may comprise an SNS tool 32a for performing an SNS analysis in the FD.

Each of the encoder apparatus 30 and 30a may comprise a temporal noise shaping, TNS, tool 33, to control the temporal shape of noise within each window
15 of the information signal (e.g., as output by the MDCT or MDST tool) in the FD.

Each of the encoder apparatus 30 and 30a may comprise a spectral quantizer 34 processing signals in the in the FD. The signal as output by the TNS tool 33 may be quantized, e.g., using dead-zone plus uniform thresholds scalar quantization. A gain index may be chosen so that the number of bits needed to encode the
20 quantized FD signal is as close as possible to an available bit budget.

Each of the encoder apparatus 30 and 30a may comprise a coder 35 processing signals in the FD, for example, to perform entropy coding, e.g., to compress a bitstream. The coder 35 may, for example, perform residual coding and/or arithmetic coding.

25 Each of the encoder apparatus 30 and 30a may comprise, for example, a noise level estimator tool 36, processing signals in the FD, to estimate the noise, quantize it, and/or transmit it in a bitstream.

In examples, the level estimator tool 36 may be placed upstream or downstream to the coder 35.

Each of the encoder apparatus 30 and 30a may comprise tools which process signals in the time domain, TD. For example, the encoder apparatus 30 or 30a
5 may comprise a re-sampling tool 38a (e.g., a downsampler) and/or a long term postfiltering, LTPF, tool 38b, for controlling an LTPF active in TD at the decoder.

Each of the encoder apparatus 30 and 30a may comprise a bitstream multiplexer tool 37 to prepare a bitstream with data obtained from TD and/or FD tools placed upstream. The bitstream may comprise a digital representation of an information
10 signal together with control data (including, for example, a bandwidth information for selecting the bandwidth at some tools of the decoder) to be used at the decoder. The bitstream may be compressed or include portions which are compressed.

Therefore, each of the encoder apparatus 30 and 30a may comprise FD tools
15 (e.g., 31-36) and, in case, TD tools (e.g., 38a, 38b).

The encoder bandwidth detector and controller 39 may control the bandwidth of FD tools forming a first group (subgroup), such as the temporal noise shaping, TNS, tool 33, and/or the noise estimator tool 36. The TNS tool 33 may be used to control the quantization noise. The bandwidth at which FD tools which are not in
20 the subgroup (such as at least one of the LPC tool 32 and/or the SNS tool 32a, the spectrum quantizer 34, and the coder 35) perform signal processing may therefore be different from the bandwidth at which the tools of the subgroup (e.g., 33, 36) perform signal processing. For example, the bandwidth for the FD tools which are not in the subgroup may be greater, e.g., may be a full bandwidth.

25 In examples, the encoder bandwidth detector and controller 39 may be a part of a digital signal processor which, for example, implements also other tools of the encoder apparatus.

Figs. 4 and 4a show decoder apparatus 40 and 40a, respectively, each of which may decode a digital representation of an information signal as encoded by the

encoder 30 or 30a, for example. Each of the decoder apparatus 40 and 40a may comprise FD tools and, in case, TD tools.

Each of the decoder apparatus 40 and 40a may comprise a bitstream multiplex tool 41 to obtain a bitstream (e.g., by transmission) from an encoder apparatus
5 (e.g., the apparatus 30 or 30a). For example, an output from the encoder apparatus 30 or 30a may be provided as an input signal to the decoder apparatus 40 or 40a.

Each of the decoder apparatus 40 and 40a may comprise a decoder 42 which may, for example, decompress data in the bitstream. Arithmetic decoding may be
10 performed. Residual decoding may be performed.

Each of the decoder apparatus 40 and 40a may comprise a noise filling tool 43 processing signals in the FD.

Each of the decoder apparatus 40 and 40a may comprise a global gain tool 44 processing signals in the FD.

15 Each of the decoder apparatus 40 and 40a may comprise a TNS decoder tool 45 processing signals in the FD. TNS can be briefly described as follows. At the encoder-side and before quantization, a signal is filtered in the frequency domain (FD) using linear prediction, LP, in order to flatten the signal in the time-domain. At the decoder-side and after inverse quantization, the signal is filtered back in the
20 frequency-domain using the inverse prediction filter, in order to shape the quantization noise in the time-domain such that it is masked by the signal.

Each of the decoder apparatus 40 and 40a may comprise an MDCT or MDST shaping tool 46 (other kinds of shaping tools may be used). Notably, the MDCT or MDST shaping tool 46 may process signals by applying scale factors (or quantized
25 scale factors) obtained from the encoder SNS tool 32a or gain factors computed from decoded LP filter coefficients (obtained from an LPC decoding tool 47) transformed to the MDCT or MDST spectrum.

Each of the decoder apparatus 40 and 40a may comprise an inverse low delay inverse MDCT or MDST tool 48a to transform signal representations from FD to TD (tools based on other kinds of inverse transform may be used).

Each of the decoder apparatus 40 and 40a may comprise an LTPF tool 48b for
5 performing a postfilter in the TD, e.g., on the basis of the parameters provided by the component 38b at the decoder.

Each of the decoder apparatus 40 and 40a may comprise a decoder bandwidth controller 49 configured to select the bandwidth of at least one of the FD tools. In particular, the bandwidth of a subgroup (e.g., formed by the tools 43 and 45) may
10 be controlled so as to be different from the bandwidth at which other FD tools (42, 44, 46, 47) process signals. The bandwidth controller 49 may be input with a signal 39a which has been prepared at the encoder side (e.g., by the bandwidth detector and controller 39) to indicate the selected bandwidth for at least one of the subgroups (33, 36, 43, 45).

15 In examples, the decoder bandwidth controller 49 may perform operations similar to those processed by the encoder bandwidth detector and controller 39. However, in some examples, the decoder bandwidth controller 49 may be intended as a component which obtains control data (e.g., encoded in a bitstream) from the encoder bandwidth detector and controller 39 and provides the control data (e.g.,
20 bandwidth information) to the tools of the subgroup (e.g., decoder noise filling tool 43 and/or TNS decoder tool 45). In examples, the controller 39 is a master and the controller 49 is a slave. In examples, the decoder bandwidth controller 49 may be a part or a section of a digital signal processor which, for example, implements also other tools of the decoder.

25 In general, the bandwidth controllers 39 and 49 may operate so that the FD tools of the subgroups (e.g., 33 and 36 for the encoder apparatus and/or 43 and 45 for the decoder apparatus) have a same frequency band, while the other FD tools of the decoder and/or encoder have another frequency band (e.g., a broader band).

It has been noted, in fact, that accordingly it is possible to reduce impairments of
30 the prior art. While for some FD tools (e.g., TNS tools, noise filling tools) it may be

preferable to actually perform a band selection, for other FD tools (e.g., 32, 34, 35, 42, 44, 46, 47) it may be preferable to process signals at a broader band (e.g. full band), Accordingly, it is possible to avoid spectral holes that would be present in case of hard selection of the bandwidth for all the tools (in particular when a wrong
5 band is selected).

In examples, the bandwidth that is selected by the decoder bandwidth controller 49 may be one of a finite number of choices (e.g., a finite number of bandwidths). In examples, it is possible to choose among narrow band NB (e.g., 4KHz), wide band WB (e.g., 8KHz), semi-super wide band SSWB (e.g., 12KHz), super wide
10 band SWB (e.g., 16KHz) or full band FB (e.g., 20KHz).

The selection may be encoded in a data field by the encoder apparatus, so that the decoder apparatus knows which bandwidths have been selected (e.g., according to a selection performed by the encoder bandwidth detector and controller 39).

15 Fig. 6a shows a method 60. The method 60 may comprise steps which may be performed, at least in part, by at least one of the controllers 39 and 49. The method 60 may be looped so as to perform operations in association to each frame of the information signal.

At step S61, an energy per band may be estimated (e.g., by the bandwidth
20 detector and controller 39).

At step S62, the bandwidth may be detected (e.g., by the bandwidth detector and controller 39).

At step S63, the detected bandwidth may be selected for at least one of the TNS tool 33 and noise estimation tool 36: these tools will perform their processes at the
25 bandwidth detected at S62.

In addition or in alternative, at step S64 parameters may be defined (and/or encoded) in the bitstream to be stored and/or transmitted and to be used by a decoder. Among the parameters, a bandwidth selection information (e.g., 39a)

may be encoded, so that the decoder will know the detected and selected bandwidth for the subgroup (e.g., TNS and noise filling/estimation).

Then, a new frame of the information signal may be examined. Method 60 may therefore cycle by moving to S61. Therefore, a decision may be carried out frame
5 by frame.

Notably, in accordance to the detected bandwidth, a different number of bits may be encoded in the bitstream. In examples, if a bandwidth 8KHz (NB) is detected, no bits will be encoded in the bitstream. However, the decoder will understand that the bandwidth is 8KHz.

- 10 Each of the encoder apparatus 30 and 30a of Fig. 3 and 3a may comprise:
- a plurality of frequency domain, FD, encoder tools (31-36) for encoding an information signal, the information signal presenting a plurality of frames; and
 - an encoder bandwidth detector and controller 39 configured to select a
15 bandwidth (e.g., at S63) for at least a subgroup (e.g., TNS tool 33, and noise level estimator tool 36) of the plurality of FD encoder tools on the basis of information signal characteristics so that at least one (e.g., 33, 36) of the FD encoder tools of the subgroup has a different bandwidth with respect to at least one of the FD encoder tools (e.g., 31, 32, 34, 35) which
20 are not in the subgroup.

In particular, the encoder bandwidth detector and controller 39 may be configured to select the bandwidth of the at least one FD encoder tool of the subgroup (33, 36) between at least a first bandwidth (e.g., Nyquist frequency) common to at least one (or more) of the FD encoder tools which are not in the subgroup and a second
25 bandwidth (e.g., NB, WB, SSWB, SWB) different from the bandwidth of the at least one (or more) of the FD encoder tools which are not in the subgroup.

Therefore, some tools may operate at bandwidths different from each other and/or perform signal processing using bandwidths different from each other.

The tools which are not in the subgroup (e.g., global gain, spectral noise shaping, and so on) may operate in open chain which respect to the bandwidth selection.

In examples, the encoder bandwidth detector and controller 39 is configured to select (e.g., at S62) the bandwidth of the at least one of the plurality of FD encoder tools (31-36) on the basis of at least one energy estimation (e.g., at S61) on the information signal.

