

Description

This invention relates to an electronic circuit comprising a main stage for producing a digital code formed with a plurality of binary bits that make transitions between a pair of binary values as a function of an input parameter, the transitions ideally occurring in accordance with a prescribed algorithm as the input parameter traverses an input range, the bits divided into a first set of at least one bit and a second set of at least one bit.

In an analog-to-digital converter (ADC) of the folding type, the bits that form the digital output code are typically generated along two largely separate paths. A coarse conversion circuit produces a set of coarse (higher-order) bits directly in response to an analog input voltage. A fine conversion circuit produces a set of fine (lower-order) bits in response to one or more signals supplied from a circuit that folds the input voltage. See van de Grift et al, "A Monolithic 8-Bit Video A/D Converter", IEEE JSSC, June 1984, pp. 374-378.

One difficulty with a folding ADC is that the bits in the coarse set sometime make transitions between binary "0" and binary "1" at values of the input voltage that are slightly different from those at which the bits in the fine set make corresponding transitions between "0" and "1". For example, consider a simple 4-bit "binary" code consisting of two bits supplied from the coarse converter and two bits supplied from the fine converter. Assume that the input voltage is in that portion of the input voltage range where the code is (0011). The left-most bit is the most significant bit (MSB). The right-most bit is the least significant bit (LSB).

If the input voltage is increased by an amount equivalent to one LSB, the code should change to (0100). This is not always the case. Because the coarse bits are generated along a separate path from the fine bits, phenomena such as input offset errors, delay differences, and/or noise may cause the second most significant bit to switch from "0" to "1" either before or after the two fine bits switch from "1" to "0". The code changes to either (0111) or to (0000). The actual value produced by the code is nowhere near the intended (0100). A similar problem occurs with the MSB.

The present invention is an electronic circuit that utilizes a bit synchronization scheme to overcome bit transition problems of the foregoing type. The starting point for the invention is a main stage that produces a digital code formed with a plurality of bits that make transitions between a pair of binary values as a function of an input parameter. The transitions ideally occur in accordance with a prescribed algorithm as the input parameter traverses a suitable input range. The bits are divided into first and second sets, each consisting of one or more bits. The bits in the second set

are usually generated along a different electronic path through the main stage than those in the first.

A synchronization stage synchronizes transitions of bits in the first set with corresponding transitions of bits in the second set. The basic approach is to define transition regions for bits in the first set. A transition region is basically that portion of the input range where phenomena of the type mentioned above can cause a bit in the first set to go to a wrong value. Using signals keyed to the transition regions, the values of bits in the first set are suitably replaced with information based on bits in the second. A single bit in the second set is preferably used in making every replacement in a transition region. No replacements are made outside the transition regions.

In particular, the synchronization stage implements this scheme with a control circuit and a substitution circuit. The control circuit generates at least one synchronization signal. Each synchronization signal reaches a specified signal level when the input parameter is in a sub-range of the input range where a bit in the first set associated with that synchronization signal makes a single transition and where a corresponding bit in the second set makes a single transition. When each synchronization signal reaches its specified signal level, the substitution circuit replaces the value of each associated bit in the first set either with the value of the corresponding bit in the second set or with the value of the complement of the corresponding bit in the second set in such a way as to satisfy the prescribed algorithm.

UK Patent Application GB-A-2 090 707 discloses an analogue-to-digital converter having a flash converter for providing a first set of coarse bits and a dual folding converter for providing a second set of fine bits. Transitions in the coarse bits are replaced with transitions from the fine bits in substitution means. The substitution means, however, are not responsive to synchronization signals indicating that the input parameter is in a sub-range of the input range where a bit in the first set of coarse bits associated with that synchronization signal makes a single transition and where a corresponding bit in the second set of fine bits also makes a single transition.

The algorithm is typically a binary code in which each bit in the first set is a coarse bit and in which each bit in the second set is a fine bit. In this case, only the most significant fine bit is used in performing the replacements. The substitution circuit specifically replaces the value of each coarse bit either with the value of the complement of the most significant fine bit if that coarse bit is the most significant bit making a transition within the appropriate sub-range or with the value of the most significant fine bit if that coarse bit is not the most significant bit making a transition within the sub-range.

The invention eliminates synchronization errors in an easy and dependable way. In folding ADCs, the

synchronization can be achieved with a small amount of circuitry.

