

METHOD AND SYSTEM FOR DISTRIBUTING ENCRYPTION AND
DECRYPTION KEYS IN
SECURE BROADCAST COMMUNICATION

5 Background of the Invention

1. Field of the Invention

The present invention relates to a technique of distributing a key for decryption in secure broadcast communication to restricted receivers.

2. Description of Related Art

10 Conventionally, as techniques of secure broadcast communication, some systems have been proposed.

For example, there is known a copied key method described in an article, S.J. Kent: "Security Requirements and Protocols for a Broadcast Scenario", IEEE Trans. Commun., COM-29, 6, pp. 778 - 786 (1981).

15 This copied key method is a fundamental system for secure broadcast communication, and is simple extension of the conventional one-to-one individual cipher communication. Namely, in this method, copies of one kind of key are distributed to a sender and a plurality of regular receivers. The sender encrypts information by use of the distributed copy key and sends
20 the encrypted information. Each of the regular receivers uses the distributed copy key to decrypt the encrypted information.

Further, there is known a technique of distributing a common key using an individual key assigned for every receiver. For example, a system of key distribution by multiplexing-demultiplexing of an information
25 sequence using the Chinese remainder theorem is described in an article, Lee, Tokiwa, et al: "A Multi-Address Communication Using a Method of



Multiplexing and Demultiplexing, 1986 Symposium on encryption and information security". Also, Mambo et al: "A Secure Broadcast Communication Method with Short Messages", IEICE TRNS. FUNDAMENTALS, VOL. E77-A, No. 8 AUGUST 1994 describes such a
5 technique.

The system of multiplexing-demultiplexing an information sequence by use of the Chinese remainder theorem performs the following processes.

(1) Key generation process:

10 For receivers i ($1 \leq i \leq s$), s integers $g_1, g_2, \dots, g_s$ ($r \leq s$) which are relatively prime to one another are generated, and g_i is distributed in advance to i -th receiver as secret key of that receiver i .

(2) Encryption process:

An information sequence of s items to be multiplexed is expressed
15 as $M_1, M_2, \dots, M_s$. The sender calculates a multiplexed sentence F to be sent by:

$$F = \sum_{i=1}^k A_i G_i M_i \text{ mod } G$$

and broadcasts it. Here, G , G_i , and A_i are defined as follows:

$$\begin{aligned} G &= \prod_{i=1}^k g_i \\ G_i &= G / g_i \\ A_i G_i &\equiv 1(\text{mod } g_i) \end{aligned}$$

20 and A_i is the smallest integer which satisfies the above equations.

(3) Decryption process:

Using g_i , the receiver i demultiplexes M_i from F by the following equation,



$$M_i = F \bmod g_i$$

Here, M_i is a common key to be distributed to the receiver i . Thus, according to this system, the copy key can be distributed in secret to the restricted receivers only.

5 Next, the system described in the literature, Mambo et al: "A Secure Broadcast Communication Method with Short Messages", IEICE TRNS. FUNDAMENTALS, VOL. E77-A, No. 8 AUGUST 1994 performs following processes.

(1) Key generation process:

10 A trusted center produces following information.

Secret key:

$$P = 2p + 1, Q = 2q + 1: \text{prime number } (p, q: \text{prime number})$$

$$e_i \in \mathbb{Z}, 0 < e_i < L (1 \leq i \leq m)$$

Public key:

$$g \in \mathbb{Z}, 0 < g < N$$

$$N = PQ$$

$$v_i = g^{e_i} \bmod N (1 \leq i \leq m)$$

15 For $\sigma \in S$, the center calculates s_σ that satisfies:

$$s_\sigma \sum_{i=1}^k e_{\sigma(i)} \equiv 1 \pmod{L}$$

and distributes it as secret key of a receiver U_σ . Here, the set S is defined as $S = \{f \mid \text{one to one map } f: A\{1, 2, \dots, k\} \rightarrow B = \{1, 2, \dots, m\}, m > k\}$.

(2) Key distribution process:

20 A sender selects an integer r at random, and, in order to have a common key K defined as:

$$K = g^r \bmod N$$

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in common with restricted receivers, calculates z_i that satisfies:

$$z_i = v_i' \bmod N \quad (1 \leq i \leq m)$$

and broadcasts z_i ($1 \leq i \leq m$).

The receiver $U\sigma$ uses the following equation:

5

$$K = \left(\prod_{i=1}^k z_{\sigma(i)} \right)^{s\sigma} \bmod N$$

to calculate the common key K .

Summary of the Invention

In the above-described conventional method of key distribution by
10 multiplexing using the Chinese remainder theorem, common key data for
respective receivers are arranged serially and transmitted. Thus, broadcast
data length becomes longer in proportion to the number of the receivers.
Accordingly, it is not suitable for communication for millions of receivers or
more, such as satellite broadcasting.

15 On the other hand, according to the system described in the above-
referred literature, Mambo et al: "A Secure Broadcast Communication
Method with Short Messages", IEICE TRNS. FUNDAMENTALS, VOL. E77-
A, No. 8 AUGUST 1994, data for distributing the common key can be
shortened. In this system, however, it is not possible to have the key
20 commonly used among restricted receivers belonging to any set of receivers.

Further, in all systems, when there are a plurality of senders,
receivers must obtain the above-described user secret key for respective
senders, and manage them.

Thus, an object of the present invention is to provide a key
25 distribution system in which a receiver can receive delivery of respective



common keys from a plurality of senders by means of only one secret key. A further object is to make it possible that, in such a key distribution system, a common key for data decryption can be commonly owned only among any sender and any set of receivers, and to shorten broadcast data for distributing
5 a common key even when the number of receivers are large.

To attain the above objects, the present invention provides a method of key distribution having an encryption-decryption key of secure broadcast communication commonly used among a plurality of senders and a plurality of receivers in a communication system comprising apparatuses of
10 said plurality of senders and apparatuses of said plurality of receivers, wherein:

in an apparatus of a trusted center provided separately from said senders and said receivers, first key information, which is common to said plurality of senders and is used for a key commonly used among said senders
15 and said receivers, is produced and distributed to prescribed receivers, while second key information which is used for a key commonly used among said senders and said receivers is produced and distributed to said senders;

in apparatuses of said senders, third key information, which is used by said receivers together with said first key information to calculate an
20 encryption-decryption key used by the sender in question in broadcast communication, is produced using the second key information distributed from said trusted center and is sent to said receivers who have a key in common with said sender; and

in an apparatus of a receiver, the encryption-decryption key which
25 is used for broadcast communication by the sender who distributed said third key information is obtained using the first key information distributed from



said trusted center and the third key information distributed from said sender.

According to the present method of key distribution, it is sufficient for each receiver to have only the first key information distributed from the trusted center as a user secret key. When an encryption-decryption key of a new sender is to be obtained, it is not necessary to receive distribution of new user secret key.

In detail, the present invention provides a method of key distribution, in which an encryption-decryption key for secure broadcast communication performed by said sender is distributed to receivers, in a communication system comprising apparatuses of a plurality of senders and apparatuses of a plurality of receivers, and an apparatus of a trusted center, comprising:

a step, in which, in the apparatus of the trusted center,

$e_i \in Z \quad (1 \leq i \leq m)$

are produced as trusted center secret key,

$t_A \in Z$

is produced as sender registration key of a sender A, and, as receiver secret key, $\sigma_x \in S_{km}$ and

$s_x(\sigma_x) \in Z$

are produced (here, when $\sigma, \sigma' \in S_{km}$ for a set S_{km} defined as:

$$S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m \}$$

it is expressed as:

$$\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A),$$



with “ $\sim$ ” becoming an equivalence relation on S_{km} , leading to an expression:

$$\tilde{S}_{km} = S_{km} / \sim$$

), and the receiver secret key σ_x , $s_x(\sigma_x)$ are distributed to the receiver x;

a step, in which, in the apparatus of the sender A,

5 sender secret key g_A , L_A , and a finite Abelian group G_A satisfying:

$$\begin{aligned} g_A &\in G_A \\ L_A &= \text{ord}_{G_A}(g_A) \end{aligned}$$

are produced (here,

$$\text{ord}_{G_A}(g)$$

means the least positive integer a satisfying:

10
$$g^a = 1 (\in G_A)$$

), and the sender secret key g_A is sent to said trusted center;

a step, in which, in the apparatus of the trusted center, receiver registration data:

$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

15 is calculated from the trusted center secret key e_i , the sender registration key t_A , and the receiver secret key $s_x(\sigma_x)$, σ_x , and the receiver registration data $s_x(\sigma_x, A)$ is sent to the sender A, and key distribution data:

$$y_{Ai} = g_A^{t_A e_i} (\in G_A) \quad (1 \leq i \leq m)$$

is calculated from g_A received from the sender A, the trusted center secret key e_i , the sender registration key t_A , and this key distribution data y_{Ai} is sent to
20 the sender A;

a step, in which, in the apparatus of the sender A, random integers r , r' are produced, receiver registration key $r_x(\sigma_x, A)$ satisfying:

$$r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$



is calculated from the receiver registration data $s_x(\sigma_x, A)$ received from said trusted center, the sender secret key L_A , and the random integer r' ; the receiver registration key $r_x(\sigma_x, A)$ is sent to the receiver x ; and key distribution data defined as:

$$5 \quad z_{Ai} = y_{Ai}^{r'} \quad (\in G_A) \quad (1 \leq i \leq m)$$

is calculated from the key distribution data y_{Ai} received from said trusted center, and the random integer r , and the key distribution data z_{Ai} is broadcasted to each receiver; and

a step, in which, in the apparatus of the receiver x , an encryption-
10 decryption key K_A used for broadcast communication is calculated using:

$$K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A)s_x(\sigma_x)} \quad (\in G_A)$$

from the receiver registration key $r_x(\sigma_x, A)$ received from the sender, the key distribution data z_{Ai} , and the receiver secret key $\sigma_x, s_x(\sigma_x)$, with the encryption-decryption key K_A being produced by the sender A , using:

$$15 \quad K_A = g_A^{r'} \quad (\in G_A).$$

According to this method of key distribution, it is not necessary for the receivers to own private keys which are different for respective senders. Further, in secure broadcast communication, even when the number of receivers is large, the length of the key distribution data can be made short.