The decoder apparatus 40 of Fig. 4 comprises comprising a plurality of FD decoder tools (43-48a) for decoding an information signal encoded in a bitstream, wherein:

10 the FD decoder tools are divided:

- in a subgroup comprising at least one FD decoder tool (e.g., 43, 45);
- in remaining FD decoder tools comprising at least one FD decoder tool (e.g., 44, 46, 48a);

wherein the decoder apparatus 40 or 40a is configured so as to choose a
15 bandwidth for at least one of the plurality of decoder tools of the subgroup (e.g., 43, 45) on the basis of bandwidth information included in the bitstream so that the at least one of the plurality of decoder tools of the subgroup (e.g., 43, 45) performs signal processing a different bandwidth with respect to at least one of the remaining FD decoder tools of the plurality of decoder tools (e.g., 44, 46, 48a).

20 Fig. 6b shows a method 60b. The method 60b may be a method for encoding an information signal according to at least a plurality of operations in the frequency domain, FD, the method comprising:

selecting a bandwidth for a subgroup of FD operations (e.g., S61b);

performing first signal processing operations at the a bandwidth for the
25 subgroup of FD operations (e.g., S62b);

performing second signal processing operations at a different bandwidth for FD operations which are not in the subgroup (e.g., S63b).

It is not necessary, e.g., to perform the steps S61b and S62b in this temporal order. For example, S62b may be performed before S61b. S61b and S62b may also be performed in parallel (e.g., using time-sharing techniques or similar).

Fig. 6c shows a method 60c. The method 60c may be a method for decoding a
 5 bitstream with an information signal and control data (e.g., 39a), the method comprising a plurality of signal processing operations in the frequency domain, FD, the method comprising:

choosing a bandwidth selection for a subgroup of FD operations on the basis of the control data (S61c);

10 performing first signal processing operations at the a bandwidth for the subgroup of FD operations (S62c);

performing second signal processing operations at a different bandwidth for FD operations which are not in the subgroup (S63c).

It is not necessary, e.g., to perform the steps S61c and S62c in this temporal
 15 order. For example, S62c may be performed before S61c. S61c and S62c may also be performed in parallel (e.g., using time-sharing techniques or similar).

According to an example, the encoder bandwidth detector and controller 39 may detect the energy per band, e.g., using an equation such as:

$$E_B(n) = \sum_{k=I_{fs}(n)}^{I_{fs}(n+1)-1} \frac{X(k)^2}{I_{fs}(n+1) - I_{fs}(n)} \text{ for } n = 0 \dots N_B - 1$$

where $X(k)$ are the MDCT or MDST coefficients (or any other representation of the
 20 signal in the FD), N_B (e.g., 64) is the number of bands and $I_{fs}(n)$ are the indices associated to the band (each index being associated to a bin).

It is therefore possible to detect (e.g., at S62) the bandwidth (e.g., among a finite number of bandwidths). The encoder bandwidth detector and controller 39 may be able to detect the commonly used bandwidth in speech communication, i.e. 4kHz,
 25 8kHz, 12kHz and 16kHz. For example, it is possible to detect the quietness of

each bandwidth. In case of a positive detection of quietness for a bandwidth, a dedicated cut-off characteristics on the spectrum is further detected. For example, a flag (or in any case a data) regarding the detection of quietness may be obtained as:

$$F_Q(bw) = \sum_{n=I_{bw\ start}(bw)}^{I_{bw\ stop}(bw)} \frac{E_b(n)}{I_{bw\ stop}(bw) - I_{bw\ start}(bw) + 1} < T_Q(bw)$$

for $bw = N_{bw} - 1 \dots 0$

- 5 The $F_Q(bw)$ is a binary value which is 1 if the summation is less than $T_Q(bw)$, and 0 if the summation is greater than $T_Q(bw)$. $F_Q(bw)$, associated to a particular bandwidth bw , indicates quietness (e.g., with logical value "1") when the summation of the energy values is less than a threshold for the particular bandwidth bw (and "0" otherwise). The summation relates to the sum of energy
- 10 values at different indexes (e.g., energy per bin or band), e.g., for n from a first index of the bandwidth associated to the index $I_{bw\ start}(bw)$ to a last index of the bandwidth associated to the index $I_{bw\ stop}(bw)$. The number of the examined bandwidths is N_{bw} .

- The procedure may stop when $F_Q(bw) == 0$ (energy greater than the threshold for
- 15 the bandwidth bw). In case $F_Q(bw + 1) == 1$, the flags $F_C(b)$ indicating the cut-off characteristic of the spectrum may be detected by

$$F_C(b) = [10\log_{10}(E_b(b - D)) - 10\log_{10}(E_b(b))] < T_C(bw)$$

for $b = I_{bw\ start}(bw) \dots I_{bw\ stop}(bw) - D$

where D defines the distance between the bands where the cut-off characteristic should be checked, i.e. $D(bw)$.

- Then, it is possible to define a final information (bandwidth information or
- 20 bandwidth selection information) to be used to control a subgroup (e.g., TNS tool 33 and/or noise level estimation tool 36 and/or the TNS decoder tool 45 and/or noise filling tool 43). The final information may be, for example, encoded in some bits and may take the form of such as

$$P_{bw} = \begin{cases} bw, & \text{if } \sum F_C(b) > 0 \\ N_{bw} - 1, & \text{else} \end{cases}$$

The parameter bandwidth P_{bw} (bandwidth selection information) may be used to control the TNS and the noise filling tool, e.g., at the decoder and embody the signal 39a. The parameter P_{bw} may be stored and/or transmitted in a bitstream using the number of bits $nbits_{bw}$. Notably, the number of bits is not necessarily constant and may vary according to the chosen sample rate f_s , hence, reducing the payload for the bitstream where not necessary.

A table such as the following one may be used:

f_s	N_{bw}	$I_{bw\ start}$	$I_{bw\ stop}$	Bandwidth (P_{bw}) 39a	$nbits_{bw}$
8000	0	-	-	{NB}	0
16000	1	{53, 0, 0, 0}	{63, 0, 0, 0}	{NB, WB}	1
24000	2	{47, 59, 0, 0}	{56, 63, 0, 0}	{NB, WB, SSWB}	2
32000	3	{44, 54, 60, 0}	{52, 59, 63, 0}	{NB, WB, SSWB, SWB}	2
48000	4	{41, 51, 57, 61}	{49, 55, 60, 63}	{NB, WB, SSWB, SWB, FB}	3

Table 1

- 10 f_s is a given sampling rate (e.g., 8 KHz, 16 KHz, 24 KHz, 32 KHz, and/or 48 KHz) and, for each f_s , the number of possible modes is $N_{bw} + 1$.

Therefore, it is possible to 0 data encode a control data field including:

- 0 data bits corresponding to (signalling the choice of) NB bandwidth;
- 1 data bit corresponding to (signalling the choice of one of) NB and WB bandwidth;

- 2 data bits corresponding to (signalling the choice of one of) NB, WB, and SSWB bandwidth;
- 2 data bits corresponding to (signalling the choice of one of) NB, WB, SSWB, and SWB bandwidth;
- 5 - 3 data bits corresponding to (signalling the choice of one of) NB, WB, SSWB, SWB, and FB bandwidth.

An electronic version of at least some portions of Table 1 may be stored in the encoder and/or decoder. Accordingly, when the parameter bandwidth P_{bw} , it is possible to automatically know control information for the TNS and noise filling
 10 operations. For example, $I_{bw\ start}$ may refer to the start index associated to the lower end of the bandwidth $I_{bw\ stop}$ may refer to the final index associated to the higher end of the bandwidth. The bandwidth choice and parameters based on this choice may, therefore, derived from a table such as Table 1.

In examples, when $f_s = 8000$, the bandwidth detector is not needed and we have
 15 $P_{bw} = 0$ and $nbits_{bw} = 0$, i.e. the parameter P_{bw} is not placed in the bitstream. However, the decoder will understand that the chosen bandwidth is NB (e.g., on the basis of electronic instruments such as an electronic version of Table 1).

Other methods may be used. One of the bandwidths NB, WB, SSWB, SWB, FB may be identified and transmitted to the FD tools of the encoder subgroup, such
 20 as the TNS shaping tool 33 and the noise estimator tool 36. Information such as the parameter P_{bw} (39a) may be encoded and transmitted to the decoder apparatus 40 or 40a, so that the decoder noise estimator tool 43 and the TNS decoder tool 45 make use of the information regarding the selected bandwidth.

In general terms, the information signal characteristics which constitute the basis
 25 for the selection of the bandwidth may comprise, inter alia, one or more of the signal bandwidth, at least one energy estimation of the information signal, cut-off characteristics on the spectrum, information on the detection of quietness in some particular bands, $F_Q(bw)$, etc.

The examples above permit to obtain a soft bandwidth switching.

5.2 MDCT or MDST (or other transform) at the encoder

A modified discrete cosine transform (MDCT) or modified discrete sine transform (MDST) (or another modulated lapped transform) tool 31 may convert a digital representation in the TD into a digital representation in the FD. Other examples
 5 (maybe based on other transformations, such as lapped transformations) may be notwithstanding used. An example is provided here.

The input signal $x(n)$ of a current frame b in the TD may consist of N_F audio samples, where the newest one is located at $x(N_F - 1)$. Audio samples of past frames are accessed by negative indexing, e.g. $x(-1)$ is the newest of the
 10 previous frame.

The time input buffer for the MDCT t may be updated according to

- $t(n) = x(Z - N_F + n)$ for $n = 0..2N_F - 1 - Z$
- $t(2N_F - Z + n) = 0$ for $n = 0..Z - 1$ (initialization just required for consistency)

15 A block of N_F time samples may be transformed to the frequency coefficients $X(k)$ using the following equation:

$$X(k) = \sqrt{\frac{2}{N_F}} \sum_{n=0}^{2N_F-1} w_N(n) \cdot t(n) \cos \left[\frac{\pi}{N_F} \left(n + \frac{1}{2} + \frac{N_F}{2} \right) \left(k + \frac{1}{2} \right) \right] \text{ for } k = 0..N_F - 1$$

where w_N is the Low Delay MDCT window according to the used frame size. The window may be optimized for $N_F = 480$ and other versions for different frame sizes may be generated by means of interpolation. The window shape may be the result
 20 of an optimization procedure and may be provided point by point.

It is also possible to apply MDST or other transformations.

5.3.1 LPC at the encoder

A linear predictive coding (LPC) analysis may be performed by an LPC tool 32. LPC is a used representing the spectral envelope of a digital signal in compressed form, using the information of a linear predictive model.