Embodiments of the invention will now be described by way of example with reference to the accompanying drawings in which:

Fig. 1 is a block diagram of an ADC employing a bit synchronization scheme in accordance with the invention;

Figs. 2 and 3 are graphs indicating how the synchronization signals and bits vary as a function of input voltage for the ADC of Fig. 1;

Fig. 4 is a block diagram of a folding-type embodiment of the ADC of Fig. 1;

Fig. 5 is a circuit diagram of the input amplifier and summing arrays in the ADC of Fig. 4, Fig. 6 is a graph showing how the output voltage from a typical input amplifier in the amplifier array of Fig. 5 varies with input voltage and Fig. 7 is a circuit diagram of this amplifier;

Fig. 8 is a graph showing how a typical folding voltage supplied from the interpolation circuit in the ADC of Fig. 4 varies with input voltage;

Fig. 9 is a logic diagram for the bit substitution circuit in the ADC of Fig. 4.

Like reference symbols are employed in the drawings and in the description of the preferred embodiments to represent the same or very similar item or items. "N" is used as a subscript to indicate signals complementary to previously defined signals.

Referring to the drawings, Fig. 1 illustrates an analog-to-digital converter containing a main stage that converts an analog input voltage V_i into an M-bit digital code divided into a set of K first binary bits $B_0, B_1, \dots, B_{K-1}$ and a set of M-K second binary bits $B_K, B_{K+1}, \dots, B_{M-1}$. K is typically two or more but may be as small as 1. The same applies to M-K.

Bits $B_0 - B_{M-1}$ make transitions between "0" and "1" as input voltage V_i traverses an input range extending from a low voltage level V_0 to a high voltage level V_F . The code is fixed at one value, usually "all 0s", when input V_i is at or below V_0 . The code is fixed at another value, usually "all 1s", when input V_i is at or above V_F . The binary transitions ideally occur in accordance with a prescribed algorithm.

The main ADC stage consists of a first conversion circuit 10 and a second conversion circuit 12. Conversion circuit 10 generates bits $B_0 - B_{K-1}$ along one general electronic path. Conversion circuit 12 similarly generates bits $B_K - B_{M-1}$ along another general electronic path.

Slight, but significant, differences in signal transmission delay typically exist between the two paths. If not suitably accounted for, these delay differences can cause transitions in certain of bits $B_0 - B_{M-1}$ to occur at wrong points. The same applies to noise. Small offset errors in the input circuitry of circuits 10 and 12 affect the values of input V_i at which bits $B_0 - B_{M-1}$ make their transitions. A bit synchronization stage op-

erating according to the teachings of the invention synchronizes transitions of bits $B_0 - B_{K-1}$ in the first-mentioned set with corresponding transitions of bits $B_K - B_{M-1}$ in the second-mentioned set so as to avoid code errors due to the preceding phenomena.

The synchronization stage is formed with a synchronization control circuit 14 and a bit substitution circuit 16. Control circuit 14 generates one or more synchronization voltage signals V_A that indicate when bits $B_0 - B_{K-1}$ in the first set are near values of input V_i at which these bits make transitions that are ideally intended to be simultaneous with transitions of bits $B_K - B_{M-1}$ in the second set. When one of the V_A signals provides such an indication for one of the bits in the first set, substitution circuit 16 replaces the current value of that bit with information derived from the bits in the second set. Circuit 16 thus provides a set of K bits $B'_0, B'_1, \dots, B'_{K-1}$ synchronized with the M-K bits in the second set. Except where a replacement is performed, each bit in the synchronized set is the same as the like-numbered bit in the original set.

The specific operation of the synchronization stage is facilitated with the assistance of the graph shown in Fig. 2. Depending on the algorithm followed by the code, one or two situations exists for each value of input V_i at which a bit in the first set makes a transition: (1) at least one bit in the second set ideally makes a simultaneous transition, or (2) no bit in the second set ideally makes a simultaneous transition. Starting with the last-mentioned situation, this case is represented by transition T of bit B_1 in Fig. 2. None of bits $B_K - B_{M-1}$ ideally make transitions simultaneous with transition T. Consequently, no bit synchronization problem exists here.

Turning to the first-mentioned situation, a voltage transition region is defined around each value of input V_i at which a bit in the first set makes a transition ideally intended to be simultaneous with the transition of a corresponding bit in the second set. The transition regions are sufficiently narrow that each bit within a transition region makes only a single transition between "0" and "1". However, the regions are sufficiently wide that the likelihood of a transition occurring outside the intended region is usually very small.