20 Further, in the present method of key distribution, the receiver secret key is made secret from the senders and the sender secret key is more reliably made secret from the receivers since the sender registration key is owned only by the trusted center in secret. Accordingly, security against impropriety is improved.

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Brief Description of the Drawings

Fig. 1 is a block diagram showing a configuration of a key distribution system according to an embodiment of the present invention;

Fig. 2 is a block diagram showing a configuration of a trusted center side apparatus according to a first embodiment of the present invention;

Fig. 3 is a block diagram showing a configuration of a sender side apparatus according to the first embodiment of the present invention;

Fig. 4 is a block diagram showing a configuration a receiver side apparatus according to the first embodiment of the present invention;

Fig. 5 is a view showing information flow caused by key distribution process in the first embodiment of the present invention;

Fig. 6 is a block diagram showing a configuration of a receiver side apparatus according to a second embodiment of the present invention;

Fig. 7 is a block diagram showing a configuration of a trusted center side apparatus according to a fifth embodiment of the present invention;

Fig. 8 is a block diagram showing a configuration of a sender side apparatus according to the fifth embodiment of the present invention; and

Fig. 9 is a block diagram showing a configuration of a receiver side apparatus according to the fifth embodiment of the present invention.

Detailed Description

In the following, there will be described a key distribution system according to an embodiment of the present invention.

Firstly, the first embodiment will be described.

Fig. 1 shows a configuration of a key distribution system according to the present embodiment of the invention.

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As shown in the figure, the present system comprises a trusted center side apparatus 100, sender side apparatuses 200, and receiver side apparatuses 300. These apparatuses are connected to one another via communication lines 400. The trusted center side apparatus 100 is the one
5 and only apparatus used by the trusted center organization in the system, while the system has a plurality of sender side apparatuses 200 and a plurality of receiver side apparatuses 300.

Fig. 2 shows a configuration of the trusted center side apparatus 100.

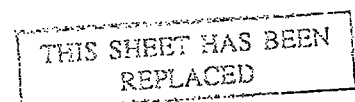
10 As shown in the figure, the trusted center side apparatus 100 comprises a random number generator 101, a prime number generator 102, a power multiplier 103, a remainder computing unit 104, an arithmetic unit 105, a memory 106, and a communication unit 107. The trusted center side apparatus 100 is connected to receiver side portable apparatuses 306 which
15 are delivered to receivers off line.

Fig. 3 shows a configuration of the sender side apparatus 200.

As shown in the figure, the sender side apparatus 200 comprises a random number generator 201, a prime number generator 202, a power multiplier 203, a remainder computing unit 204, an arithmetic unit 205, a
20 memory 206, a communication unit 207, an encryption-decryption unit 208, an authentication unit 209, and an accounting unit 210.

Fig. 4 shows a configuration of the receiver side apparatus 300.

As shown in the figure, the receiver side apparatus 300 comprises a power multiplier 301, a remainder computing unit 302, an arithmetic unit
25 303, a memory 304, a communication unit 305, a receiver side portable apparatus 306 which has been delivered off line from the trusted center, an



encryption-decryption unit 307, and an authentication unit 308.

In the following, three processes will be described, i.e., a preparation process, a key distribution process, and an encryption-decryption process.

5 First, the preparation process will be described.

(1) Preparation process

(i) A sender A uses the random number generator 201, the prime number generator 202, the power multiplier 203, the remainder computing unit 204 and the arithmetic unit 205 within the sender side apparatus 200, to
10 produce the following keys, and makes only public key open to the public.

Secret key:

$$\begin{aligned} P_A, Q_A : & \text{ prime numbers} \\ L_A = & \text{lcm}(\text{ord}_{P_A}(g_A), \text{ord}_{Q_A}(g_A)) \\ g_A \in & \mathbb{Z}, \quad 0 < g_A < N_A \\ r, r' \in & \mathbb{Z}, \quad 0 < r, r' < L_A \end{aligned}$$

Public key:

$$N_A (= P_A Q_A)$$

15 The secret key is stored in the memory 206. Further, pieces of the secret key g_A, L_A are sent to the trusted center by means of the communication unit 207.

(ii) The reliable trusted center uses the arithmetic unit 105 within the trusted center side apparatus 100 to generate the following information.

20 Trusted center secret key:

$$e_i \in \mathbb{Z} \quad (1 \leq i \leq m)$$

Sender registration key for the sender A:

$$t_A \in \mathbb{Z}, \quad 0 < t_A < L_A$$

Receiver secret key for a receiver x:



$$s_x(\sigma_x) \in Z, \quad 0 < t_A < L_A$$

All these are stored in the memory 106 together with σ_x .

Here, for a set:

$$S_{km} = \{\sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m\}$$

5 when $\sigma, \sigma' \in S_{km}$, it is assumed that:

$$\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A)$$

Here, " $\sim$ " becomes an equivalence relation on S_{km} , and

$$\tilde{S}_{km} = S_{km} / \sim$$

Further, the trusted center takes the receiver secret key $s_x(\sigma_x)$ from
10 the memory 106, stores it in the receiver side portable apparatus 306
together with σ_x , and delivers that apparatus to the receiver x off line. Of
course, it may be delivered to the receiver in another way.

Next, the key distribution process will be described.

(2) Key distribution process

15 Fig. 5 shows information flow among the trusted center, senders
and receivers in the present key distribution process.

(i) The trusted center uses the remainder computing unit 104 and the
arithmetic unit 105 within the trusted center side apparatus 100 to calculate
receiver registration key $s_x(\sigma_x, A)$ defined as:

$$20 \quad s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)} \bmod L_A$$

from L_A received from the sender A, trusted center secret key e_i , σ_x , receiver
secret key $s_x(\sigma_x)$ and the sender registration key t_A , and sends it to the sender
A using the communication unit 107.



Further, the trusted center uses the power multiplier 103, the remainder computing unit 104 and the arithmetic unit 105 to calculate sender key distribution data y_{Ai} of the sender A:

$$y_{Ai} = g_A^{u_i} \bmod N_A \quad (1 \leq i \leq m)$$

5 from g_A received from the sender A, sender public key N_A of the sender A, and trusted center public key e_i , and sends it to the sender A by means of the communication unit 107.

(ii) The sender A generates random integers r, r' by means of the random number generator 201 within the sender side apparatus 200, and
10 stores them in the memory 206. Further, the sender A uses the power multiplier 203, the remainder computing unit 204 and the arithmetic unit 205 to calculate data encryption key K_A defined by:

$$K_A = g_A^{r'} \bmod N_A$$

from these random integers r, r' , his own secret key g_A and his own public key
15 N_A , and stores it in the memory 206. Then, for the purpose of having the key K_A in common with the receiver x , the sender A uses the remainder computing unit 204 and the arithmetic unit 205 to calculate receiver registration key $r_x(\sigma_x, A)$ of the receiver x , satisfying:

$$r_x(\sigma_x, A)s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

20 from his own secret key L_A , the random number r' and the receiver registration key $s_x(\sigma_x, A)$, and sends it to the receiver by means of the communication unit 207.

Further, the sender A uses the power multiplier 203, the remainder computing unit 204 and the arithmetic unit 205 within the sender side
25 apparatus 200 to calculate receiver key distribution data z_{Ai} defined by:

$$z_{Ai} = y_{Ai}^r \bmod N_A \quad (1 \leq i \leq m)$$



from his own public key N_A , the random number r and y_{Ai} received from the trusted center, and broadcasts it to the receivers by means of the communication unit 207.

In the above process, the trusted center performs a part of the processes for producing $r_x(\sigma_x, A)$ and z_{Ai} , for one thing, to keep σ_x of the receiver x , the receiver secret key $s_x(\sigma_x)$ and the sender registration number t_A secret from the sender A so as to prevent improprieties by the sender A .

(iii) The receiver x uses the power multiplier 301, the remainder computing unit 302 and the arithmetic unit 303 within the receiver side apparatus 300 to calculate the data encryption key K_A by the expression:

$$K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A) s_x(\sigma_x)} \mod N_A$$

from the key distribution data z_{Ai} received from the sender A , the receiver registration key $r_x(\sigma_x, A)$, σ_x and the receiver secret key $s_x(\sigma_x)$ delivered from the trusted center and the sender public key N_A opened to the public, and stores it in the memory 304.

By the above-described process, the sender A and the receiver x can have the key K_A in common. Similarly, for any other sender and any other receiver, it is possible to have a key in common. In that case, receiver key distribution data z_{Ai} broadcasted by each sender is the same for each receiver.