5 An LPC filter may be derived in a warped frequency domain and therefore psychoacoustically optimized. To obtain the autocorrelation function, the Energy $E_B(b)$, as defined above, may be pre-emphasized by

$$E_{Pre}(b) = E_B(b) \cdot 10^{\frac{b \cdot g_{tilt}}{10(N_B-1)}} \text{ for } b = 0..N_B - 1$$

where

f_s	g_{tilt}
16000	18
24000	22
32000	26
48000	30

and transformed to time domain using, for example, an inverse odd DFT

$$R_{Pre}(n) = Re \left(\sum_{b=0}^{N_B-1} E_{Pre}(b) \cdot e^{j \frac{\pi \cdot n}{N_B} (b + \frac{1}{2})} \right) \text{ for } n = 0..N_B - 1$$

10

$$R_{Pre}(0) = R_{Pre}(0) * 1.0001$$

In case $R_{Pre}(0) = 0$, set $R_{Pre}(0) = 1$ and $R_{Pre}(1..N_B - 1) = 0$. The first N_L samples are extracted into the vector $R_L = R_{Pre}(0..N_L - 1)$, where N_L stands for the LP filter order, i.e. $N_L = 16$.

The LP filter coefficients may be calculated, for example, based on the vector R_L through the Levinson-Durbin procedure. This procedure may be described by the following pseudo code:

```

 $e = R_L(0)$ 
 $a^0(0) = 1$ 
for  $k = 1$  to  $N_L$  do
     $rc = \frac{-\sum_{n=0}^{k-1} a^{k-1}(n)R_L(k-n)}{e}$ 
5     $a^k(0) = 1$ 
    for  $n = 1$  to  $k - 1$  do
         $a^k(n) = a^{k-1}(n) + rc \cdot a^{k-1}(k - n)$ 
     $a^k(k) = rc$ 
     $e = (1 - rc^2)e$ 
10 with  $a(k) = a^{N_L}(k), k = 0..N_L$  are the estimated LPC coefficients and  $e$  is the
    prediction error.

```

The LPC coefficients may be weighted, in examples, by equation such as:

$$a_w(k) = a(k) \cdot 0.94^k \text{ for } k = 0..N_L$$

The LPC coefficients may be quantized.

For example, the weighted LPC coefficients $a_w(k)$ are first convolved with the

15 coefficients $b(i)$ using

$$a_c(k) = \sum_{i=0}^2 a'_w(k-i)b(i) \text{ for } k = 0..N_L + 2$$

with

$$a'_w(k) = \begin{cases} a_w(k) & \text{if } 0 \leq k \leq N_L \\ 0 & \text{otherwise} \end{cases}$$

and

$$b(i) = \begin{cases} \sum_{k=0}^{N_L} a_w(k) - \sum_{k=0}^{N_L} (-1)^k a_w(k) & \text{if } i = 0 \text{ or } i = 2 \\ -2 \left(\sum_{k=0}^{N_L} a_w(k) + \sum_{k=0}^{N_L} (-1)^k a_w(k) \right) & \text{if } i = 1 \end{cases}$$

The coefficients $a_c(k)$ may then be transformed to the frequency domain using

$$A(k) = \sum_{n=0}^{N_L+2} a_c(n) e^{\frac{-i2\pi k \left(n - \frac{N_L+1+2}{2} \right)}{N_T}} \text{ for } k = 0..N_T - 1$$

where $N_T = 256$ is the transform length. Note that this transform can be efficiently implemented using a pruned FFT. The real and imaginary parts of $A(k)$ are then extracted

$$A_r(k) = \text{Re}(A(k)) \text{ for } k = 0..\frac{N_T}{2}$$

$$A_i(k) = \text{Im}(A(k)) \text{ for } k = 0..\frac{N_T}{2}$$

- 5 LSFs may be obtained by a zero-crossing search of $A_r(k)$ and $A_i(k)$ that can be described with the following pseudo-code

```

specix = 1;
lsfix = 0;
while ((specix <= 128) && lsfix <= 15)
10 {
    while (specix <= 128 && Ar[specix-1]*Ar[specix] >= 0)
    {
        specix++;
    }
15 if (specix <= 128)
    {
        tmp = specix-1 + Ar[specix-1]/(Ar[specix-1]-Ar[specix]);
        lsf[lsfix++] = tmp/128;
    }
20 while (specix <= 128 && Ai[specix-1]*Ai[specix] >= 0)
    {
        specix++;
    }
    if (specix <= 128)
25 {
        tmp = specix-1 + Ai[specix-1]/(Ai[specix-1]-Ai[specix]);
        lsf[lsfix++] = tmp/128;
    }

```

```

    }
}

```

If less than 16 LSFs are found, the LSFs are set according to

$$\text{lsf}(k) = \frac{k+1}{N_L+1} \quad \text{for } k = 0..N_L-1$$

5 An LPC shaping may be performed in the MDCT or MDST (FD) domain by applying gain factors computed from the weighted and quantized LP filter coefficients transformed to the MDCT or MDST spectrum.

To compute $N_B = 64$ LPC shaping gains, weighted LP filter coefficients $\tilde{a}$ are first transformed into the frequency domain using an odd DFT.

$$G_{LPC}(b) = \sum_{k=0}^{N_L} \tilde{a}(k) \cdot e^{-j\frac{\pi k}{N_B}(b+\frac{1}{2})} \quad \text{for } b = 0..N_B-1$$

10 LPC shaping gains $g_{LPC}(b)$ may then be obtained as the absolute values of $G_{LPC}(b)$.

$$g_{LPC}(b) = |G_{LPC}(b)| \quad \text{for } b = 0..N_B-1$$

The LPC shaping gains $g_{LPC}(b)$ may be applied on the MDCT or MDST frequency lines for each band separately in order to generate the shaped spectrum $X_s(k)$ as outlined by the following code.

```

15 for b=0 to  $N_B-1$  do
    for  $k=I_{fs}(b)$  to  $I_{fs}(b+1)-1$ 
         $X_s(k) = X(k) \cdot g_{LPC}(b)$ 

```

As can be seen from above, the LPC tool, for performing the LPC analysis, is not controlled by the controller 39: for example, there is no selection of a particular bandwidth.

20 5.3.2 SNS at the encoder

With reference to Fig. 4a, it is possible to use a spectral noise shaper tool 32a.

Spectral noise shaping (SNS) shapes the quantization noise in the frequency domain such that it is minimally perceived by the human ear, maximizing the perceptual quality of the decoded output.

Spectral noise shaping may be performed using, for example, 16 scaling
5 parameters. These parameters may be obtained in the encoder by first computing
the energy of the MDCT (or MDST, or another transform) spectrum in 64 non-
uniform bands, then by applying some processing to the 64 energies (smoothing,
pre-emphasis, noise-floor, log-conversion), then by downsampling the 64
processed energies by a factor of 4 to obtain 16 parameters which are finally
10 normalized and scaled. These 16 parameters may be then quantized using vector.
The quantized parameters may then be interpolated to obtain 64 interpolated
scaling parameters. These 64 scaling parameters are then used to directly shape
the MDCT (or MDST...) spectrum in the 64 non-uniform bands. The scaled MDCT
(or MDST...) coefficients may then be quantized using a scalar quantizer with a
15 step size controlled by a global gain. At the decoder, inverse scaling is performed
in every 64 bands, shaping the quantization noise introduced by the scalar
quantizer. An SNS technique here disclosed may use, for example, only 16+1
parameters as side-information and the parameters can be efficiently encoded
with a low number of bits using vector quantization. Consequently, the number of
20 side-information bits is reduced, which may lead to a significant advantage at low
bitrate and/or low delay. A non-linear frequency scaling may be used. In this
examples, none of the LPC-related functions are used to reduce complexity. The
required processing functions (smoothing, pre-emphasis, noise-floor, log-
conversion, normalization, scaling, interpolation) need very small complexity in
25 comparison. Only the vector quantization still has relatively high complexity.
However, some low complexity vector quantization techniques can be used with
small loss in performance (multi-split/multi-stage approaches). This SNS technique
is not relying on a LPC-based perceptual filter. It uses 16 scaling parameters
which can be computed with a lot of freedom. Flexibility is therefore increased.

30 At the encoder 30a, the SNS tool 32 may perform at least one of the following
passages:

- Step 1: Energy per band

The energy per band $E_B(n)$ may be computed as follows

$$E_B(b) = \sum_{k=Ind(b)}^{Ind(b+1)-1} \frac{X(k)^2}{Ind(b+1)-Ind(b)} \quad for \quad b = 0 \dots N_B - 1$$

- with $X(k)$ are the MDCT (or MDST, or another transform) coefficients, $N_B = 64$ is the number of bands and $I_{fs}(n)$ are the band indices. The bands may be non-uniform and follow the perceptually-relevant bark scale (smaller in low-frequencies, larger in high-frequencies).

- Step 2: Smoothing

The energy per band $E_B(b)$ is smoothed using

$$E_S(b) = \begin{cases} 0.75 \cdot E_B(0) + 0.25 \cdot E_B(1) & , if \quad b = 0 \\ 0.25 \cdot E_B(62) + 0.75 \cdot E_B(63) & , if \quad b = 63 \\ 0.25 \cdot E_B(b-1) + 0.5 \cdot E_B(b) + 0.25 \cdot E_B(b+1) & , otherwise \end{cases}$$

- 10 This step may be mainly used to smooth the possible instabilities that can appear in the vector $E_B(b)$. If not smoothed, these instabilities are amplified when converted to log-domain (see step 5), especially in the valleys where the energy is close to 0.

- Step 3: Pre-emphasis

- 15 The smoothed energy per band $E_S(b)$ is then pre-emphasized using

$$E_P(b) = E_S(b) \cdot 10^{\frac{b \cdot g_{tilt}}{10 \cdot 63}} \quad for \quad b = 0 \dots 63$$

- with g_{tilt} controls the pre-emphasis tilt and depends on the sampling frequency. It is for example 18 at 16kHz and 30 at 48kHz. The pre-emphasis used in this step has the same purpose as the pre-emphasis used in the LPC-based perceptual filter of prior art 2, it increases the amplitude of the shaped Spectrum in the low-frequencies, resulting in reduced quantization noise in the low-frequencies.