Fig. 2 illustrates three transition regions X, Y and Z generally centered around respective transitions in bits B_0, B_1 , and B_{K-1} . Bit B_{K-1} is also shown as making a transition within region Y. The transitions for the corresponding bits in the second set can be seen by inspection. Note that there are two possible corresponding bits in the second set for the B_{K-1} transition in region Z.

Control circuit 14 supplies one V_A synchronization signal for each transition region. Each synchronization signal is thereby associated with one or more particular bits in the first set. Each V_A signal reaches a specified condition when input V_i is in the transition region for the associated bit or bits in the first set.

Fig. 2 depicts three such synchronization signals V_{AX} , V_{AY} , and V_{AZ} respectively associated with transition regions X, Y, and Z. Each of signals V_{AX} - V_{AZ} is normally at a low voltage V_L and goes to a high voltage V_H when input V_i enters the associated transition region. The boundaries of regions X - Z are set at points where the V_A signals are approximately equal to $(V_L + V_H)/2$. That is, the specified V_A condition mentioned in the previous paragraph occurs when each V_A signal rises above a threshold voltage approximately equal to $(V_L + V_H)/2$.

When a V_A signal is at the specified condition, substitution circuit 16 replaces the value of each associated bit in the first set either (a) with the value of the corresponding bit in the second set or (b) with the value of the binary complement of the corresponding bit in the second set in such a way as to meet the prescribed algorithm. For example, circuit 16 uses the true value of the corresponding bit in a second set when the transitions of the two bits are in the same direction. The reverse applies when the transitions are in opposite directions. The dotted lines in Fig. 2 indicate the replaced bit values. In this way, the invention synchronizes transitions in bits B_0 - B_{K-1} with transitions in bits B_K - B_{M-1} .

Synchronization errors of lesser significance may exist among transitions of bits in the second set. The effect of these errors can be minimized by utilizing a single bit in the second set for performing all the replacements within a transition region. That is, circuit 16 performs each replacement in the region using either the true or complementary value of this bit. Depending on the algorithm prescribed for bits B_0 - B_{M-1} , one bit in the second set may be used in performing all the replacements in all the transition regions.

The prescribed algorithm is preferably a binary code in which bits B_0 - B_{M-1} decrease progressively in significance as the bit number increases. Bits B_0 - B_{K-1} are thereby the coarse bits. Bits B_K - B_{M-1} are the fine bits. Fig. 3 illustrates how a binary code varies as a function of input V_i for the case where both K and M-K are equal to 2.

To obtain the greatest synchronization accuracy in the binary-code situation, substitution circuit 16 only uses the most significant fine bit B_K for replacement purposes. This is bit B_2 in Fig. 3. More specifically, circuit 16 employs the following replacement procedure in each transition region. The value of the complement of the most significant fine bit is substituted for the value of each coarse bit if it is the most significant bit making a transition within that region. The true value of the most significant fine bit is substituted for the value of each other coarse bit that makes a transition within that region. The replacement algorithm can be expressed as:

$$B_j' = \begin{cases} B_j & j > p \\ B_{NK} & j = p \\ B_K & j < p \end{cases}$$

where the subscript j is an integer from 0 to K-1, p is an integer equal in value to the subscript of the most significant coarse bit B_p making a transition within the transition region, and B_{NK} is the complement of B_K .

The transition regions have largely the same width. As indicated in Fig. 3, the width is preferably equal to one half of the average spacing between consecutive transitions of the least significant coarse bit B_{K-1} . In addition, the number of transition regions is equal to the number of transitions made by bit B_{K-1} .

The invention is particularly useful in ADCs of the folding type. Turning to Fig. 4, it depicts an application of the invention to a multiple-folding ADC that generates an 8-bit binary code consisting of two coarse bits supplied along one analog path and six fine bits supplied along another analog path. Coarse conversion circuit 10 of Fig. 1 is implemented with an input amplifier array 18, a summing array 20, and a group of coarse comparators 22. Fine conversion circuit 12 of Fig. 1 is formed with arrays 18 and 20, an interpolation circuit 24, a group of fine comparators 26, and an encoder 28.

The V_A synchronization signals consist of part of the amplifier signals normally supplied from amplifier array 18. No additional circuitry beyond that already present in array 18 is needed to implement control circuit 14 of Fig. 1. This provides a particularly efficient arrangement.

Fig. 5 shows details of arrays 18 and 20. Amplifier array 18 contains 64 input amplifiers A_0 - A_{63} arranged in 8 rows by 8 columns. Letting i be a running integer, each amplifier A_i amplifies the difference between input voltage V_i and a corresponding reference voltage V_{Ri} to produce an intermediate voltage V_{Ai} . Reference voltages V_{R0} , V_{R1} ... V_{R63} are supplied from a resistive divider at equally spaced intervals across the input range from V_{R0} to V_{R63} .