Further, in the above-described process, the receiver x receives the key distribution data z_{Ai} and the receiver registration key $r_x(\sigma_x, A)$ from the sender A . These z_{Ai} and $r_x(\sigma_x, A)$ have been produced under the effect of the sender registration number t_A of the sender A , and t_A is kept secret from the receiver. Thus, owing to the use of the above-described sender registration number t_A , it becomes very difficult for the receiver x to deduce a data



encryption key K_B of another sender B from the key distribution data z_{Ai} and the receiver registration key $r_x(\sigma_x, A)$ received from the sender A and key distribution data z_{Bi} broadcasted by the sender B.

In the following, the encryption-decryption process will be
5 described.

(3) Encryption-decryption process

(i) The sender A uses the encryption-decryption unit 208 within the sender side apparatus 200 to encrypt data P with the common key K_A produced in the key distribution process. Then, the sender A sends a cipher
10 text $C=E(K_A: P)$ to a receiver by means of the communication unit 207.

(ii) The receiver uses the communication unit 305 within the receiver side apparatus 300 to receive the public key C, and uses the encryption-decryption unit 307 to decrypt the public key C with the common key K_A stored in the memory 304, so as to obtain the original data.

15 The above is the first embodiment of the present invention.

In the conventional technique, when a sender newly enters a key distribution system, it is necessary for the sender to produce σ_x and receiver secret key $s_x(\sigma_x)$ by himself and to deliver them to a receiver off line. On the other hand, according to the key distribution system of the first embodiment,
20 it suffices for a sender who newly enters the system to produce the sender secret key P_A, Q_A, L_A, g_A and the sender public key N_A . Further, receiver secret key owned by a receiver is same for all the senders. Thus, when a receiver wishes to receive data from a new sender, it is not necessary to obtain new receiver secret key.

25 Further, the receiver registration key $r_x(\sigma_x, A)$ makes it possible to have a common key in common with receivers belonging to any set of



receivers, for data encryption and decryption concerning them only. Even when the number of receivers is large, it is not necessary to correspondingly lengthen the key distribution data z_{Ai} broadcasted for distribution of a common keys.

5 In the above processes, it is possible alternatively that the sender sends the sender secret key L_A to the trusted center in advance, and in the step of the preparation process, the trusted center generates the trusted center secret key e_i , the sender registration key t_A , and the receiver secret key $s_x(\sigma_x)$ satisfying:

$$\begin{aligned} e_i &\in Z, \quad 0 < e_i < L_A, \quad (1 \leq i \leq m) \\ t_A &\in Z, \quad 0 < t_A < L_A \\ 10 \quad s_x(\sigma_x) &\in Z, \quad 0 < s_x, \sigma_x < L_A, \end{aligned}$$

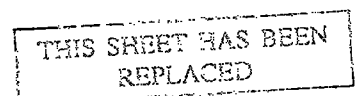
and, in the step of the key distribution process, the trusted center calculates the receiver registration key $s_x(\sigma_x, A)$ by:

$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)} \mod L_A$$

from the sender secret key L_A , the trusted center secret key e_i , the sender registration key t_A and the receiver secret key $s_x(\sigma_x)$.

Next, a second embodiment of the present invention will be described.

In the second embodiment of the present invention, the receiver side portable apparatus 300 of the above first embodiment is provided with a power multiplier 3064, a remainder computing unit 3063, an arithmetic unit 20 3061 and a memory 3062 as shown in Fig. 6. Further, the process of calculating the data encryption key K_A in the receiver side apparatus 300 of the above first embodiment is partly performed in the receiver side portable



apparatus 306.

Namely, in the second embodiment, (1) in the step of the preparation process, the trusted center stores σ_x and the receiver secret key $s_x(\sigma_x)$ in the receiver side portable apparatus 306 (such as an IC card (a smart
5 card)), and delivers them to the receiver x.

Then, (2) in the key distribution process, the receiver x outputs the key distribution data z_{Ai} stored in the memory within the receiver side apparatus 300 to the receiver side portable apparatus 306. Next, within the receiver side portable apparatus 306, the receiver x uses the power multiplier
10 3064 and the remainder computing unit 3063 to calculate:

$$\xi_x(\sigma_x, A) = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{s_x(\sigma_x)} \bmod N_A$$

from σ_x , the receiver secret key $s_x(\sigma_x)$, the key distribution data z_{Ai} and the sender public key N_A , and outputs the calculated result $\xi_x(\sigma_x, A)$ to the receiver side apparatus 300.

15 Next, the receiver x uses the power multiplier 301, the remainder computing unit 302 and the arithmetic unit 303 within the receiver side apparatus 300 to calculate the data encryption key K_A by:

$$K_A = \xi_x(\sigma_x, A)^{r_x(\sigma_x, A)} \bmod N_A$$

from $\xi_x(\sigma_x, A)$ outputted to the receiver side apparatus 300, the receiver
20 registration key $r_x(\sigma_x, A)$ stored in the memory 304, and the sender public key N_A , and stores it in the memory 304.

Thus, σ_x and the receiver secret key $s_x(\sigma_x)$ are prevented from being outputted outside the receiver side portable apparatus 306, so that these are prevented from being stolen by electronic duplicating or the like.

25 The above is the second embodiment of the present invention.



Next, a third embodiment of the present invention will be described.

The third embodiment of the present invention is different from the above-described first embodiment in that the value of r in the equation:

$$K_A = g_A^{r'} \bmod N_A$$

5 is changed periodically every short time period, and z_{Ai} obtained by using the changed r is broadcasted periodically for update of the data encryption-decryption key K_A calculated in the sender side apparatus 200 and in the receiver side apparatus 300.

Further, in the present embodiment, the value of r' is made to be
10 characteristic to sending data so as to identify data broadcasted by a sender. Namely, a registration key $r_x(\sigma_x, A)$ which the receiver x received from a sender to decrypt certain broadcast data or a set of broadcast data is characteristic to that broadcast data. To obtain another broadcast data or another set of broadcast data, the receiver x must receive another $r_x(\sigma_x, A)$ for
15 that broadcast data or that set of broadcast data.

The above is the third embodiment of the present invention.

Next, a fourth embodiment of the present invention will be described.

The fourth embodiment is different from the above-described first
20 embodiment in that authentication of receivers is performed so as to charge receivers who have a common key K_A in common for fee-charging data P encrypted with the key K_A and sent by a sender.

Namely, in the fourth embodiment of the present invention, following processes are further performed.

25 (1) Preparation process

(i) The trusted center uses the arithmetic unit 105 within the trusted



center side apparatus 100 to produce a number UID_x for the receiver x in advance, and stores it in the receiver side portable apparatus 306 together with the receiver secret key $s_x(\sigma_x)$, and delivers that apparatus. Further UID_x is stored in the memory 106, so as to correspond to $s_x(\sigma_x)$.

5 (ii) The sender A uses the arithmetic unit 205 within the sender side apparatus 200 to produce his own number $BIDA$, and stores it in the memory 206. The sender A sends $BIDA$ to the trusted center by means of the communication unit 207.

(iii) The trusted center receives $BIDA$, the number of the sender A ,
10 using the communication unit 107 within the trusted center side apparatus 100, and stores it in the memory 106 so as to correspond to the sender registration key t_A .

(iv) The receiver x uses the authentication unit 308 within the receiver side apparatus 300 to produce information for authentication from the
15 receiver secret key $s_x(\sigma_x)$, and sends it to the sender A .

(v) The sender A confirms the information for authentication by means of the authentication unit 209 within the sender side apparatus 200.

As the method of authentication, any conventionally-known authentication method may be employed in which authentication is not
20 effected as long as the receiver does not know $s_x(\sigma_x)$. However, it is necessary for the sender to be prevented from knowing the receiver private key $s_x(\sigma_x)$ itself.

For example, authentication of the receiver x may be performed as in the following processes according to a method employing a signature using
25 the RSA algorithm (described in an article by, R. L. Rivest, A. Shamir, L. Adelman: "A Method for Obtaining Digital Signatures and Public Key



Cryptosystems", Commun. of the ACM, Vol. 21, No. 2, pp. 120-126, 1978).

(1) The trusted center uses the arithmetic unit 105 within the trusted center side apparatus 100 to produce (y_x, n_x) satisfying:

$$\begin{aligned} s'_x y_x &\equiv 1 \pmod{\text{lcm}(p_x - 1, q_x - 1)} \\ n_x &= p_x q_x \quad (p_x, q_x : \text{prime numbers}) \end{aligned}$$

5 for the receiver x , and deliver them to the sender in advance. Here, for a function π which has been opened, s'_x is defined as $s'_x = \pi(s_x(\sigma_x))$.

(2) The receiver x uses the authentication unit 308 within the receiver side apparatus 300 to calculate a hash value $(h(w); 0 < h(W) < n_x)$ of broadcast data W using a unidirectional hash function h as public key, and produces a
10 signature for $h(W)$ by the following expression:

$$\text{sgn}_x(h(W)) = h(W)^{s'_x} \pmod{n_x}$$

using the private key s'_x , and sends it to the sender by means of the communication unit 305 together with a data sending request.

(3) The sender uses the authentication unit 209 within the sender side
15 apparatus 200 to confirm that:

$$\text{sgn}_x(h(W))^{y_x} = h(W) \pmod{n_x}$$

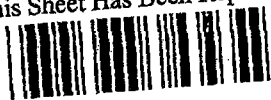
is satisfied.