- Step 4: Noise floor

A noise floor at -40dB is added to $E_P(b)$ using

$$E_P(b) = \max(E_P(b), \text{noiseFloor}) \quad \text{for } b = 0..63$$

with the noise floor being calculated by

$$\text{noiseFloor} = \max\left(\frac{\sum_{b=0}^{63} E_P(b)}{64} \cdot 10^{-\frac{40}{10}}, 2^{-32}\right)$$

This step improves quality of signals containing very high spectral dynamics such as e.g. glockenspiel, by limiting the amplitude amplification of the shaped spectrum in the valleys, which has the indirect effect of reducing the quantization noise in the peaks (an increase of quantization noise in the valleys is not perceptible).

- Step 5: Logarithm

A transformation into the logarithm domain is then performed using

$$E_L(b) = \frac{\log_2(E_P(b))}{2} \quad \text{for } b = 0..63$$

- 10 • Step 6: Downsampling

The vector $E_L(b)$ is then downsampled by a factor of 4 using

$$E_4(b) = \begin{cases} w(0)E_L(0) + \sum_{k=1}^5 w(k) E_L(4b + k - 1) & , \text{if } b = 0 \\ \sum_{k=0}^4 w(k) E_L(4b + k - 1) + w(5)E_L(63) & , \text{if } b = 15 \\ \sum_{k=0}^5 w(k) E_L(4b + k - 1) & , \text{otherwise} \end{cases}$$

with

$$w(k) = \left\{ \frac{1}{12}, \frac{2}{12}, \frac{3}{12}, \frac{3}{12}, \frac{2}{12}, \frac{1}{12} \right\}$$

This step applies a low-pass filter ($w(k)$) on the vector $E_L(b)$ before decimation.

This low-pass filter has a similar effect as the spreading function used in

psychoacoustic models: it reduces the quantization noise at the peaks, at the cost of an increase of quantization noise around the peaks where it is anyway perceptually masked.

- Step 7: Mean Removal and Scaling

- 5 The final scale factors are obtained after mean removal and scaling by a factor of 0.85

$$scf(n) = 0.85 \left(E_4(n) - \frac{\sum_{b=0}^{15} E_4(b)}{16} \right) \text{ for } n = 0..15$$

- Since the codec has an additional global-gain, the mean can be removed without any loss of information. Removing the mean also allows more efficient vector quantization. The scaling of 0.85 slightly compress the amplitude of the noise
- 10 shaping curve. It has a similar perceptual effect as the spreading function mentioned in Step 6: reduced quantization noise at the peaks and increased quantization noise in the valleys.

- Step 8: Quantization

- The scale factors are quantized using vector quantization, producing indices which
- 15 are then packed into the bitstream and sent to the decoder, and quantized scale factors $scfQ(n)$.

- Step 9: Interpolation

The quantized scale factors $scfQ(n)$ are interpolated using:

$$\begin{aligned}
scfQint(0) &= scfQ(0) \\
scfQint(1) &= scfQ(0) \\
scfQint(4n+2) &= scfQ(n) + \frac{1}{8}(scfQ(n+1) - scfQ(n)) \quad \text{for } n = 0..14 \\
scfQint(4n+3) &= scfQ(n) + \frac{3}{8}(scfQ(n+1) - scfQ(n)) \quad \text{for } n = 0..14 \\
scfQint(4n+4) &= scfQ(n) + \frac{5}{8}(scfQ(n+1) - scfQ(n)) \quad \text{for } n = 0..14 \\
scfQint(4n+5) &= scfQ(n) + \frac{7}{8}(scfQ(n+1) - scfQ(n)) \quad \text{for } n = 0..14 \\
scfQint(62) &= scfQ(15) + \frac{1}{8}(scfQ(15) - scfQ(14)) \\
scfQint(63) &= scfQ(15) + \frac{3}{8}(scfQ(15) - scfQ(14))\#
\end{aligned}$$

and transformed back into linear domain using

$$g_{SNS}(b) = 2^{scfQint(b)} \quad \text{for } b = 0..63$$

Interpolation may be used to get a smooth noise shaping curve and thus to avoid any big amplitude jumps between adjacent bands.

- Step 10: Spectral Shaping

- 5 The SNS scale factors $g_{SNS}(b)$ are applied on the MDCT (or MDST, or another transform) frequency lines for each band separately in order to generate the shaped spectrum $X_s(k)$

$$X_s(k) = \frac{X(k)}{g_{SNS}(b)} \quad \text{for } k = I_{fs}(b)..I_{fs}(b+1) - 1, \quad \text{for } b = 0..63$$

5.4 TNS at the encoder

- Fig. 7 shows a method 70 indicating operations of a TNS tool such as the TNS tool
10 33 of the encoder 30 or 30a.

At step S71, selection information regarding the selected bandwidth (e.g., parameter P_{bw}) may be obtained from the encoder bandwidth detector and controller 39, for example.

According to the selection information (bandwidth information), the behaviour of the TNS is different for different bandwidths (NB, WB, SSWB, SWB, FB). An example is provided by the following table:

Band-width	num_tns_filters	start_freq(f)	stop_freq(f)	sub_start(f,s)	sub_stop(f,s)
NB	1	{12}	{80}	{{12,34,57}}	{{34,57,80}}
WB	1	{12}	{160}	{{12,61,110}}	{{61,110,160}}
SSWB	1	{12}	{240}	{{12,88,164}}	{{88,164,240}}
SWB	2	{12,160}	{160,320}	{{12,61,110}, {160,213,266}}	{{61,110,160}, {213,266,320}}
FB	2	{12,200}	{200,400}	{{12,74,137}, {200,266,333}}	{{74,137,200}, {266,333,400}}

Table 2

- 5 For example, when the selection information is SWB, the TNS will perform a filtering twice (see num_tns_filters). As can be seen from the tables, different indexes are associated to different bandwidths (e.g., for NB the stop frequency is different than for WB, and so on).

Therefore, as can be seen, the TNS tool 33 may operate at a different bandwidth
 10 on the basis of the selection set out by the controller 39. Notably, other FD tools of the same encoder apparatus 40 or 40a may continue perform processes at a different frequency.

The TNS encoding steps are described below. First, an analysis estimates a set of reflection coefficients for each TNS Filter (step S72). Then, these reflection
 15 coefficients are quantized (step S73). And finally, the MDCT- or MDST-spectrum is filtered using the quantized reflection coefficients (step S73).

With reference to the step S72, a complete TNS analysis described below may be repeated for every TNS filter f , with $f = 0..num_tns_filters-1$ (num_filters is given

in Table 2). Other TNS analysis operations may be performed, which provide reflection coefficients.

The TNS tool may be configured to perform an autocorrelation on a TNS input value. A normalized autocorrelation function may be calculated as follows, for

5 each $k = 0..8$ (for example)

$$r(k) = \begin{cases} r_0(k) & , \text{ if } \prod_{s=0}^2 e(s) = 0 \\ \sum_{s=0}^2 \frac{\sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1-k} X_s(n)X_s(n+k)}{e(s)} & , \text{ otherwise} \end{cases}$$

with

$$r_0(k) = \begin{cases} 1 & , \text{ if } k = 0 \\ 0 & , \text{ otherwise} \end{cases}$$

and

$$e(s) = \sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1} X_s(n)^2 \quad \text{for } s = 0..2$$

10 with $\text{sub_start}(f, s)$ and $\text{sub_stop}(f, s)$ given Table 2. $e(s)$ is an energy sum over a spectral subsection (a normalization factor between the start and the stop frequency of each filter).

The normalized autocorrelation function may be lag-windowed using, for example:

$$r(k) = r(k) \exp \left[-\frac{1}{2} (0.02\pi k)^2 \right] \quad \text{for } k = 0..8$$

In some examples, it is possible to perform a decision to turn on/off the TNS filter f in the current frame is based on the prediction gain

15 If $\text{predGain} > \text{thresh}$, then turn on the TNS filter f

with $\text{thresh} = 1.5$ and the prediction gain may be computed by

$$\text{predGain} = \frac{r(0)}{e}$$

The additional steps described below are performed only if the TNS filter f is turned on (or in the example which do not use the turning on/off).

In some examples, a weighting factor may be computed by

$$\gamma = \begin{cases} 1 - (1 - \gamma_{\min}) \frac{\text{thresh2} - \text{predGain}}{\text{thresh2} - \text{thresh}} & , \text{ if tns_lpc_weighting} = 1 \text{ and } \text{predGain} < \text{thresh2} \\ 1 & , \text{ otherwise} \end{cases}$$

5 with $\text{thresh2} = 2$, $\gamma_{\min} = 0.85$ and

$$\text{tns_lpc_weighting} = \begin{cases} 1 & , \text{ if nbits} < 480 \\ 0 & , \text{ otherwise} \end{cases}$$

The LPC coefficients may be weighted using the factor γ

$$a_w(k) = \gamma^k a(k) \quad \text{for } k = 0..8$$

The weighted LPC coefficients may be converted to reflection coefficients using the following procedure:

$$\begin{aligned} & a^K(k) = a_w(k), k = 0, \dots, K \\ & \text{for } k = K \text{ to } 1 \text{ do} \\ & \quad \text{rc}(k) = a^k(k) \\ 10 \quad & \quad e = (1 - \text{rc}(k)^2) \\ & \quad \text{for } n = 1 \text{ to } k - 1 \text{ do} \\ & \quad \quad a^{k-1}(n) = \frac{a^k(n) - \text{rc}(k)a^k(k-n)}{e} \end{aligned}$$

wherein $\text{rc}(k, f) = \text{rc}(k)$ are the final estimated reflection coefficients for the TNS filter f .

15 If the TNS filter f is turned off, then the reflection coefficients may be simply set to 0: $\text{rc}(k, f) = 0$, $k = 0..8$.

At step S73, a quantization step may be performed. For example, for each TNS filter f , reflection coefficients (e.g., as obtained at step S72) may quantized. For example, scalar uniform quantization in the arcsine domain may be used:

$$rc_i(k, f) = \text{nint} \left[\frac{\arcsin(rc(k, f))}{\Delta} \right] + 8 \quad \text{for } k = 0..8$$

and/or

$$rc_q(k, f) = \sin[\Delta(rc_i(k, f) - 8)] \quad \text{for } k = 0..8$$

- 5 with $\Delta = \frac{\pi}{17}$ and $\text{nint}(\cdot)$ being the rounding-to-nearest-integer function, for example; $rc_i(k, f)$ the quantizer output indices; and $rc_q(k, f)$ the quantized reflection coefficients.