Fig. 6 illustrates the general shape for typical voltage V_{Ai} as a function of V_i . V_{Ai} would ideally have the triangular shape shown in dashed line. Due to practical amplifier characteristics, V_{Ai} actually has the more rounded shape indicated in solid line. Each voltage V_{Ai} exceeds threshold voltage $(V_L + V_H)/2$ when V_i is between V_{Ri-8} and V_{Ri} .

The internal configuration for typical amplifier A_i is depicted in Fig. 7. The collectors of emitter-coupled NPN transistors Q_{Li} and Q_{Ri} are respectively connected to the collectors of transistors Q_{Ri-8} and Q_{Li+8} in amplifiers A_{i-8} and A_{i+8} . Amplifier A_i operates in a differential manner in conjunction with amplifier A_{i-8} . When V_i equals V_{Ri} , the differential pair Q_{Li} and Q_{Ri}

is balanced so that V_{A1} has a zero crossing. The differential pair Q_{L1-8} and Q_{R1-8} in amplifier A_{1-8} is similarly balanced when V_i equals V_{RL-8} . Due to the collector cross coupling, V_{A1} has another zero crossing at that point. The result is that V_{A1} reaches V_H when V_i equals V_{RL-4} and is constant at V_L when V_i is less than V_{RL-12} or greater than R_{RH4} .

Summing array 20 is subdivided into a folding array 30 and a coarse summing section 32 as indicated back in Fig. 5. Folding array 30 electrically combines every sixteenth intermediate signal V_{A1} to produce eight pairs of complementary folded voltages V_B . Summing section 32 electrically combines certain of the V_{A1} signals to generate two pairs of complementary voltages V_C . The circles in Fig. 5 represent the particular couplings to the A_i amplifiers.

Returning to Fig. 4, coarse comparator group 22 consists of two master-slave flip-flops that generate coarse bits B_0 and B_1 by comparing the two true V_C voltages with their respective complements.

Interpolation circuit 24 interpolates by a factor of four between consecutive V_B voltages to produce 32 pairs of complementary folded voltages V_D . Circuit 24 is preferably embodied in the manner disclosed in EP-A-0 319 097. Fig. 8 depicts a typical one of voltages V_D as a function of V_i .

Fine comparator group 26 consists of a flip-flop ladder that compares the 32 true V_D voltages with their respective complements to generate a string of 32 bits B in a circular code. The bit signals supplied from the ladder include the most significant fine bit B_2 and its complement B_{N2} as indicated in Fig. 4. The ladder is preferably implemented in the manner described in EP-A-0 328 215, entitled "Error Correction Circuit Suitable for Thermometer or Circular Code".

Encoder 28 converts bit string B into fine bits $B_3 - B_7$. Encoder 28 is formed with suitable input logic and an appropriately programmed read-only memory. The additional digital processing used to generate bits $B_3 - B_7$ beyond that experienced by bit B_2 usually does not in itself introduce any significant synchronization problems among the fine bits.

Fig. 9 shows a functional equivalent of a preferred embodiment for the internal circuitry of substitution circuit 16. This functional equivalent consists of a complementary-output buffer 34, an OR/NOR gate 36, AND gates 38, 40, 42, 44, and 46, and OR gates 48 and 50 interconnected in the indicated way. Gates 34 - 50 operate according to conventional combinatorial principles. Gates 38 - 50 are preferably implemented in current-mode logic (CML) using one CML gate for gates 38, 40, and 48 and another CML gate for gates 42 - 46 and 50.

Synchronization signals V_A supplied to substitution circuit 16 in Fig. 9 consist of amplifier voltages V_{A19} , V_{A35} , and V_{A51} provided from input amplifier array 18 in Fig. 4. Although the ADC in Fig. 4 generates an 8-bit binary code, only coarse bits B_0 and B_1 and

fine bit B_2 are specifically involved in the synchronization scheme. The same applies to the 4-bit binary code represented in Fig. 3. Consequently, the three synchronization signals that respectively identify when input V_i is in the low, middle, and high transition regions shown in Fig. 3 are roughly equivalent to signals V_{A19} , V_{A35} , and V_{A51} .