After the authentication, the sender A uses the communication unit 207 within the sender side apparatus 200 to send his own number BIDA and
20 UID $_x$, the number of the receiver x , to the trusted center.

The trusted center uses the arithmetic unit 105 within the trusted center side apparatus 100 to calculate the receiver registration data $s_x(\sigma_x, A)$ expressed by:

$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)} \pmod{LA}$$

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from the sender registration key t_A and the receiver secret key $s_x(\sigma_x)$ corresponding to the sender number BIDA and the receiver number UIDx received respectively, and sends it to the sender A.

The sender uses the remainder computing unit 204 and the
5 arithmetic unit 205 within the sender side apparatus 200 to calculate the receiver registration key $r_x(\sigma_x, A)$ satisfying:

$$r_x(\sigma_x, A)s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

from the receiver registration key $s_x(\sigma_x, A)$ received from the trusted center, and sends it to the receiver by means of the communication unit 207. At
10 that time, in the case that data to be delivered being encrypted with the common key K_A (namely, the data which the receiver requested to the sender) is chargeable, the sender A charges the receiver x by means of the accounting unit 210.

The above is the fourth embodiment of the present invention.

15 Next, a fifth embodiment of the present invention will be described.

The fifth embodiment of the present invention is an extension of the copy key system described in the Description of Related Art to the system shown in Fig. 1 which includes a plurality of senders, a plurality of receivers and a trusted center.

20 Fig. 7 shows a configuration of a trusted center side apparatus 100 according to the fifth embodiment of the present invention. As shown in the figure, the trusted center side apparatus 100 comprises a random number generator 111, an arithmetic unit 112, a memory 113, and a communication unit 114. Further, Fig. 8 shows a configuration of a sender side apparatus
25 200 according to the fifth embodiment of the present invention. As shown in the figure, the sender side apparatus 200 comprises a random number



generator 211, an arithmetic unit 212, an encryption-decryption unit 213, a memory 214, and a communication unit 215. Further, Fig. 9 shows a configuration of a receiver side apparatus 300 according to the fifth embodiment of the present invention. As shown in the figure, the receiver side apparatus 300 comprises an arithmetic unit 311, an encryption-decryption unit 312, a memory 313, and a communication unit 314.

In the present embodiment, firstly, the following process is performed as a preparation process.

(1) Preparation process

10 (i) The trusted center uses the random number generator 111 within the trusted center side apparatus 100 to generate a secret key K_0 and stores it in the memory 113. Also, the trusted center distributes it to senders and receivers.

(ii) The sender A uses the random number generator 211 within the sender side apparatus 200 to generate sender registration key BIDA, and sends it to the receivers by means of the communication unit 215.

Thereafter, the following process is performed as a key distribution process.

(2) Key distribution process

20 (i) The sender A uses the arithmetic unit 212 within the sender side apparatus 200 to calculate a common key $KA' = F(K_0, BIDA)$ from the secret key K_0 and the sender registration key BIDA, using a suitable one-way function F .

(ii) The receiver uses the arithmetic unit 312 within the receiver side apparatus 300 to calculate the common key $KA' = F(K_0, BIDA)$ from the secret key K_0 distributed from the trusted center and the sender registration key



BIDA received from the sender, using the one-way function F .

(iii) The sender A uses the random number generator 211 within the sender side apparatus 200 to produce a suitable integer r and sends it to the receiver by means of the communication unit 215. Further, the sender A
5 calculates a data encryption-decryption key $DK=F'(r, KA)$ from the integer r and the common key KA using a suitable function F' , and stores it in the memory 214.

(iv) The receiver uses the arithmetic unit 312 within the receiver side apparatus 300 to calculate the data encryption-decryption key $DK=F'(r, KA)$
10 from the integer r received from the sender A and the common key KA , and stores it in the memory 314.

After the distribution of the data encryption-decryption key DK by the key distribution process, an encryption-decryption process is performed as described in the following.

15 (3) Encryption-decryption process

(i) The sender A encrypts data P to be delivered with the data encryption-decryption key DK using the encryption-decryption unit 213 within the sender side apparatus 200, and sends it to the receiver x by means of the communication unit 215.

20 (ii) The receiver receives the encrypted delivered data P by the communication unit 314 within the receiver side apparatus 300, and decrypts it with the data encryption-decryption key DK using the encryption-decryption unit 312.

In the above-described fifth embodiment of the present invention
25 also, a single piece of secret key suffices for a receiver to have respective common keys of a plurality of senders in common with them.



Next, a sixth embodiment of the present invention will be described.

In the sixth embodiment of the present invention, senders and receivers have a common key in common among them, through cipher communication using respective keys of the receivers.

5 Configurations of the whole system, a trusted center side apparatus 100, a sender side apparatus 200, and a receiver side apparatus 300 according to the present embodiment are similar to those in the above-described fifth embodiment of the present invention.

 In the present embodiment, the following process is performed as a
10 preparation process.

(1) Preparation process

(i) The trusted center uses the random number generator 111 within the trusted center side apparatus 100 to produce secret key s_x of the receiver x , and delivers it to the receiver. Further, the trusted center uses the
15 random number generator 111 to produce BIDA, the number of the sender A , and distributes it to the sender A and the receiver x .

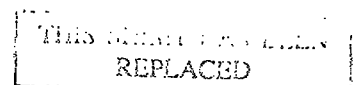
(ii) The sender A uses the random number generator 211 within the sender side apparatus 200 to produce the common key K_A .

 Thereafter, the following process is performed as a key distribution
20 process.

(2) Key distribution process

(i) The trusted center uses the arithmetic unit 112 within the trusted center side apparatus 100 to calculate a session key K_{Ax} , which is to be used in a session between the receiver x and the sender A and defined as $K_{Ax}=F(s_x,$
25 $K_A)$ using a suitable one-way function F . The trusted center calculates K_{Ax} from BIDA, the number of the sender A , and the receiver secret key s_x , and

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stores it in the memory 113.

(ii) The receiver x uses the arithmetic unit 312 within the receiver side apparatus 300 to calculate $K_{Ax}=F(s_x, BIDA)$ from the number of the sender A BIDA and his own secret key s_x using the one-way function F , and stores it in the memory 313.

(iii) The sender A uses the encryption-decryption unit 213 within the sender side apparatus 200 to encrypt the common key K_A with the session key K_{Ax} to obtain key distribution data K_{Cx} , and sends the key distribution data K_{Cx} to the receiver by means of the communication unit 215. Further, the sender A uses the random number generator 211 to produce a suitable integer r , and sends it to the receiver by means of the communication unit 215. In addition, the sender A uses the arithmetic unit 212 to calculate a data encryption-decryption key $DK=F'(r, K_A)$ from the integer r and the common key K_A using a suitable function F' , and stores it in the memory 214.

(iv) The receiver x receives the key distribution data K_{Cx} from the sender A, and uses the encryption-decryption apparatus 303 within the receiver side apparatus 300 to decrypt the key distribution data K_{Cx} with the session key K_{Ax} . Then, the receiver x uses the arithmetic unit 302 to calculate the data encryption-decryption key $DK=F'(r, K_A)$ from the decrypted common key K_A and the received integer r , using the function F' , and stores it in the memory 303.

After the distribution of the data encryption-decryption key DK by the key distribution process, an encryption-decryption process is performed as described in the following.

(3) Encryption-decryption process

(i) The sender A uses the encryption-decryption unit 203 within the

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sender side apparatus 200 to encrypt data P to be delivered with the data encryption-decryption key DK, and sends it to the receiver x by means of the communication unit 205.

(ii) The receiver x receives the encrypted delivered data P through the communication unit 305 within the receiver side apparatus 300. Then, the receiver x uses the encryption-decryption unit 303 to decrypt the received data P with the data encryption-decryption key DK.

The sender A may change the value of the common key K_A periodically so as to change the value of the data encryption-decryption key DK periodically.

According to the above-described sixth embodiment of the present invention also, the receiver can have respective common keys of a plurality of senders in common with them, by means of single piece of secret key.

Next, a seventh embodiment of the present invention will be described.

Similarly to the above-described sixth embodiment, in the seventh embodiment of the present invention, senders and receivers have a common key in common among them by cipher communication using respective keys of the receivers.

Configurations of the whole system, a trusted center side apparatus 100, a sender side apparatus 200, and a receiver side apparatus 300 according to the present embodiment are similar to those in the above-described fifth embodiment of the present invention.

In the present seventh embodiment, the following process is performed as a preparation process.

(1) Preparation process

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(i) The trusted center uses the random number generator 111 and the arithmetic unit 112 within the trusted center side apparatus 100 to produce a private key s_x and public key p_x of the receiver x according to a suitable open key cryptographic system E and sends s_x to the receiver. Further, the
5 trusted center produces $BIDA$, a number of the sender A , using the random number generator 111, and delivers or sends it to the sender A and the receiver x .

(ii) The sender A uses the random number generator 211 within the sender side apparatus 200 to produce a common key K_A , and stores it in the
10 memory 114.

Thereafter, the following process is performed as a key distribution process.

(2) Key distribution process

(i) The sender A uses the encryption-decryption unit 213 within the sender side apparatus 200 to calculate a key distribution data $K_{Cx}=E(p_x, K_A)$
15 from the receiver public key p_x and the common key K_A employing a suitable above-described public key cipher system E . Then, the sender A sends K_{Cx} to the receiver x by means of the communication unit 215.