An order of the quantized reflection coefficients may be calculated using

$$k = 7$$

while $k \geq 0$ and $rc_q(k, f) = 0$ do

$$10 \quad k = k - 1$$

$$rc_{order}(f) = k + 1$$

A total number of bits consumed by TNS in the current frame may be computed as follows

$$nbits_{TNS} = \left\lceil \sum_{f=0}^{\text{num_tns_filters}-1} \frac{2048 + nbits_{TNS_{order}}(f) + nbits_{TNS_{rc}}(f)}{2048} \right\rceil$$

with

$$nbits_{TNS_{order}}(f) = \begin{cases} \text{ac_tns_order_bits}[\text{tns_lpc_weighting}][rc_{order}(f) - 1] & , \text{if } rc_{order}(f) > 0 \\ 0 & , \text{otherwise} \end{cases}$$

and/or

$$nbits_{TNS_{coef}}(f) = \begin{cases} \sum_{k=0}^{rc_{order}(f)-1} ac_tns_coef_bits[k][rc_i(k, f)] & , if \ rc_{order}(f) > 0 \\ 0 & , otherwise \end{cases}$$

[...] means a rounding operation to the integer over.

The tables tab_nbits_TNS_order and tab_nbits_TNS_coef may be pre-defined.

At step S74, a digital representation of an information signal in the FD (e.g., as provided by the LPC tool 32 or SNS tool 32a) may be filtered. This representation
 5 may be, in examples, in the form of a modified discrete cosine or sine transform (MDCT or MDST). The MDCT spectrum $X_s(n)$ may filtered using the following algorithm, for example:

$$s^0(\text{start_freq}(0) - 1) = s^1(\text{start_freq}(0) - 1) = \dots = s^7(\text{start_freq}(0) - 1) = 0$$

for $f = 0$ to num_tns_filters-1 do

 for $n = \text{start_freq}(f)$ to stop_freq(f) - 1 do

$$t^0(n) = s^0(n) = X_s(n)$$

10 for $k = 0$ to 7 do

$$t^{k+1}(n) = t^k(n) + rc_q(k)s^k(n-1)$$

$$s^{k+1}(n) = rc_q(k)t^k(n) + s^k(n-1)$$

$$X_f(n) = t^8(n)$$

where $X_f(n)$ is the TNS filtered MDCT or MDST spectrum.

15 Other filtering techniques may be used. However, it may be seen that the TNS is applied to the particular bandwidth (e.g., NB, WB, SSWB, SWB, FB) chosen by the controller 39 on the basis of the signal characteristics.

5.5 Spectral quantization at the encoder

A spectrum quantizer tool 34 is here discussed. The MDCT or MDST spectrum
 20 after TNS filtering ($X_f(n)$) may be quantized using dead-zone plus uniform threshold scalar quantization and the quantized MDCT or MDST spectrum $X_q(n)$ may then be encoded using arithmetic encoding. A global gain gg may control the

step size of the quantizer. This global gain is quantized with 7 bits and the quantized global gain index gg_{ind} is then an integer, for example, between 0 and 127. The global gain index may be chosen such that the number of bits needed to encode the quantized MDCT or MDST spectrum is as close as possible to the available bit budget.

In one example, a number of bits available for coding the spectrum may be given by

$$nbits_{spec} = nbits - nbits_{bw} - nbits_{TNS} - nbits_{LTPF} - nbits_{LPC/SNS} - nbits_{gain} - nbits_{nf} - \left\lceil \log_2 \left(\frac{N_E}{2} \right) \right\rceil - 1$$

with $nbits$ being the number of bits available in one TD frame for the original information signal, $nbits_{bw}$ provided in Table 1, $nbits_{TNS}$ provided by the TNS (total number of bits consumed by TNS in a current frame), $nbits_{LTPF}$ being associated to the LTPF 38b (number of bits consumed by LTPF), $nbits_{LPC/SNS} = 38$, $nbits_{gain} = 7$ and $nbits_{nf} = 3$, for example. In examples, also protection bits (e.g., cyclical redundancy code, CRC, bits) may be taken into consideration.

An offset may first be computed using

$$nbits_{offset} = 0.8 * nbits_{offset}^{old} + 0.2 * \min(40, \max(-40, nbits_{offset}^{old} + nbits_{spec}^{old} - nbits_{est}^{old}))$$

with $nbits_{offset}^{old}$ is the value of $nbits_{offset}$ in the previous frame, $nbits_{spec}^{old}$ is the value of $nbits_{spec}$ in the previous frame and $nbits_{est}^{old}$ is the value of $nbits_{est}$ in the previous frame.

This offset may then be used to adjust the number of bits available for coding the spectrum

$$nbits_{spec} = \text{nint}(nbits_{spec} + nbits_{offset})$$

A global gain index may then be estimated such that the number of bits needed to encode the quantized MDCT or MDST spectrum is as close as possible to the available bit budget. This estimation is based on a low-complexity bisection search which coarsely approximates the number of bits needed to encode the quantized spectrum. The algorithm can be described as follows

```

fac = 128;
ggind = 127;
for (iter = 0; iter < 7; iter++)
{
    fac >>= 1;
    ggind -= fac;
    tmp = 0;
    for (i = 0; i < NE/4; i++)
    {
        if (E[i]*28/20 < ggind)
        {
            tmp += 2.7*28/20;
        }
        else
        {
            tmp += E[i]*28/20 - ggind + 7*28/20;
        }
    }
    if (tmp > nbbitsspec*1.4*28/20)
    {
        ggind += fac;
    }
}

```

with E[k] being the energy (in dB) of blocks of 4 MDCT or MDST coefficients given
by

$$E(k) = 10 * \log_{10} \left(\sum_{n=0}^3 X_f(4 * k + n)^2 \right) \quad \text{for } k = 0.. \frac{N_E}{4}$$

The global gain index above is first unquantized using

$$gg = 10^{\frac{gg_{ind}}{28}}$$

The spectrum X_f may then be quantized using, for example:

$$X_q(n) = \begin{cases} \min \left(\left\lfloor \frac{X_f(n)}{gg} + 0.375 \right\rfloor, 32767 \right) & , \text{ if } X_f(n) \geq 0 \\ \max \left(\left\lfloor \frac{X_f(n)}{gg} - 0.375 \right\rfloor, -32768 \right) & , \text{ otherwise} \end{cases}$$

The number of bits nb_{bits}_{est} needed to encode the quantized MDCT or MDST (or,
anyway, FD) spectrum $X_q(n)$ can be accurately estimated using the algorithm
below.

A bitrate flag is first computed using, for example:

```

get_rateFlag( $f_s$ , nbits)
if (nbits > (160 + min(4, ( $f_s$ /8000-1)) * 160))
{
5       rateFlag = 512;
}
else
{
10      rateFlag = 0;
}
return rateFlag;

```

Then the index of the last non-zeroed 2-tuple is obtained by

```

get_lastnz( $X_q$ [],  $N_E$ )
lastnz =  $N_E$ ;
15 while (lastnz>2 &&  $X_q$ [lastnz-1] == 0 &&  $X_q$ [lastnz-2] == 0)
{
    lastnz -= 2;
}
return lastnz;

```

20 The number of bits $nbits_{est}$ may be then computed as follows

```

 $nbits_{est}$  = 0;
c = 0;
for (n = 0; n < lastnz; n++)
{
25     t = c + rateFlag;
    if (k >  $N_E$ /2)
    {
        t += 256;
    }
30     a = abs( $X_q$ [k]);
    b = abs( $X_q$ [k+1]);
     $nbits_{est}$  += (min(a,1) + min(b,1)) * 2048;
    lev = 0;
    while (max(a,b) >= 4)
35     {
        pki = ac_spec_lookup[t+lev*1024];
         $nbits_{est}$  += 2*2048 + ac_spec_bits[pki][16];
        a >>= 1;
        b >>= 1;
40         lev = min(lev+1,3);
    }
    pki = ac_spec_lookup[t+lev*1024];
    sym = a + 4*b;
     $nbits_{est}$  += ac_spec_bits[pki][sym];
45     if (lev <= 1)

```

```

    {
        t = 1 + (a+b)*(lev+1);
    }
    else
5      {
        t = 12 + lev;
      }
      c = (c&15)*16 + t;
    }
10  nbitsest = ceil(nbitsest/2048);

```

with `ac_lookup` and `ac_bits` are tables which may be predefined.

The number of bits $nbits_{est}$ may be compared with the available bit budget $nbits_{spec}$. If they are far from each other, then the quantized global gain index gg_{ind} is adjusted and the spectrum is requantized. A procedure used to adjust the

15 quantized global gain index gg_{ind} is given below

```

if ((ggind < 127 && nbitsest > nbitsspec) ||
    (ggind > 0 && nbitsest < nbitsspec - 20))
{
    if (nbitsest < nbitsspec - 20)
20    {
        ggind -= 1;
    }
    else if (ggind == 126 || nbitsest < nbitsspec + 20)
    {
25        ggind += 1;
    }
    else
    {
30        ggind += 2;
    }
}

```

As can see from above, the spectral quantization is not controlled by the controller

39: there is no restriction to a particular band.

5.6 Entropy coding

All or part of the encoded data (TNS data, LTPF data, global gain, quantized spectrum...) may be entropy coded, e.g., by compression according to any algorithm.

- 5 A portion of this data may be composed by pure bits which are directly put in the bitstream starting from the end of the bitstream and going backward.

The rest of data may be encoded using arithmetic encoding starting from the beginning of the bitstream and going forward.

- The two data fields above may be exchanged regarding starting point and
- 10 direction of reading/writing of the bit stream.