The specific operations performed by circuit 16 in Fig. 4 consist of (a) replacing the value of bits B_0 and B_1 respectively with the values of bits B_{N2} and B_2 if V_{A35} is greater than $(V_L + V_H)/2$ and (b) replacing the value of bit B_1 with the value of bit B_{N2} if V_{A19} or V_{A51} is greater than $(V_L + V_H)/2$. A noteworthy feature of this substitution is that the value of bit B_1 is replaced with the value of bit B_{N2} when V_i is in either the low or high transition region shown in Fig. 3. Consequently, a logical OR combinatorial operation can be performed on voltages V_{A19} and V_{A35} as indicated by OR/NOR gate 36 in Fig. 9. More generally, circuit 16 can use a gate that generates the logical OR of each group of synchronization signals for which the same replacements are performed.

While the invention has been described with reference to particular embodiments, this description is solely for the purpose of illustration and is not to be construed as limiting the scope of the invention claimed below. For example, the definitions of binary "0" and "1" are arbitrary and could be reversed from what is described above. Where the same replacements are made in a group of transition regions, a single synchronization signal could be initially defined instead of a group of synchronization signals combined by a logical OR operation.

The algorithm for the bits could consist of two or more sub-algorithms such as a binary code for the bits in the first set and a Gray code for the bits in the second set. Multiple applications of the invention can be made if there are three or more generally different sets of bits. Various modifications and applications may thus be made by those skilled in the art without departing from the true scope of the invention as defined in the appended claims.

Claims

1. An electronic circuit comprising a main stage (10, 12) for producing a digital code formed with a plurality of binary bits ($B_0..B_{M-1}$) that make transitions between a pair of binary values as a function of an input parameter (V_i), the transitions ideally occurring in accordance with a prescribed algorithm as the input parameter (V_i) traverses an input range ($V_0..V_F$), the bits being divided into a first set ($B_0..B_{K-1}$) of at least one bit and a second set ($B_K..B_{M-1}$) of at least one bit, characterized by a synchronization stage (14, 16) comprising: control means (14) responsive to the input



is balanced so that V_{A1} has a zero crossing. The differential pair Q_{L1-8} and Q_{R1-8} in amplifier A_{1-8} is similarly balanced when V_i equals V_{R1-8} . Due to the collector cross coupling, V_{A1} has another zero crossing at that point. The result is that V_{A1} reaches V_H when V_i equals V_{R1-4} and is constant at V_L when V_i is less than V_{R1-12} or greater than V_{R11-4} .

Summing array 20 is subdivided into a folding array 30 and a coarse summing section 32 as indicated back in Fig. 5. Folding array 30 electrically combines every sixteenth intermediate signal V_{A1} to produce eight pairs of complementary folded voltages V_B . Summing section 32 electrically combines certain of the V_{A1} signals to generate two pairs of complementary voltages V_C . The circles in Fig. 5 represent the particular couplings to the A_i amplifiers.

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Fig. 9 shows a functional equivalent of a preferred embodiment for the internal circuitry of substitution circuit 16. This functional equivalent consists of a complementary-output buffer 34, an OR/NOR gate 36, AND gates 38, 40, 42, 44, and 46, and OR gates 48 and 50 interconnected in the indicated way. Gates 34 - 50 operate according to conventional combinatorial principles. Gates 38 - 50 are preferably implemented in current-mode logic (CMJ) using one CML gate for gates 38, 40, and 48 and another CML gate for gates 42 - 46 and 50.

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fine bit B_2 are specifically involved in the synchronization scheme. The same applies to the 4-bit binary code represented in Fig. 3. Consequently, the three synchronization signals that respectively identify when input V_i is in the low, middle, and high transition regions shown in Fig. 3 are roughly equivalent to signals V_{A19} , V_{A35} , and V_{A51} .

The specific operations performed by circuit 16 in Fig. 4 consist of (a) replacing the value of bits B_0 and B_1 respectively with the values of bits B_{N2} and B_2 if V_{A35} is greater than $(V_L + V_H)/2$ and (b) replacing the value of bit B_1 with the value of bit B_{N2} if V_{A19} or V_{A51} is greater than $(V_L + V_H)/2$. A noteworthy feature of this substitution is that the value of bit B_1 is replaced with the value of bit B_{N2} when V_i is in either the low or high transition region shown in Fig. 3. Consequently, a logical OR combinatorial operation can be performed on voltages V_{A19} and V_{A35} as indicated by OR/NOR gate 36 in Fig. 9. More generally, circuit 16 can use a gate that generates the logical OR or each group of synchronization signals for which the same replacements are performed.