(ii) The receiver x uses the encryption-decryption unit 313 within the receiver side apparatus 300 to encrypt, with his own private key s_x , the key
20 distribution data K_{Cx} received from the sender A , and stores the encrypted common key K_A in the memory 314.

(iii) The sender A uses the random number generator 211 within the sender side apparatus 200 to generate a suitable integer r , and sends it to the receiver x by means of the communication unit 215. Further, the sender A
25 uses the arithmetic unit 212 to calculate the data encryption-decryption key $DK=F(r, K_A)$ from the integer r and the common key K_A using a suitable

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function F , and stores it in the memory 214.

- (iv) The receiver x uses the arithmetic unit 312 within the receiver side apparatus 300 to calculate the data encryption-decryption key $DK=F(r, K_A)$ from the common key K_A and the integer r received from the sender, using the function F , and stores it in the memory 314.

After the distribution of the data encryption-decryption key DK by the key distribution process, an encryption-decryption process is performed as described in the following.

(3) Encryption-decryption process

- (i) The sender A uses the encryption-decryption unit 213 within the sender side apparatus 200 to encrypt data to be delivered P with the data encryption-decryption key DK , and sends it to the receiver x by means of the communication unit 213.
- (ii) The receiver x receives the encrypted delivered data P through the communication unit 314 within the receiver side apparatus 300. Then, the receiver x decrypts the data P with the data encryption-decryption key DK using the encryption-decryption unit 312.
- (iii) The sender A changes the value of the common key K_A periodically so as to change the data encryption-decryption key DK .

According to the above-described seventh embodiment of the present invention also, the receiver can have respective common keys of a plurality of senders in common with them, using a single piece of secret key.

Next, an eighth embodiment of the present invention will be described.

The eighth embodiment of the present invention is similar to the above-described sixth and seventh embodiments except that sending of the

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key distribution data K_{Cx} from the sender to the receiver is modified as described in the following.

Namely, in the present eighth embodiment, the sender A separates the set of all of the receivers into some subsets. Next, in sending key
5 distribution data K_{Ci} ($1 \leq i \leq n$, n : the number of the receivers) to the receivers, the sender assigns communication channels to the above-described subsets of the receivers respectively, and sends the key distribution data K_{Ci} to the receivers through the communication unit 215 within the sender side apparatus 200.

10 On the other hand, any receiver x receives the key distribution data K_{Cx} by the communication unit 314 within the receiver side apparatus 300 through the communication channel assigned to the subset to which he belongs.

The embodiments of the present invention have been described
15 above.

Each of the above-described processes performed by the trusted center side apparatus 100, the sender side apparatus 200, and the receiver side apparatus 300 may be performed by a computer which executes programs describing procedures for performing respective processes. In that
20 case, each program describing a procedure for performing a process may be stored in a storing medium to be supplied to each computer.

As described above, according to the present invention, a receiver can receive distribution of respective common keys of a plurality of senders, using only one receiver secret key.

25 Further, in some inventions disclosed herein, it is possible to have a common key for data encryption and decryption in common among any

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sender and receivers belonging to any set of receivers. Further, even if the number of the receivers is large, it is possible to realize common key distribution using short broadcast communication data for distribution of a common key.

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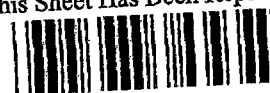


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What is claimed is:

- 1 1. A method of key distribution having an encryption-decryption
2 key of secure broadcast communication in common among a plurality of
3 senders and a plurality of receivers in a communication system
4 comprising apparatuses of said plurality of senders and apparatuses of
5 said plurality of receivers, wherein:
 - 6 in an apparatus of a trusted center provided separately from
7 said senders and said receivers, first key information which is common
8 to said plurality of senders and is used for having a key in common
9 among said senders and said receivers, is produced and distributed to
10 prescribed receivers, while second key information which is used for a
11 key in common among said senders and said receivers is produced and
12 distributed to said senders;
 - 13 in an apparatus of a sender, third key information, which is
14 used by said receivers together with said first key information to
15 calculate an encryption-decryption key used by the sender in question
16 in broadcast communication, is produced using the second key
17 information distributed from said trusted center and is sent to said
18 receivers who have the encryption-decryption key in common with said
19 sender; and
 - 20 in an apparatus of a receiver, the encryption-decryption key
21 which is used for broadcast communication by the sender who
22 distributed said third key information is obtained using the first key
23 information distributed from said trusted center and the third key
24 information distributed from said sender.

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1 2. A method of key distribution having an encryption-decryption
2 key of secure broadcast communication in common among a plurality of
3 senders and a plurality of receivers in a communication system
4 comprising an apparatus of a trusted center, apparatuses of said
5 plurality of senders and apparatuses of said plurality of receivers,
6 comprising:
7 a first step, in the apparatus of the trusted center, of producing
8 receiver secret key s_x of receiver x so as to correspond to a subset S' of a
9 finite set S , and distributing the receiver secret key s_x to said receiver x ,
10 while producing sender registration key t_A of a sender A ;
11 a second step, in the apparatus of the sender A , of producing
12 sender secret key g_A of the sender A and calculating an encryption-
13 decryption key K_A from the sender secret key g_A ;
14 a third step, in cooperation between the apparatus of the
15 trusted center and the apparatus of the sender A , of producing key
16 distribution data W by a function of at least the sender secret key g_A and
17 the sender registration key t_A , while producing receiver registration key
18 r_x of the receiver x for the sender A by a function of at least the receiver
19 private key information s_x and the sender registration key t_A ;
20 a fourth step, in the apparatus of the sender A , of sending the
21 receiver registration key r_x to the receiver x and broadcasting the key
22 distribution data W to each receiver; and
23 a fifth step, in the apparatus of the receiver x , of calculating
24 the encryption-decryption key K_A from the receiver private key
25 information s_x distributed from the trusted center and the key
26 distribution data W and the receiver registration key r_x received from
27 the sender A , the encryption-decryption key K_A being to be used for
28 secure broadcast communication performed by the sender A ;

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29 wherein said receiver secret key s_x , the key distribution data W ,
30 and the receiver registration key r_x are produced in such a relation that
31 the encryption-decryption key K_A can be derived from said receiver
32 secret key s_x , the key distribution data W , and the receiver registration
33 key r_x , and other information opened to each receiver.

1 3. The method of key distribution according to Claim 2, wherein:
2 said second step is one in which, in the apparatus of the sender
3 A, variable information r_A and r_A' for changing the encryption-
4 decryption key are produced, and the encryption-decryption key K_A is
5 calculated by a function of the variable information r_A and r_A' acting on
6 the secret key g_A , the variable information r_A being for changing the
7 encryption-decryption key periodically, and the variable information r_A'
8 being for changing the encryption-decryption key for every broadcast
9 content; and

10 said third step is one in which:
11 in the apparatus of the trusted center, data W' is produced by a
12 function of the sender secret key g_A and the sender registration key t_A ,
13 and sent to the sender A, while data r_x' is produced by a function of the
14 receiver secret key s_x of the receiver x and the sender registration key t_A ,
15 and sent to the sender A; and

16 in the apparatus of the sender A, the key distribution data W is
17 produced by a function of the variable information r_A acting on the data
18 W' , and the receiver registration key r_x of the receiver x for the sender A
19 is produced by a function of the variable information r_A' acting on the
20 data r_x' .

1 4. A method of key distribution for having an encryption-

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2 decryption key of secure broadcast communication in common among a
3 plurality of senders and a plurality of receivers in a communication
4 system comprising an apparatus of a trusted center, apparatuses of said
5 plurality of senders and apparatuses of said plurality of receivers,
6 comprising:

7 a first step for distributing receiver secret key σ_x , $s_x(\sigma_x)$ to the
8 receiver x, in which, in the apparatus of the trusted center,

$$9 \quad e_i \in Z \quad (1 \leq i \leq m)$$

10 are produced as a trusted center secret key;

$$11 \quad t_A \in Z$$

12 is produced as a sender registration key of a sender A; $\sigma_x \in Skm$ and

$$13 \quad s_x(\sigma_x) \in Z$$

14 are produced (here, when $\sigma, \sigma' \in Skm$ for S_{km} :

$$15 \quad S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \rightarrow B = \{1, 2, \Lambda, m\}, 0 < k < m \}$$

16 it is expressed by the following expression:

$$17 \quad \sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A)$$

18 “ $\sim$ ” being an equivalence relation on S_{km} , to lead to the following
19 expression:

$$20 \quad \tilde{S}_{km} = S_{km} / \sim$$

21), a second step, in which, in an apparatus of the sender A, are
22 produced sender secret keys g_A , L_A , and a finite Abelian group G_A
23 satisfying:

$$24 \quad g_A \in G_A \\ L_A = \text{ord}_{G_A}(g_A)$$

25 (here,

$$26 \quad \text{ord}_{G_A}(g)$$

27 means a least positive integer a which satisfies:



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28 $g^a = 1 (\in G_A)$
 29), and the sender secret key g_A is sent to the trusted center;
 30 a third step, in which, in the apparatus of the trusted center,
 31 receiver registration data defined by the following expression:

$$32 \quad s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

33 is calculated from the trusted center secret key e_i , the sender
 34 registration key t_A , and the receiver secret key $s_x(\sigma_x)$, σ_x ; and the
 35 receiver registration data $s_x(\sigma_x, A)$ is sent to the sender A, while key
 36 distribution data defined as:

$$37 \quad y_{Ai} = g_A^{u_i} (\in G_A) \quad (1 \leq i \leq m)$$

38 is calculated from the g_A received from the sender A, the trusted center
 39 secret key e_i , and the sender registration key t_A ; and the key
 40 distribution data y_{Ai} is sent to the sender A;

41 a fourth step, in which, in the apparatus of the sender A,
 42 integers r, r' are produced; a receiver registration key $r_x(\sigma_x, A)$
 43 satisfying:

$$44 \quad r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

45 is calculated from the receiver registration data $s_x(\sigma_x, A)$ received from
 46 the trusted center, the sender secret key L_A , and the integer r' ; and the
 47 receiver registration key $r_x(\sigma_x, A)$ is sent to the receiver x; while key
 48 distribution data defined as:

$$49 \quad z_{Ai} = y_{Ai}^r (\in G_A) \quad (1 \leq i \leq m)$$

50 is calculated from the key distribution data y_{Ai} received from the
 51 trusted center and the integer r ; and the key distribution data z_{Ai} is
 52 broadcasted to each receiver; and

53 a fifth step, in which, in an apparatus of the receiver x, an
 54 encryption-decryption key K_A is calculated by:





55
$$K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A) s_x(\sigma_x)} (\in G_A)$$

56 from the receiver registration key $r_x(\sigma_x, A)$ received from the sender, the
 57 key distribution data z_{Ai} , the receiver secret key σ_x , $s_x(\sigma_x)$ distributed
 58 from the trusted center; the encryption-decryption key K_A being
 59 generated by the sender A employing:

60
$$K_A = g_A^{r'} (\in G_A)$$

61 to be used for broadcast communication.

1 5. A method of key distribution having an encryption-decryption
 2 key of secure broadcast communication in common among a plurality of
 3 senders and a plurality of receivers in a communication system
 4 comprising an apparatus of a trusted center, apparatuses of said
 5 plurality of senders and apparatuses of said plurality of receivers,
 6 comprising:

7 a first step, in which, in the apparatus of the trusted center,

8
$$e_i \in Z \quad (1 \leq i \leq m)$$

9 is produced as a trusted center secret key;

10
$$t_A \in Z$$

11 is produced as a sender registration key of a sender A; $\sigma_x \in S_{km}$ and

12
$$s_x(\sigma_x) \in Z$$

13 are produced as receiver secret keys of a receiver x;

14 (here, when $\sigma, \sigma' \in S_{km}$ for a set defined as:

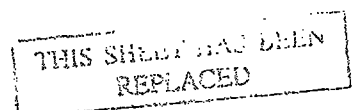
15
$$S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m \}$$

16 it is expressed by the following expression:

17
$$\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A)$$

18 and, in that case, " $\sim$ " becomes an equivalence relation on S_{km} , which

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19 leads to the following expression:

20
$$\tilde{S}_{km} = S_{km} / \sim$$

21); and the receiver secret keys σ_x , $s_x(\sigma_x)$ are sent to the receiver x;

22 a second step, in which, in an apparatus of the sender A,

23
$$\begin{aligned} P_A, Q_A: & \text{ prime numbers} \\ L_A &= \text{lcm}(\text{ord}_{P_A}(g_A), \text{ord}_{Q_A}(g_A)) \\ g_A &\in \mathbb{Z}, \quad 0 < g_A < N_A \\ r, r' &\in \mathbb{Z}, \quad 0 < r, r' < L_A \end{aligned}$$

24 are produced as sender secret keys;

25
$$N_A (= P_A Q_A)$$

26 is produced as a sender public key; and the sender secret key g_A is sent
27 to the trusted center;

28 a fourth step, in which, in the apparatus of the trusted center,
29 receiver registration data defined as:

30
$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

31 is calculated from the trusted center secret key e_i , the sender
32 registration key t_A , and the receiver secret keys $s_x(\sigma_x)$, σ_x ; and the
33 receiver registration data $s_x(\sigma_x, A)$ is sent to the sender A; while key
34 distribution data defined as:

35
$$y_{Ai} = g_A^{u_i} \bmod N_A \quad (1 \leq i \leq m)$$

36 is calculated from g_A received from the sender A, the trusted center
37 secret key e_i , and the sender registration key t_A ; and the key
38 distribution data y_{Ai} is sent to the sender A;

39 a fifth step, in which, in the apparatus of the sender A, integers
40 r, r' are produced; and a receiver registration key $r_x(\sigma_x, A)$ satisfying:

41
$$r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

42 is calculated from the receiver registration data $s_x(\sigma_x, A)$ received from
43 the trusted center, the sender secret key L_A , and the integer r' ; and the





44 receiver registration key $r_x(\sigma_x, A)$ is sent to the receiver x ; while key
45 distribution data defined as:

$$46 \quad z_{Ai} = y_{Ai}^r \bmod N_A \quad (1 \leq i \leq m)$$

47 is calculated from the key distribution data y_{Ai} received from the
48 trusted center and the integer r ; and the key distribution data z_{Ai} is
49 broadcasted to each receiver; and

50 a sixth step, in which, in an apparatus of the receiver x , an
51 encryption-decryption key K_A is calculated by:

$$52 \quad K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A) s_x(\sigma_x)} \bmod N_A$$

53 from the receiver registration key $r_x(\sigma_x, A)$ received from the sender, the
54 key distribution data z_{Ai} , and the receiver secret keys $\sigma_x, s_x(\sigma_x)$
55 distributed from the trusted center; the encryption-decryption key
56 being generated by the sender A using:

$$57 \quad K_A = g_A^{r'} \bmod N_A$$

58 to be used for broadcast communication.

1 6. The method of key distribution according to Claim 5, further
2 comprising:

3 a seventh step, in which, in said apparatus of the sender A ,
4 said secret key L_A is generated and sent to the trusted center before the
5 apparatus of the trusted center performs the first step of generating the
6 trusted center secret key e_i , the sender identification t_A , and the
7 receiver secret key $s_x(\sigma_x)$;

8 wherein:

9 in said first step, in the apparatus of the trusted center, the
10 trusted center information e_i , the sender identification t_A , and the
11 receiver secret key $s_x(\sigma_x)$ are produced from the sender secret key L_A

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12 received from the sender so as to satisfy:

$$\begin{aligned} e_i &\in Z, \quad 0 < e_i < L_A, \quad (1 \leq i \leq m) \\ t_A &\in Z, \quad 0 < t_A < L_A \\ s_x(\sigma_x) &\in Z, \quad 0 < s_x(\sigma_x) < L_A, \end{aligned}$$

14 and,

15 in said fourth step, in the apparatus of the trusted center, the
16 receiver registration data $s_x(\sigma_x, A)$ is calculated from the sender secret
17 key L_A , the trusted center secret key e_i , the sender registration key t_A ,
18 and the receiver secret key $s_x(\sigma_x)$ so as to satisfy:

$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)} \mod L_A.$$

1 7. The method of key distribution according to Claim 5, further
2 comprising:

3 an eighth step, in which, in the apparatus of the sender A, a
4 value of the integer r is changed; and a value of the encryption-
5 decryption key K_A is recalculated in accordance with the changed
6 integer r ; while the key distribution data z_{Ai} is recalculated in
7 accordance with the changed integer r and broadcasted to each receiver
8 again; and

9 a ninth step, in which, in the apparatus of the receiver x , the
10 encryption-decryption key K_A is recalculated in accordance with the key
11 distribution data z_{Ai} broadcasted again.

1 8. The method of key distribution according to Claim 5, wherein:
2 in said fifth step, in the apparatus of the sender A, a value of
3 the integer r' is generated in accordance with each piece of information
4 or each kind of information to be encrypted and broadcasted; and





5 encryption-decryption keys K_A are calculated for respective information
6 or respective kinds of information in accordance with respective values
7 of the integer r' generated in accordance with the information or kinds
8 of information; while pieces of receiver registration key $r_x(\sigma_x, A)$ are
9 calculated in accordance with respective integers r' generated in
10 accordance with respective said information or kinds of information on
11 which the receiver x is allowed to perform decryption, and the said
12 pieces of receiver registration key $r_x(\sigma_x, A)$ are sent to the receiver x ; and
13 in said sixth step, in the apparatus of the receiver x ,
14 encryption-decryption key K_A for said information or kinds of
15 information on which said receiver x is allowed to perform decryption
16 are calculated in accordance with respective pieces of receiver
17 registration key $r_x(s_x, A)$ received.

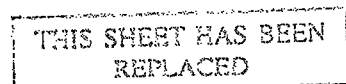
1 9. The method of key distribution according to Claim 5, further
2 comprising:

3 a tenth step in which, in the apparatuses of the sender A and
4 the receiver x , possession of the receiver secret key $s_x(\sigma_x)$ by said
5 receiver x is authenticated according to a system by which the receiver
6 secret key $s_x(\sigma_x)$ is made secret from the sender A , before the fifth step
7 for sending the receiver registration key $r_x(\sigma_x, A)$ from the sender A to
8 the receiver x .