An example in pseudo code may be:

```

bp = 0;
bp_side = nbytes - 1;
mask_side = 1;
15 nbits_written = 2 << 11;
c = 0;
lastnz = get_lastnz( $X_q$ ,  $N_E$ );
rateFlag = get_rateFlag( $f_s$ , nbits);
/* Bandwidth */
20 if (nbitsbw > 0)
{
    write_uint_backward(bytes, &bp_side, &mask_side,  $P_{bw}$ , nbitsbw);
    nbits_written += nbitsbw << 11;
}
25 /* Global Gain */
write_uint_backward(bytes, &bp_side, &mask_side,  $gg_{ind}$ , 7);
nbits_written += 7 << 11;
30 /* Noise Factor */
write_uint_backward(bytes, &bp_side, &mask_side, 0, 3);
nbits_written += 3 << 11;

/* TNS activation flag */
35 for (f = 0; f < num_tns_filters; f++)
{
    write_bit_backward(bytes, &bp_side, &mask_side, min( $rc_{order}(f)$ , 1));
    nbits_written += 1 << 11;
}
40

```

```

/* LTPF data */
write_bit_backward(bytes, &bp_side, &mask_side, pitch_present);
nbits_written += 1 << 11;
if (pitch_present != 0)
5  {
    write_uint_backward(bytes, &bp_side, &mask_side, pitch_index, 9);
    write_uint_backward(bytes, &bp_side, &mask_side, ltpf_active, 1);
    nbits_written += 10 << 11;
}
10
/* Env-VQ integer bits */
write_uint_backward(bytes, &bp_side, &mask_side, L_lsf_idx[0], 10 >> 1);
write_uint_backward(bytes, &bp_side, &mask_side, L_lsf_idx[1], 10 >> 1);
write_bit_backward(bytes, &bp_side, &mask_side, lsf_submode_flag);
15 write_uint_backward(bytes, &bp_side, &mask_side, L_lsf_idx[3], fgBits);
write_bit_backward(bytes, &bp_side, &mask_side, L_lsf_idx[4]);
nbits_written += (12 + fgBits) << 11;

/* Last non-zero tuple */
20 nbits_lastnz = ceil(log2(NE/2));
bp_side_lastnz = bp_side;
mask_side_lastnz = mask_side;
write_uint_backward(bytes, &bp_side_lastnz, &mask_side_lastnz, (lastnz >>
1) - 1, nbits_lastnz);
25 nbits_written += nbits_lastnz << 11;

```

5.7 Noise estimation at the encoder

A noise estimation tool 36 (noise level estimator) may control the noise filing on the decoder side. At the encoder side, the noise level parameter may be
30 estimated, quantized and transmitted or stored in a bitstream.

The noise level may be estimated based on the spectral coefficients which have been quantized to zero, i.e. $X_q(k) == 0$. The indices for the relevant spectral coefficients are given by

$$I_{NF}(k) = \begin{cases} 1 & \text{if } 24 \leq k < bw_{stop} \text{ and } X_q(i) == 0 \text{ for all } i = k - 3.. \min(bw_{stop}, k + 3) \\ 0 & \text{otherwise} \end{cases}$$

where bw_{stop} may depend on the bandwidth detected at step S62 and/or by the
35 bandwidth detector and controller 39 as defined, for example, in the following table:

	Bandwidth(P_{bw} , 39a)				
	NB	WB	SSWB	SWB	FB
bw_{stop}	80	160	240	320	400

Table 3

For the identified indices, the mean level of missing coefficients is estimated based on the spectrum after TNS filtering ($X_f(k)$), for example, and normalized by the global gain.

$$L_{NF} = \frac{\sum_{k=0}^{N_E-1} I_{NF}(k) \cdot \frac{|X_f(k)|}{gg}}{\sum_{k=0}^{N_E-1} I_{NF}(k)}$$

The final noise level may be quantized to eight steps:

$$F_{NF} = \min(\max([8 - 16 \cdot L_{NF}], 0), 7)$$

- 5 Therefore, the noise level estimator tool 36 may be controlled by the controller 39, e.g., on the basis of bandwidth information 39a.

For example, an electronic version of Table 3 may be stored in a storage unit so that, when the bandwidth selection for a particular bandwidth is obtained, the parameter bw_{stop} is easily derived.

10 5.8 Entropy decoding at the decoder

All the encoded data (TNS data, LTPF data, global gain, quantized spectrum...) may be entropy decoded at the decoder side, e.g., using the decoder tool 42. A bitstream provided by an encoder may, therefore, be decompressed according to any algorithm.

15 5.9 Noise filling at the decoder

A decoder noise filling tool 43 is here discussed. The decoder noise filling tool 43 may be controlled, inter alia, by the decoder bandwidth controller 49 (and/or by the

controller 39 via information 39a encoded in the bitstream, such as the control data field N_{bw} and/or P_{wb} of Table 1).

The indices for the relevant spectral coefficients may be given by

$$I_{NF}(k) = \begin{cases} 1 & \text{if } 24 \leq k < bw_{stop} \text{ and } \widehat{X}_q(i) == 0 \text{ for all } i = k - 3.. \min(bw_{stop}, k + 3) \\ 0 & \text{otherwise} \end{cases}$$

where bw_{stop} may be given in Table 3.

- 5 The noise filling may be applied on the identified relevant spectral lines $I_{NF}(k)$ using a transmitted noise factor F_{NF} obtained from the encoder. F_{NF} may be calculated at the noise estimator on encoder side. F_{NF} may be a 3 bit value coded as side information in the bit stream. F_{NF} may be obtained, for example, using the following procedure:

```

10  /* Bandwidth */
    if (nbitsbw > 0)
    {
        Pbw = read_uint(bytes, &bp_side, &mask_side, nbitsbw);
    }
15  else
    {
        Pbw = 0;
    }

20  /* Global Gain */
    ggind = read_uint(bytes, &bp_side, &mask_side, 7);

    /* Noise Level */
    FNF = read_uint(bytes, &bp_side, &mask_side, 3);
25  /* TNS activation flag */
    if (Pbw < 3)
    {
        num_tns_filters = 1;
    }
30  else
    {
        num_tns_filters = 2;
    }

35  for (f = 0; f < num_tns_filters; f++)
    {
        rcorder(f) = read_bit(bytes, &bp_side, &mask_side);
    }

```

```

/* LTPF data */
pitch_present = read_bit(bytes, &bp_side, &mask_side);
if (pitch_present != 0)
{
5   pitch_index = read_uint(bytes, &bp_side, &mask_side, 9);
    ltpf_active = read_uint(bytes, &bp_side, &mask_side, 1);
}
else
{
10   pitch_index = 0;
    ltpf_active = 0;
}

/* LSF-VQ integer bits */
15 L_lsf_idx[0] = read_uint(bytes, &bp_side, &mask_side, 10 >> 1);
   L_lsf_idx[1] = read_uint(bytes, &bp_side, &mask_side, 10 >> 1);
   lsf_submode_flag = read_bit(bytes, &bp_side, &mask_side);
   L_lsf_idx[3] = read_uint(bytes, &bp_side, &mask_side, fgBits);
20   L_lsf_idx[4] = read_bit(bytes, &bp_side, &mask_side);

/* Last non-zero tuple */
nbits_lastnz = ceil(log2( $N_E/2$ ));
lastnz = read_uint(bytes, &bp_side, &mask_side, nbits_lastnz);
lastnz = (lastnz + 1) << 1;

```

25 A procedure is here provided:

```

 $\widehat{L}_{NF} = (8 - F_{NF}) / 16;$ 
for k=0.. $bw_{stop} - 1$ 
  if  $I_{NF}(k) == 1$ 
30   nf_seed = (13849 + nf_seed * 31821) & 0xFFFF;
    if nf_seed >= 0x8000
       $\widehat{X}_q(k) = \widehat{L}_{NF};$ 
    else
       $\widehat{X}_q(k) = -\widehat{L}_{NF};$ 

```

How to obtain the nf_seed may be described, for example, by the following
35 pseudocode:

```

{
   $\widehat{X}_q[k] = 0;$ 
}

40 /* Noise Filling Seed */
tmp = 0;
for (k = 0; k <  $N_E$ ; k++)
{
  tmp += abs( $\widehat{X}_q[k]$ ) * k;
45 }
nf_seed = tmp & 0xFFFF;

```

As can be seen from above, the decoder noise filter tool 43 may make use of the parameter bw_{stop} .

In some examples, the parameter bw_{stop} explicitly obtained as a value in the bitstream. In examples, the parameter bw_{stop} is obtained by the controller 49 on the basis of the bandwidth information 39a (P_{bw}) in a control field of the bitstream encoded by the encoder. The decoder may have an electronic version of Table 3 stored in a non-transitory storage unit. Accordingly, the bitstream length is reduced.

Therefore, the bandwidth controller 49 (and/or the bandwidth detector and controller 39 of the decoder via the control data 39a) may control the decoder noise filling tool 43.

5.9 Global gain at the decoder

A global gain may be applied on the spectrum after the noise filling has been applied using, for example, a formula such as

$$\widehat{X}_f(k) = \widehat{X}_q(k) \cdot 10^{\left(\frac{gg_{ind}}{28}\right)} \quad \text{for } k = 0..N_E - 1$$

where gg_{ind} is a global gain index, e.g., obtained from the encoder.

5.10 TNS at the decoder

A TNS decoder tool 45 is here discussed. The quantized reflection coefficients may be obtained for each TNS filter f using

$$rc_q(k, f) = \sin[\Delta(rc_i(k, f) - 8)] \quad k = 0..8$$

where $rc_i(k, f)$ are the quantizer output indices.

The MDCT or MDST spectrum $\widehat{X}_f(n)$ (e.g., as generated by the global gain tool) may then be filtered using a following procedure such as:

```

s0(start_freq(0) - 1) = s1(start_freq(0) - 1) = ... = s7(start_freq(0) - 1) = 0
for f = 0 to num_tns_filters-1 do
    for n = start_freq(f) to stop_freq(f) - 1 do
        tK(n) =  $\widehat{X}_f(n)$ 

```

```

for k = 7 to 0 do
    tk(n) = tk+1(n) - rcq(k)sk(n - 1)
    sk+1(n) = rcq(k)tk(n) + sk(n - 1)
Xŝ(n) = s0(n) = t0(n)

```

5 where $\widehat{X}_s(n)$ is the output of the TNS decoder.

The parameters num_tns_filters, start_freq and stop_freq may be provided, on the basis of control information provided by the encoder.

In some examples num_tns_filters, start_freq and/or stop_freq are not explicitly provided in the bitstream. In examples, num_tns_filters, start_freq and stop_freq
10 are derived on the basis of the N_{bw} value in a control field of the bitstream encoded by the encoder. For example, the decoder may have an electronic version of Table 2 (or at least a portion thereof) stored therein. Accordingly, the bitstream length is reduced.

Therefore, the TNS decoder tool 45 may be controlled by the bandwidth detected
15 at the encoder side.

5.11.1 MDCT or MDST shaping at the decoder

An MDCT or MDST shaping tool 46 is here discussed. The LPC or SNS shaping may be performed in the MDCT (FD) domain by applying gain factors computed from the decoded LP filter coefficients transformed to the MDCT or MDST
20 spectrum.