While the invention has been described with reference to particular embodiments, this description is solely for the purpose of illustration and is not to be construed as limiting the scope of the invention claimed below. For example, the definitions of binary "0" and "1" are arbitrary and could be reversed from what is described above. Where the same replacements are made in a group of transition regions, a single synchronization signal could be initially defined instead of a group of synchronization signals combined by a logical OR operation.

The algorithm for the bits could consist of two or more sub-algorithms such as a binary code for the bits in the first set and a Gray code for the bits in the second set. Multiple applications of the invention can be made if there are three or more generally different sets of bits. Various modifications and applications may thus be made by those skilled in the art without departing from the true scope of the invention as defined in the appended claims.

Claims

1. An electronic circuit comprising a main stage (10, 12) for producing a digital code formed with a plurality of binary bits ($B_0..B_{M-1}$) that make transitions between a pair of binary values as a function of an input parameter (V_i), the transitions ideally occurring in accordance with a prescribed algorithm as the input parameter (V_i) traverses an input range ($V_0..V_F$), the bits being divided into a first set ($B_0..B_{K-1}$) of at least one bit and a second set ($B_K..B_{M-1}$) of at least one bit, characterized by a synchronization stage (14, 16) comprising: control means (14) responsive to the input



parameter (V_i) for generating at least one synchronization signal (V_A), each synchronization signal having a specified signal level indicating that the input parameter (V_i) is in a sub-range of the input range (V_0-V_F) where a bit in the first set ($B_0..B_{K-1}$) associated with that synchronization signal makes a single transition and where a corresponding bit in the second set ($B_K..B_{M-1}$) also makes a single transition; and

substitution means (16) responsive to each synchronization signal (V_A) when it has the specified signal level for replacing the value of each associated bit in the first set ($B_0..B_{K-1}$) either with the value of the corresponding bit in the second set ($B_K..B_{M-1}$) or with the value of the binary complement of the corresponding bit in the second set ($B_K..B_{M-1}$) in such a manner as to satisfy the prescribed algorithm.

2. A circuit as claimed in claim 1, characterized in that the substitution means (16) performs every replacement in a subrange using either the value of a single one of the bits in the second set ($B_K..B_{M-1}$) or the value of the binary complement of that single bit.
3. A circuit as claimed in claim 1, characterized in that there are first (10) and second (12) bit generation paths through the main stage (10, 12), each bit in the first set ($B_0..B_{K-1}$) being generated along the first path (10), each bit in the second set ($B_K..B_{M-1}$) being generated along the second path (12) and in that the specified signal level for each synchronization signal (V_A) consists of it being greater than or being less than a threshold voltage.
4. A circuit as claimed in claim 3, characterized in that there is only one subrange for each synchronization signal (V_A).
5. A circuit as claimed in claim 3 wherein there are at least two bits in each set, characterized in that the substitution means (16) performs every replacement in a subrange using either the value of a single one of the bits in the second set ($B_K..B_{M-1}$) or the value of the binary complement of that single bit.
6. A circuit as claimed in claim 1, characterized in that:

the main stage comprises means (18) responsive to the input parameter (V_i) for producing a plurality of intermediate signals ($V_{A0}-V_{A83}$) that vary between first (V_L) and second (V_H) levels as the input parameter (V_i) traverses the input range ($V_{R0}-V_{R83}$), each intermediate signal ($V_{A0}-V_{A83}$) being at the first level (V_L) for most of the traverse of the input parameter (V_i) across the input range

($V_{R0}-V_{R83}$) and reaching the second level (V_H) at a different value ($V_{Ri,4}$) of the input parameter (V_i) than each other intermediate signal, and means (20, 24, 26, 28) responsive to the intermediate signals ($V_{A0}-V_{A83}$) for generating each bit in the second set ($B_2..B_7$); and

the synchronisation stage comprises part (18, 16) of the main stage, each synchronization signal ($V_{A19}, V_{A35}, V_{A51}$) being one of the intermediate signals ($V_{A0}-V_{A83}$).