1 10. The method of key distribution according to Claim 5, further
2 comprising:

3 an eleventh step in which, in the apparatus of the sender A , a
4 fee for information or a kind of information, which is encrypted with the
5 encryption-decryption key K_A corresponding to the receiver registration

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6 key $r_x(\sigma_x, A)$ and broadcasted, is calculated as accounting information
7 for the receiver A, said eleventh step being performed when the fifth
8 step is performed for sending the receiver registration key $r_x(\sigma_x, A)$ from
9 the sender A to the receiver x.

1 11. A method of key distribution having an encryption-decryption
2 key of secure broadcast communication in common among a plurality of
3 senders and a plurality of receivers in a communication system
4 comprising apparatuses of said plurality of senders and apparatuses of
5 said plurality of receivers, comprising:

6 a first step, in which, in an apparatus of a trusted center
7 provided separately from said apparatuses of senders and said receivers,
8 a secret key K0 is generated and distributed to each sender and each
9 receiver;

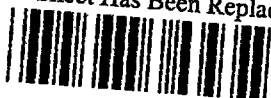
10 a second step, in which, in an apparatus of a sender A, a sender
11 registration key BIDA is generated and sent to the receivers, while a
12 common key $K_A' = F(K0, BIDA)$ is calculated from the secret key K0
13 distributed from the trusted center and the sender registration key
14 BIDA using a given one-way function F;

15 a third step, in which, in the apparatuses of the receivers, the
16 common key K_A' is calculated from the secret key K0 distributed from
17 the trusted center and the sender registration key BIDA received from
18 the sender A using said one-way function F;

19 a fourth step, in which, in the apparatus of the sender A, an
20 integer r is produced and sent to the receivers; and

21 a fifth step, in which, in the apparatuses of the receivers, an
22 encryption-decryption key K_A is calculated by $K_A = F'(r, K_A')$ from the
23 integer r received from the sender A and the common key K_A' using a

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24 given function F' , the encryption-decryption key K_A being calculated by
25 the sender A using $K_A = F'(r, K_A')$.

1 12. A method of key distribution having an encryption-decryption
2 key of secure broadcast communication in common among a plurality of
3 senders and a plurality of receivers in a communication system
4 comprising apparatuses of said plurality of senders and apparatuses of
5 said plurality of receivers, comprising:

6 a first step, in which, in an apparatus of a trusted center
7 provided separately from said apparatuses of senders and said receivers,
8 a personal key s_x for a receiver x is produced, and distributed to the
9 sender A and the receiver x ;

10 a second step, in which, in an apparatus of the sender A, a
11 common key K_A is produced, and said common key K_A is owned in
12 common with the receiver x by cipher communication using the personal
13 key s_x for the receiver x , which has been distributed from the trusted
14 center;

15 a third step, in which, in the apparatus of the sender A,
16 information r is produced and sent to the receiver x ; and

17 a fourth step, in which, in each of the apparatuses of the
18 receiver x and the sender A, an encryption-decryption key DK is
19 calculated in accordance with an expression, $DK = F(r, K_A)$.

1 13. The method of key distribution according to Claim 12, further
2 comprising:

3 a fifth step, in which, in the apparatus of the sender A, the
4 common key K_A is changed, and the changed common key K_A is owned in
5 common with the receiver x by cipher communication using the personal

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6 key s_x for the receiver x ; and

7 a sixth step, in which, in each of the apparatuses of the sender
8 A and the receiver x , the encryption-decryption key DK is recalculated
9 by $DK=F(r, K_A)$ in accordance with the changed common key K_A .

1 14. The method of key distribution according to Claim 12, further
2 comprising:

3 a seventh step, in which, in the apparatus of the sender A, the
4 information r is changed, and the changed information r is informed to
5 the receivers; and

6 an eighth step, in which, in each of the apparatuses of the
7 sender A and the receivers, the encryption-decryption key DK is
8 recalculated by $DK=F(r, K_A)$ in accordance with the changed
9 information r .

1 15. The method of key distribution according to Claim 12, further
2 comprising:

3 a ninth step, in which, in the apparatus of the sender A, key
4 information for the receiver x to calculate said encryption-decryption
5 key DK is sent through a given communication channel which is
6 different from a communication channel for broadcasting information
7 encrypted with the encryption-decryption key DK .

1 16. A key distribution system for having an encryption-decryption
2 key used for secure broadcast communication in common among senders
3 and receivers, wherein:

4 said system comprises a trusted center apparatus, a plurality
5 of sender apparatuses, and a plurality of receiver apparatuses;

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6 said trusted center apparatus comprises:

7 means for producing:

$$8 \quad e_i \in Z \quad (1 \leq i \leq m)$$

9 as a trusted center secret key;

$$10 \quad t_A \in Z$$

11 as a sender registration key of a sender A; and $\sigma_x \in S_{km}$ and

$$12 \quad s_x(\sigma_x) \in Z$$

13 (here, when $\sigma, \sigma' \in S_{km}$ for a set S_{km} defined as:

$$14 \quad S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m \}$$

15 it is expressed by the following expression:

$$16 \quad \sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A)$$

17 and, in that case, " $\sim$ " becomes an equivalence relation on S_{km} , leading
18 to the following definition:

$$19 \quad \tilde{S}_{km} = S_{km} / \sim$$

20); and for distributing the receiver secret keys $\sigma_x, s_x(\sigma_x)$ to said receiver
21 apparatus used by the receiver x;

22 means for calculating receiver registration data defined by:

$$23 \quad s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

24 from the trusted center secret key e_i , the sender registration key t_A , and
25 the receiver secret keys $s_x(\sigma_x)$, σ_x , and for sending the receiver
26 registration data $s_x(\sigma_x, A)$ to said sender apparatus used by the sender
27 A; and

28 means for calculating key distribution data defined by:

$$29 \quad y_{Ai} = g_A^{t_A e_i} (\in G_A) \quad (1 \leq i \leq m)$$

30 from g_A received from said sender apparatus used by the sender A, the
31 trusted center secret key e_i , and the sender registration key t_A , and for





32 sending the key distribution data y_{Ai} to said sender apparatus used by
33 the sender A;

34 said sender apparatus used by the sender A comprises:

35 means for producing sender secret keys g_A , L_A , and a finite
36 Abelian group G_A , satisfying:

$$\begin{aligned} 37 \quad & g_A \in G_A \\ & L_A = \text{ord}_{G_A}(g_A) \end{aligned}$$

38 (here,

$$39 \quad \text{ord}_{G_A}(g))$$

40 means a least positive integer a satisfying:

$$41 \quad g^a = 1 (\in G_A)$$

42), and for sending the sender secret key g_A to said trusted center
43 apparatus;

44 means for producing integers r , r' , calculating receiver
45 registration key $r_x(\sigma_x, A)$ satisfying:

$$46 \quad r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

47 from the receiver registration data $s_x(\sigma_x, A)$ received from said trusted
48 center apparatus, the sender secret key L_A , and the integer r' ; and for
49 sending the receiver registration key $r_x(\sigma_x, A)$ to said receiver apparatus
50 of the receiver x ; and

51 means for calculating key distribution data defined as:

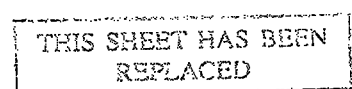
$$52 \quad z_{Ai} = y_{Ai}^r (\in G_A) \quad (1 \leq i \leq m)$$

53 from the key distribution data y_{Ai} received from said trusted center
54 apparatus, and the integer r ; and for broadcasting the key distribution
55 data z_{Ai} to said plurality of receiver apparatuses; and

56 said receiver apparatus of the receiver x comprises:

57 means for calculating encryption-decryption key K_A by the
58 following expression:

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59
$$K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A) s_x(\sigma_x)} (\in G_A),$$

60 said encryption-decryption key K_A being generated by said sender
61 apparatus of the sender A by the following expression:

62
$$K_A = g_A^{r'} (\in G_A).$$

1 17. A trusted center apparatus which supports distribution of an
2 encryption-decryption key from a sender apparatus to a receiver
3 apparatus, having said encryption-decryption key used in secure
4 broadcast communication in common among senders and receivers,
5 comprising:

6 means for producing:

7
$$e_i \in Z \quad (1 \leq i \leq m)$$

8 as a trusted center secret key;

9
$$t_A \in Z$$

10 as a sender registration key of a sender A; and, as receiver secret keys of
11 a receiver x, $\sigma_x \in S_{km}$ and:

12
$$s_x(\sigma_x) \in Z$$

13 (here, when $\sigma, \sigma' \in S_{km}$ for a set S_{km} :

14

15
$$S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m \},$$

16 it is expressed by the following expression:

17
$$\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A),$$

18 with " $\sim$ " being an equivalence relation on S_{km} , which lead to an
19 expression:

20

21
$$\tilde{S}_{km} = S_{km} / \sim,$$





22 and for distributing the receiver secret keys σ_x , $s_x(\sigma_x)$ to said receiver
 23 apparatus used by the receiver x;

24 means for calculating receiver registration data defined by:

$$25 \quad s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

26 from the trusted center secret key e_i , the sender registration key t_A , and
 27 the receiver secret keys $s_x(\sigma_x)$, σ_x ; and for sending the receiver
 28 registration data $s_x(\sigma_x, A)$ to said sender apparatus used by the sender
 29 A; and

30 means for calculating key distribution data defined as:

$$31 \quad y_{Ai} = g_A^{u_i} (\in G_A) \quad (1 \leq i \leq m)$$

32 from g_A received from said sender apparatus of the sender A, the trusted
 33 center secret key e_i , and the sender registration key t_A ; and for sending
 34 the key distribution data y_{Ai} to said sender apparatus of the sender A.