To compute the N_B LPC shaping gains, the decoded LP filter coefficients $\tilde{a}$ may be first transformed into the frequency domain using an odd DFT.

$$G_{LPC}(b) = \sum_{k=0}^{N_L} \tilde{a}(k) \cdot e^{-j\frac{\pi k}{N_B}(b+\frac{1}{2})} \quad \text{for } b = 0..N_B - 1$$

The LPC shaping gains $g_{LPC}(b)$ may then be computed as the reciprocal absolute values of $G_{LPC}(b)$.

$$g_{LPC}(b) = \frac{1}{|G_{LPC}(b)|} \quad \text{for } b = 0..N_B - 1$$

The LPC shaping gains $g_{LPC}(b)$ may be applied on the TNS filtered MDCT frequency lines for each band separately as outlined in order to generate the shaped spectrum $\hat{X}(k)$ as outlined, for example, by the following code:

```

5   for (b=0; b<NB; b++) {
      for (k=Ifs(b); k<Ifs(b+1); k++) {
           $\hat{X}(k) = \hat{X}_f(k) \cdot g_{LPC}(b)$ 
      }
  }

```

As can be seen above, the MDCT or MDST shaping tool 46 does not need to be restricted to a particular bandwidth and, therefore, does not need to be controlled by the controller 49 or 39.

5.11.2 SNS at the decoder

The following steps may be performed at the noise shaper decoder, SNS, tool 46a:

- 15 • Step 1: Quantization

The vector quantizer indices produced in encoder step 8 (see section 5.3.2) are read from the bitstream and used to decode the quantized scale factors $scfQ(n)$.

- Step 2: Interpolation

Same as Step 9 at section 5.3.2.

- 20 • Step 3: Spectral Shaping

The SNS scale factors $g_{SNS}(b)$ are applied on the quantized MDCT (or MDST, or another transform) frequency lines for each band separately in order to generate the decoded spectrum $\hat{X}(k)$ as outlined by the following code.

$$\hat{X}(k) = \hat{X}_S(k) \cdot g_{SNS}(b) \text{ for } k = I_{f_s}(b) \dots I_{f_s}(b+1) - 1, \text{ for } b = 0 \dots 63$$

5.12 MDCT or MDST synthesis at the decoder

- 25 An inverse MDCT or MDST tool 48a is here discussed (other tools based on other transformations, such as lapped transformations, may be used).

A reconstructed spectrum $\hat{X}(k)$ may be transformed to time domain by the following steps:

1. Generation of time domain aliasing buffer $\hat{t}(n)$

$$\hat{t}(n) = \sqrt{\frac{2}{N_F}} \sum_{k=0}^{N_F-1} \hat{X}(k) \cos \left[\frac{\pi}{N_F} \left(n + \frac{1}{2} + \frac{N_F}{2} \right) \left(k + \frac{1}{2} \right) \right] \text{ for } n = 0..2N_F - 1$$

2. Windowing of time-aliased buffer

$$\hat{t}(n) = w_N(2N - 1 - n) \cdot \hat{t}(n) \text{ for } n = 0..2N_F - 1$$

3. Conduct overlap-add operation to get reconstructed time samples $\hat{x}(n)$

$$\hat{x}(n) = \text{mem_ola_add}(n) + \hat{t}(Z + n) \text{ for } n = 0..N_F - Z - 1$$

$$\hat{x}(n) = \hat{t}(Z + n) \text{ for } n = N_F - Z..N_F - 1$$

$$\text{mem_ola_add}(n) = \hat{t}(N_F + Z + n) \text{ for } n = 0..N_F - Z - 1$$

with $\text{mem_ola_add}(n)$ is initialized to 0 before decoding the first frame.

With reference to step 1, an MDST may be performed by exchanging the cos function by a sine function, e.g., to have:

$$\hat{t}(n) = \sqrt{\frac{2}{N_F}} \sum_{k=0}^{N_F-1} \hat{X}(k) \sin \left[\frac{\pi}{N_F} \left(n + \frac{1}{2} + \frac{N_F}{2} \right) \left(k + \frac{1}{2} \right) \right] \text{ for } n = 0..2N_F - 1$$

- 10 As can be seen above, the inverse MDCT or MDST tool 48a is not controlled on the basis of the bandwidth determined at the encoder side.

6. Other examples

- 15 Fig. 8a shows an apparatus 110 which may implement at least some tools of the encoder apparatus 30 or 30a and/or perform at least some steps of the method 60 and/or 70. The apparatus 110 may comprise a processor 111 and a non-transitory memory unit 112 storing instructions which, when executed by the processor 111, may cause the processor 111 to implement at least one of the TD and/or FD tools

of the encoder apparatus 30 or 30a. In particular, the instructions may implement a subgroup of FD tools (e.g., TNS and/or noise filling) and other FD tools which are not in the subgroup (e.g., 31, 32, 34, 35). The instructions may also comprise instructions which, when executed by the processor 111, perform a selection of the bandwidth so that the bandwidth of the signals processed by the tools in the subgroup of FD tools (e.g., TNS and/or noise filling) differs from the bandwidth of the signals processed by the other FD tools which are not in the subgroup (e.g., 31, 32, 34, 35). The instructions may be such as to control the bandwidth selection based on energy detections associated to the different bandwidths. The instructions may also comprise instructions which, when executed by the processor 111, permit to control a decoder and, in particular, permit to control the bandwidth of a subgroup of FD tools (e.g., 43, 45) which may be different from the bandwidth of other FD tools. The bandwidth chosen for the subgroup at the encoder may be the same chosen for the subgroup at the decoder. The non-transitory memory unit 112 may also comprise other data, such as at least portions electronic versions of Tables 1, 2, and/or 3. The apparatus 110 may comprise a storage space 118 for storing, for example, a bitstream obtained from an information signal (e.g., an audio signal). The apparatus 110 may comprise an output unit 117 for transmitting data, e.g., wirelessly, e.g., using a particular protocol, such as Bluetooth. For example, the apparatus 110 may define, by executing the instructions stored in the non-transitory memory unit 112, a bitstream to be transmitted to a decoder. The apparatus 110 may also comprise an input unit 116 for obtaining data, e.g., wirelessly, e.g., using a particular protocol, such as Bluetooth.

Fig. 8b shows an apparatus 120 which may implement at least some tools of the decoder apparatus 40 or 40a. The apparatus 120 may comprise a processor 121 and a non-transitory memory unit 122 storing instructions which, when executed by the processor 121, may cause the processor 121 to implement at least one of the TD and/or FD tools of the decoder apparatus 40 or 40a. In particular, the instructions may implement a subgroup of FD tools (e.g., TNS and/or noise filling) and other FD tools which are not in the subgroup (e.g., 44, 46, etc.). The instructions may also comprise instructions which, when executed by the

processor 121, perform a selection of the bandwidth so that the bandwidth of the signals processed by the tools in the subgroup of FD tools (e.g., TNS and/or noise filling) differs from the bandwidth of the signals processed by the other FD tools which are not in the subgroup (e.g., 44, 46, etc.). The instructions may be such as

5 to control a bandwidth selection based on energy detections associated to the different bandwidths, as, for example, performed by an encoder. The instructions may also comprise instructions which, when executed by the processor 121, permit to operate as a encoder and, in particular, permit to control the bandwidth of a subgroup of FD tools (e.g., 43, 45) which may be different from the bandwidth

10 of other FD tools. The bandwidth chosen for the subgroup at the encoder may be the same chosen for the subgroup at the decoder. The non-transitory memory unit 122 may also comprise other data, such as at least portions electronic versions of Tables 1, 2, and/or 3. The apparatus 120 may comprise a storage space 128 for storing, for example, a bitstream obtained from an information signal (e.g., an

15 audio signal). The apparatus 120 may comprise an output unit 127 for transmitting data, e.g., wirelessly, e.g., using a particular protocol, such as Bluetooth. The apparatus 120 may also comprise an input unit 126 for obtaining data, e.g., wirelessly, e.g., using a particular protocol, such as Bluetooth. For example, the apparatus 120 may obtain, by executing the instructions stored in the non-

20 transitory memory unit 122, a bitstream transmitted by a decoder.

In examples, the apparatus 110 and 120 may be the same device. In examples, the composition of different apparatus 110 and 120 form a system.

Depending on certain implementation requirements, examples may be implemented in hardware. The implementation may be performed using a digital

25 storage medium, for example a floppy disk, a Digital Versatile Disc (DVD), a Blu-Ray Disc, a Compact Disc (CD), a Read-only Memory (ROM), a Programmable Read-only Memory (PROM), an Erasable and Programmable Read-only Memory (EPROM), an Electrically Erasable Programmable Read-Only Memory (EEPROM) or a flash memory, having electronically readable control signals stored thereon,

30 which cooperate (or are capable of cooperating) with a programmable computer system such that the respective method is performed. Therefore, the digital storage medium may be computer readable.

Generally, examples may be implemented as a computer program product with program instructions, the program instructions being operative for performing one of the methods when the computer program product runs on a computer. The program instructions may for example be stored on a machine readable medium.

- 5 Other examples comprise the computer program for performing one of the methods described herein, stored on a machine readable carrier. In other words, an example of method is, therefore, a computer program having a program instructions for performing one of the methods described herein, when the computer program runs on a computer.
- 10 A further example of the methods is, therefore, a data carrier medium (or a digital storage medium, or a computer-readable medium) comprising, recorded thereon, the computer program for performing one of the methods described herein. The data carrier medium, the digital storage medium or the recorded medium are tangible and/or non-transitionary, rather than signals which are intangible and
- 15 transitory.

A further example comprises a processing unit, for example a computer, or a programmable logic device performing one of the methods described herein.

A further example comprises a computer having installed thereon the computer program for performing one of the methods described herein.

- 20 A further example comprises an apparatus or a system transferring (for example, electronically or optically) a computer program for performing one of the methods described herein to a receiver. The receiver may, for example, be a computer, a mobile device, a memory device or the like. The apparatus or system may, for example, comprise a file server for transferring the computer program to the
- 25 receiver.

In some examples, a programmable logic device (for example, a field programmable gate array) may be used to perform some or all of the functionalities of the methods described herein. In some examples, a field programmable gate array may cooperate with a microprocessor in order to

perform one of the methods described herein. Generally, the methods may be performed by any appropriate hardware apparatus.

The above described examples are illustrative for the principles discussed above. It is understood that modifications and variations of the arrangements and the
5 details described herein will be apparent. It is the intent, therefore, to be limited by the scope of the impending patent claims and not by the specific details presented by way of description and explanation of the examples herein.