Patentansprüche

1. Elektronische Schaltung mit einer Hauptstufe (10, 12) zum Erzeugen eines digitalen Codes, der mit einer Vielzahl binärer Bits ($B_0..B_{M-1}$) gebildet wird, die als Funktion eines Eingangsparameters (V_i) Übergänge zwischen einem Paar binärer Werte machen, wobei die Übergänge idealerweise entsprechend einem vorgeschriebenen Algorithmus auftreten, wenn der Eingangsparameter (V_i) einen Eingangsbereich (V_0-V_F) durchläuft, wobei die Bits in einen ersten Satz ($B_0..B_{K-1}$) aus mindestens einem Bit und einen zweiten Satz ($B_K..B_M$) aus mindestens einem Bit unterteilt werden, gekennzeichnet durch eine Synchronisationsstufe (14, 16) mit:

auf den Eingangsparameter (V_i) ansprechenden Steuerungsmitteln (14) zum Erzeugen mindestens eines Synchronisationssignals (V_A), wobei jedes Synchronisationssignal einen festgelegten Signalpegel aufweist, der angibt, daß der Eingangsparameter (V_i) in einem Teilbereich des Eingangsbereichs (V_0-V_F) liegt, wo ein zu diesem Synchronisationssignal gehörendes Bit im ersten Satz ($B_0..B_{K-1}$) einen einfachen Übergang macht und ein entsprechendes Bit in dem zweiten Satz ($B_K..B_{M-1}$) auch einen einfachen Übergang macht und

einem auf jedes Synchronisationssignal (V_A) ansprechenden Substitutionsmittel (16), wenn dieses den festgelegten Signalpegel hat, um den Wert jedes zugehörigen Bits in dem ersten Satz ($B_0..B_{K-1}$) entweder durch den Wert des entsprechenden Bits in dem zweiten Satz ($B_K..B_{M-1}$) oder durch den Wert des Binärkomplements des entsprechenden Bits in dem zweiten Satz ($B_K..B_{M-1}$) auf solche Weise zu ersetzen, daß der vorgeschriebene Algorithmus erfüllt wird.

2. Schaltung nach Anspruch 1, dadurch gekennzeichnet, daß das Substitutionsmittel (16) jedes Ersetzen in einem Teilbereich ausführt, wobei es entweder den Wert eines einzelnen der Bits in dem zweiten Satz ($B_K..B_{M-1}$) oder den Wert des Binärkomplements dieses einzelnen Bits verwendet.



3. Schaltung nach Anspruch 1, dadurch gekennzeichnet, daß es erste (10) und zweite (12) Bit erzeugungspfade durch die Hauptstufe (10, 12) gibt, wobei jedes Bit in dem ersten Satz ($B_0..B_{K-1}$) entlang des ersten Pfades (10) erzeugt wird, jedes Bit in dem zweiten Satz ($B_K..B_{M-1}$) entlang des zweiten Pfades (12) erzeugt wird, und daß der festgelegte Signalpegel für jedes Synchronisationssignal (V_A) darin besteht, daß er größer oder kleiner als eine Schwellenspannung ist.

4. Schaltung nach Anspruch 3, dadurch gekennzeichnet, daß es für jedes Synchronisationssignal (V_A) nur einen Teilbereich gibt.

5. Schaltung nach Anspruch 3, wobei in jedem Satz mindestens zwei Bits enthalten sind, dadurch gekennzeichnet, daß das Substitutionsmittel (16) jedes Ersetzen in einem Teilbereich ausführt, wobei es entweder den Wert eines einzelnen der Bits in dem zweiten Satz ($B_K..B_{M-1}$) oder den Wert des Binärkomplements dieses einzelnen Bits verwendet.

6. Schaltung nach Anspruch 1, dadurch gekennzeichnet, daß:

die Hauptstufe auf den Eingangsparameter (V_I) ansprechende Mittel (18) umfaßt, zum Erzeugen einer Vielzahl von Zwischensignalen ($V_{A0}-V_{A63}$), die zwischen ersten (V_L) und zweiten (V_H) Pegeln variieren, wenn der Eingangsparameter (V_I) den Eingangsbereich ($V_{R0}-V_{R63}$) durchläuft, wobei jedes Zwischensignal ($V_{A0}-V_{A63}$) für den größten Teil des Durchlaufens des Eingangsparameters (V_I) durch den Eingangsbereich ($V_{R0}-V_{R63}$) auf dem ersten Pegel (V_L) liegt und den zweiten Pegel (V_H) bei einem anderen Wert (V_{R1-4}) des Eingangsparameters (V_I) erreicht als jedes andere Zwischensignal, sowie Mittel (20, 24, 26, 28), die auf die Zwischensignale ($V_{A0}-V_{A63}$) ansprechen, um jedes Bit in dem zweiten Satz ($B_2..B_7$) zu erzeugen; und

die Synchronisationsstufe einen Teil (18, 16) der Hauptstufe umfaßt, wobei jedes Synchronisationssignal (V_{A19} , V_{A35} , V_{A51}) eines der Zwischensignale ($V_{A0}-V_{A63}$) ist.