Claims

1. An encoder apparatus comprising:

a plurality of frequency domain, FD, encoder tools for encoding an information signal, the information signal presenting a plurality of frames; and

5 an encoder bandwidth detector and controller configured to select a bandwidth for at least a subgroup of the plurality of FD encoder tools, the subgroup including less FD encoder tools than the plurality of FD encoder tools, on the basis of information signal characteristics so that at least one of the FD encoder tools of the subgroup has a different bandwidth with respect to at least
10 one of the FD encoder tools which are not in the subgroup.

2. The encoder apparatus of claim 1, wherein:

at least one FD encoder tool of the subgroup is a temporal noise shaping, TNS, tool and/or a noise level estimator tool.

3. The encoder apparatus of claim 1 or 2, wherein:

15 at least one FD encoder tool which is not in the subgroup is chosen among at least one of linear predictive coding, LPC, based spectral shaper, a spectral noise shaping, SNS, tool, a spectral quantizer, and a residual coder.

4. The encoder apparatus of any of the preceding claims, wherein:

20 the encoder bandwidth detector and controller is configured to select the bandwidth of the at least one FD encoder tool of the subgroup between at least a first bandwidth common to at least one of the FD encoder tools which are not in the subgroup and a second bandwidth different from the bandwidth of the at least one of the FD encoder tools which are not in the subgroup.

25 5. The encoder apparatus of any of the preceding claims, wherein:

the encoder bandwidth detector and controller is configured to select the bandwidth of the at least one of the plurality of FD encoder tools on the basis of at least one energy estimation on the information signal.

6. The encoder apparatus of any of the preceding claims, wherein:

5 the encoder bandwidth detector and controller is configured to compare at least one energy estimate associated to a bandwidth of the information signal to a respective threshold to control the bandwidth for the at least one of the plurality of FD encoder tools.

7. The encoder apparatus of any of the preceding claims, wherein:

10 the at least one of the plurality of FD encoder tools of the subgroup comprises a TNS configured to autocorrelate a TNS input signal within the bandwidth chosen by the encoder bandwidth detector and controller.

8. The encoder apparatus of any of the preceding claims, wherein the at least one of the FD encoder tools which are not in the subgroup is configured to operate
15 at a full bandwidth.

9. The encoder apparatus of any of the preceding claims, wherein the encoder bandwidth detector and controller is configured to select at least one bandwidth which is within the full bandwidth at which the at least one of the FD encoder tools which are not in the subgroup is configured to operate.

20 10. The encoder apparatus of any of the preceding claims, wherein the at least one of the remaining FD encoder tools of the plurality of FD encoder tools is configured to operate in open chain with respect to the bandwidth chosen by the encoder bandwidth detector and controller.

11. The encoder apparatus of any of the preceding claims, wherein:

25 the encoder bandwidth detector and controller is configured to select a bandwidth among a finite number of bandwidths and/or among a set of pre-defined bandwidths.

12. The encoder apparatus of any of the preceding claims, wherein:

the encoder bandwidth detector and controller is configured to perform a selection among at least one or a combination of the following bandwidths: a 4 KHz, 8 KHz, 12 KHz, 16 KHz, and 24 KHz, and/or NB, WB, SSWB, SWB, FB.

5 13. The encoder apparatus of any of the preceding claims, wherein:

the encoder bandwidth detector and controller is configured to control the signalling of the bandwidth to a decoder.

14. The encoder apparatus of any of the preceding claims, further configured to:

10 encode a control data field including an information regarding the chosen bandwidth.

15. The encoder apparatus of any of the preceding claims, further configured to:

encode a control data field including:

- 15
- 0 data bits corresponding to NB bandwidth;
 - 1 data bit corresponding to NB, WB bandwidth;
 - 2 data bits corresponding to NB, WB, SSWB bandwidth;
 - 2 data bits corresponding to NB, WB, SSWB, SWB bandwidth;
 - 3 data bits corresponding to NB, WB, SSWB, SWB, FB bandwidth.

20 16. The encoder apparatus of claim 5 or 6, wherein the at least one energy estimation is performed as:

$$E_B(n) = \sum_{k=I_{fs}(n)}^{I_{fs}(n+1)-1} \frac{X(k)^2}{I_{fs}(n+1) - I_{fs}(n)} \quad for \ n = 0 \dots N_B - 1$$

where $X(k)$ are the MDCT or MDST or FD coefficients, N_B is the number of bands and $I_{f_s}(n)$ are the indices associated to the band.

17. The encoder apparatus of any of the preceding claims, further comprising a TNS tool configured to perform a filtering operation including the calculation of the filtering operation:

for each $k = 0..8$

$$r(k) = \begin{cases} r_0(k) & , \text{ if } \prod_{s=0}^2 e(s) = 0 \\ \frac{\sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1-k} X_s(n)X_s(n+k)}{e(s)} & , \text{ otherwise} \end{cases}$$

with

$$r_0(k) = \begin{cases} 1 & , \text{ if } k = 0 \\ 0 & , \text{ otherwise} \end{cases}$$

and

$$e(s) = \sum_{n=\text{sub_start}(f,s)}^{\text{sub_stop}(f,s)-1} X_s(n)^2 \quad \text{for } s = 0..2$$

- 10 where $X(k)$ are the MDCT or MDST or FD coefficients, $\text{sub_start}(f, s)$ and $\text{sub_stop}(f, s)$ are associated to the particular bandwidth as detected by the encoder bandwidth detector and controller.

18. The encoder apparatus of any of the preceding claims, further comprising a noise estimator configured to estimate a noise level using

$$L_{NF} = \frac{\sum_{k=0}^{N_E-1} I_{NF}(k) \cdot \frac{|X_f(k)|}{gg}}{\sum_{k=0}^{N_E-1} I_{NF}(k)}$$

- 15 where gg refers to the global gain, $I_{NF}(k)$ to the identification of the spectral lines on which the noise level is to be estimated, and $X_f(k)$ is the signal.

19. A decoder apparatus comprising a plurality of FD decoder tools for decoding an information signal encoded in a bitstream, wherein:

the FD decoder tools are divided:

- in a subgroup comprising at least one FD decoder tool;
- in remaining FD decoder tools comprising at least one FD decoder tool;

5 wherein the decoder apparatus is configured so that at least one of the plurality of decoder tools of the subgroup performs signal processing a different bandwidth with respect to at least one of the remaining FD decoder tools of the plurality of decoder tools.

20. The decoder apparatus of claim 19, further comprising a bandwidth controller configured to:

10 choose the bandwidth on the basis of the bandwidth information.

21. The decoder apparatus of claim 19 or 20, wherein:

the subgroup comprises at least one of a decoder noise filling tool and/or a temporal noise shape, TNS, decoder.

22. The decoder apparatus of any of claims claim 19-21, wherein:

15 the at least one of the remaining FD decoder tools is at least one of a decoder linear predictive coding, LPC, tool, a spectral noise shaping, SNS, tool, a decoder global gain tool, an MDCT or MDST shaping tool or another shaping tool based on another transformation.

23. The decoder apparatus of any of claims claim 19-22, further configured to:

20 control the bandwidth of the at least one of the plurality of decoder tools in the subgroup between:

at least a first bandwidth common to at least one of the remaining FD decoder tools; and

at least a second bandwidth different from the first bandwidth.

24. The decoder apparatus of any of claims 19-23, wherein the at least one of the FD remaining decoder tools is configured to operate at a full bandwidth.

25. The decoder apparatus of any of claims 19-24, wherein the at least one of the remaining FD decoder tools is configured to operate in open chain with respect
5 to the chosen bandwidth.

26. The decoder apparatus of any of claims 19-25, further configured to:

choose a bandwidth among a finite number of bandwidths and/or among a set of pre-defined bandwidths.

27. The decoder apparatus of any of claims 19-26, further configured to:

10 perform a choice among at least one or a combination of: a 8 KHz, 16 KHz, 24 KHz, 32 KHz, and 48 KHz and/or NB, WB, SSWB, SWB, FB.

28. The decoder apparatus of any of claims 19-27, further comprising a noise filling tool configured to apply a noise level using indices given by

$$I_{NF}(k) = \begin{cases} 1 & \text{if } 24 \leq k < bw_{stop} \text{ and } \widehat{X}_q(i) == 0 \text{ for all } i = k - 3.. \min(bw_{stop}, k + 3) \\ 0 & \text{otherwise} \end{cases}$$

15 where bw_{stop} is obtained on the basis of bandwidth information in the bitstream.

29. The decoder apparatus of any of claims 19-28, further comprising a TNS decoder configured to perform:

$$s^0(\text{start_freq}(0) - 1) = s^1(\text{start_freq}(0) - 1) = \dots = s^7(\text{start_freq}(0) - 1) = 0$$

for $f = 0$ to $\text{num_tns_filters}-1$ do

for $n = \text{start_freq}(f)$ to $\text{stop_freq}(f) - 1$ do

$$t^K(n) = \widehat{X}_f(n)$$

20 for $k = 7$ to 0 do

$$t^k(n) = t^{k+1}(n) - rc_q(k)s^k(n-1)$$

$$s^{k+1}(n) = rc_q(k)t^k(n) + s^k(n-1)$$

$$\widehat{X}_s(n) = s^0(n) = t^0(n)$$

where $\widehat{X}_s(n)$ is the output of the TNS decoder and $\widehat{X}_f(n)$ is the input of the TNS decoder, num_tns_filters, start_freq, stop_freq are obtained on the basis of bandwidth information (39a) in the bitstream.

30. A system comprising an encoder apparatus and a decoder apparatus,
 5 wherein the encoder apparatus is according to any of claims 1-18 and/or the decoder apparatus is according to any of claims 19-29.

31. A method for encoding an information signal according to at least a plurality of operations in the frequency domain, FD, the method comprising:

selecting a bandwidth for a subgroup of FD operations;

10 performing first signal processing operations at the a bandwidth for the subgroup of FD operations;

performing second signal processing operations at a different bandwidth for FD operations which are not in the subgroup.

32. A method for decoding a bitstream with an information signal and control
 15 data, the method comprising a plurality of signal processing operations in the frequency domain, FD, the method comprising:

choosing a bandwidth selection for a subgroup of FD operations on the basis of the control data;

20 performing first signal processing operations at the a bandwidth for the subgroup of FD operations;

performing second signal processing operations at a different bandwidth for FD operations which are not in the subgroup.