Revendications

1. Circuit électronique comprenant un étage principal (10, 12) pour produire un code numérique formé avec une pluralité de bits binaires ($B_0..B_{M-1}$) qui fait des transitions entre une paire de valeurs binaires en fonction d'un paramètre d'entrée (V_I), les transitions se produisant dans le cas idéal selon un algorithme prescrit lorsque le paramètre d'entrée (V_I) traverse une gamme d'entrée (V_0-

V_F), les bits étant divisés en un premier jeu ($B_0..B_{K-1}$) d'au moins un bit et un deuxième jeu ($B_K..B_M$) d'au moins un bit, caractérisé par un étage de synchronisation (14, 16) comprenant:

des moyens de commande (14) sensibles au paramètre d'entrée (V_I) pour engendrer au moins un signal de synchronisation (V_A), chaque signal de synchronisation présentant un signal de niveau spécifique indiquant que le paramètre d'entrée (V_I) se trouve dans une gamme partielle de la gamme d'entrée (V_0-V_F) alors qu'un bit dans le premier jeu ($B_0..B_{K-1}$) associé audit signal de synchronisation fait une transition simple et un bit correspondant dans le deuxième jeu ($B_K..B_{M-1}$) fait également une transition simple et

un moyen de substitution (16) sensible à chaque signal de synchronisation (V_A) lorsqu'il présente le niveau de signal spécifique pour remplacer la valeur de chaque bit associé dans le premier jeu ($B_0..B_{K-1}$) soit avec la valeur du bit correspondant dans le deuxième jeu ($B_K..B_{M-1}$) soit avec la valeur du complément binaire du bit correspondant dans le deuxième jeu ($B_K..B_{M-1}$) de façon à satisfaire à l'algorithme prescrit.

2. Circuit selon la revendication 1, caractérisé en ce que les moyens de substitution effectuent chaque remplacement dans une gamme partielle utilisant soit la valeur d'un bit simple du deuxième jeu ($B_K..B_{M-1}$) soit la valeur du complément binaire de ce bit simple.

3. Circuit selon la revendication 1, caractérisé en ce qu'il y a des premiers (10) et deuxième (12) trajets d'engendrement de bit à travers l'étage principal (10, 12) chaque bit dans le premier jeu ($B_0..B_{K-1}$) étant engendré le long du premier trajet (10), chaque bit dans le deuxième jeu ($B_K..B_{M-1}$) étant engendré le long du deuxième trajet (12) et en ce que le niveau de signal spécifique pour chaque signal de synchronisation (V_A) consiste dans le fait qu'il est supérieur ou inférieur à une tension de seuil.

4. Circuit selon la revendication 3, caractérisé en ce qu'il y a qu'une gamme partielle pour chaque signal de synchronisation (V_A).

5. Circuit selon la revendication 3, dans lequel il y a au moins deux bits dans chaque jeu, caractérisé en ce que le moyen de substitution (16) effectue chaque remplacement dans une gamme partielle utilisant soit la valeur d'un seul des bits du deuxième jeu ($B_K..B_{M-1}$) soit la valeur du complément binaire de ce seul bit.

6. Circuit selon la revendication 1, caractérisé en ce que l'étage principal comprend des moyens (18)



sensibles aux paramètres d'entrée (V_i) pour produire plusieurs signaux intermédiaires ($V_{AO}-V_{A83}$) qui varient entre des premiers (V_L) et deuxièmes (V_H) niveaux lorsque le paramètre d'entrée (V_i) traverse la gamme d'entrée ($V_{RO}-V_{R83}$), chaque signal intermédiaire ($V_{AO}-V_{A83}$) étant au premier niveau (V_L) pour la plupart des passages du paramètre d'entrée (V_i) par la gamme d'entrée ($V_{RO}-V_{R83}$) et atteignant le deuxième niveau (V_H) à une autre valeur (V_{Ri-4}) du paramètre d'entrée (V_i) que chaque autre signal intermédiaire et des moyens (20, 24, 26, 28) sensibles aux signaux intermédiaires ($V_{AO}-V_{A83}$) pour engendrer chaque bit dans le deuxième jeu ($B_2..B_7$); et

l'étage de synchronisation comprend une partie (18, 16) de l'étage principal, chaque signal de synchronisation (V_{A18} , V_{A35} , V_{A51}) étant l'un des signaux intermédiaires ($V_{AO}-V_{A83}$).

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