1 18. A sender apparatus which distributes an encryption-decryption
 2 key to a receiver apparatus being supported by a trusted center
 3 apparatus, having said encryption-decryption key used in secure
 4 broadcast communication in common among senders and receivers,
 5 comprising:

6 means for producing sender secret key g_A , L_A , and a finite
 7 Abelian group G_A satisfying:

$$8 \quad \begin{aligned} g_A &\in G_A \\ L_A &= \text{ord}_{G_A}(g_A) \end{aligned}$$

9 (here,

$$10 \quad \text{ord}_{G_A}(g))$$

11 means a least positive integer a satisfying:

$$12 \quad g^a = 1 (\in G_A)$$

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13); and for sending the sender secret key g_A to said trusted center
 14 apparatus;
 15 means for producing integers r, r' ; calculating a receiver
 16 registration key $r_x(\sigma_x, A)$ satisfying:
 17
$$r_x(\sigma_x, A)s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

 18 from the receiver registration data $s_x(\sigma_x, A)$, the sender secret key L_A ,
 19 and the integer r , and for sending the receiver registration key $r_x(\sigma_x, A)$
 20 to said receiver apparatus used by the receiver x ;
 21 means for calculating key distribution data defined as:
 22
$$z_{Ai} = y_{Ai}^r \pmod{G_A} \quad (1 \leq i \leq m)$$

 23 from key distribution data y_{Ai} received from said trusted center
 24 apparatus, and the integer r ; and for broadcasting the key distribution
 25 data z_{Ai} to a plurality of said receiver apparatuses as objects of
 26 broadcast communication; and
 27 means for calculating an encryption-decryption key K_A used for
 28 broadcast communication, in accordance with:
 29
$$K_A = g_A^{r'} \pmod{G_A}.$$

1 19. A receiver apparatus which receives an encryption-decryption
 2 key from a sender apparatus being supported by a trusted center
 3 apparatus, having said encryption-decryption key used in secure
 4 broadcast communication in common among senders and receivers,
 5 comprising:

6 means for calculating the encryption-decryption key K_A used
 7 for the broadcast communication by an expression:

$$8 \quad K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A)s_x(\sigma_x)} \pmod{G_A}$$

9 from a receiver registration data $r_x(\sigma_x, A)$ and a key distribution data z_{Ai}

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10 received from the sender apparatus used by a sender A, and receiver
 11 secret keys σ_x , $s_x(\sigma_x)$ received from the trusted center apparatus, the
 12 encryption-decryption key K_A being generated by said sender apparatus
 13 of the sender A using an expression:

14
$$K_A = g_A^{r'} (\in G_A).$$

1 20. A receiver apparatus which receives an encryption-decryption
 2 key from a sender apparatus being supported by a trusted center
 3 apparatus, having said encryption-decryption key used in secure
 4 broadcast communication in common among senders and receivers,
 5 comprising:

6 a main apparatus and an auxiliary apparatus;

7 wherein:

8 said auxiliary apparatus comprises:

9 means for connection with said trusted center apparatus;

10 storing means for taking in and storing receiver secret keys

11 $s_x(\sigma_x)$, σ_x produced by said trusted center apparatus;

12 means for connection with said main apparatus, and

13 means for taking in key distribution data z_{Ai} received by said

14 main apparatus from said sender apparatus used by a sender A;

15 calculating:

16
$$\xi_x(\sigma_x, A) = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{s_x(\sigma_x)} \bmod N_A$$

17 from the key distribution data z_{Ai} and the receiver secret keys σ_x , $s_x(\sigma_x)$
 18 stored in said storing means; and for outputting a result $\xi_x(\sigma_x, A)$ of the
 19 calculation to said main apparatus; and

20 said main apparatus comprises:

21 means for outputting the key distribution data z_{Ai} received





22 from said sender apparatus of the sender A to said auxiliary apparatus;
 23 and

24 means for calculating encryption-decryption key K_A used in
 25 broadcast communication, using an expression:

$$26 \quad K_A = \xi_x(\sigma_x, A)^{r_x(\sigma_x, A)} \bmod N_A$$

27 from a receiver registration key $r_x(\sigma_x, A)$ received from the sender
 28 apparatus of the sender A, and $\xi_x(\sigma_x, A)$ output from the auxiliary
 29 apparatus, the encryption-decryption key K_A being generated by said
 30 sender apparatus of the sender A, using an expression:

$$31 \quad K_A = g_A^{r'} (\in G_A).$$

1 21. An auxiliary apparatus used being connected with a trusted
 2 center apparatus and a plurality of receiver apparatuses, in a
 3 communication system comprising a plurality of sender apparatuses
 4 which perform broadcast communication, said plurality of receiver
 5 apparatuses, and said trusted center apparatus which supports
 6 distribution of an encryption-decryption key of a sender apparatus used
 7 in secure broadcast communication to a receiver apparatus; said
 8 auxiliary apparatus comprising:

9 means for connection with said trusted center apparatus;

10 storing means for taking in and storing receiver secret keys

11 $s_x(\sigma_x)$, σ_x for a receiver x, produced by said trusted center apparatus;

12 means for connection with said receiver apparatus used by the
 13 receiver x; and

14 means for taking in key distribution data z_{Ai} received by said
 15 receiver apparatus used by the receiver x from said sender apparatus
 16 used by the sender A; for calculating:

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17
$$\xi_x(\sigma_x, A) = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{s_x(\sigma_x)} \bmod N_A$$

18 from the key distribution data z_{Ai} and the receiver secret keys $\sigma_x, s_x(\sigma_x)$
 19 stored in said storing means; and for outputting a calculation result
 20 $\xi_x(\sigma_x, A)$ to said receiver apparatus of the receiver x.

1 22. A storing medium which stores a program for making an
 2 information processing apparatus execute a process for supporting
 3 distribution of an encryption-decryption key from sender apparatuses to
 4 receiver apparatuses, having said encryption-decryption key used in
 5 secure broadcast communication in common among senders and
 6 receivers, wherein:

7 said program makes the information processing apparatus
 8 execute:

9 a step of producing:

10
$$e_i \in Z \quad (1 \leq i \leq m)$$

11 as a trusted center secret key,

12
$$t_A \in Z$$

13 as a sender registration key of a sender A, and, as receiver secret keys of
 14 a receiver x, $\sigma_x \in S_{km}$ and:

15
$$s_x(\sigma_x) \in Z$$

16 (here, when $\sigma, \sigma' \in S_{km}$ for a set S_{km} defined as:

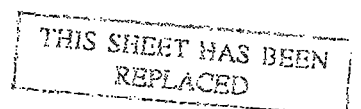
17
$$S_{km} = \{ \sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, \quad 0 < k < m \}$$

18 it is expressed by the following expression:

19
$$\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A),$$

20 with " $\sim$ " becoming an equivalence relation on S_{km} , leading to an
 21 expression:

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22 $\tilde{S}_{km} = S_{km} / \sim$

23); and of distributing the receiver secret keys σ_x , $s_x(\sigma_x)$ to said receiver
24 apparatus used by the receiver x;

25 a step of calculating receiver registration data defined as:

26
$$s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$$

27 from the trusted center secret key e_i , the sender registration key t_A , and
28 the receiver secret keys $s_x(\sigma_x)$, σ_x ; and of sending the receiver
29 registration data $s_x(\sigma_x, A)$ to said sender apparatus used by the sender
30 A; and

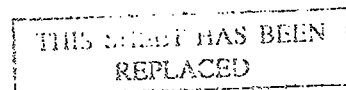
31 a step of calculating key distribution data y_{Ai} defined as:

32
$$y_{Ai} = g_A^{u_i} (\in G_A) \quad (1 \leq i \leq m)$$

33 from g_A received from said sender apparatus used by the sender A, the
34 trusted center secret key e_i , and the sender registration key t_A ; and of
35 sending the key distribution data y_{Ai} to said sender apparatus used by
36 the sender A.

1 23. A storing medium which stores a program for making an
2 information processing apparatus execute a process in sender
3 apparatuses which distribute an encryption-decryption key to receiver
4 apparatuses, being supported by a trusted center apparatus, having
5 said encryption-decryption key used in secure broadcast communication
6 in common among senders and receivers, wherein:
7 said program makes the information processing apparatus
8 execute:
9 a step of producing sender secret keys g_A , L_A , and a finite
10 Abelian group G_A satisfying:

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$$11 \quad g_A \in G_A$$

$$L_A = \text{ord}_{G_A}(g_A)$$



12 (here,

$$13 \quad \text{ord}_{G_A}(g_A)$$

14 means a least positive integer a satisfying:

$$15 \quad g_A^a = 1 (\in G_A)$$

16), and of sending the sender secret key g_A to said trusted center

17 apparatus;

18 a step of producing integers r, r' ; calculating receiver

19 registration key $r_x(\sigma_x, A)$ satisfying:

$$20 \quad r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

21 and sending the receiver registration key $r_x(\sigma_x, A)$ to said receiver

22 apparatus used by the receiver x ;

23 a step of calculating key distribution data defined as:

$$24 \quad z_{Ai} = y_{Ai}^{r'} (\in G_A) \quad (1 \leq i \leq m)$$

25 from key distribution data y_{Ai} received from said key administration

26 apparatus, and the integer r ; and broadcasting the key distribution data

27 z_{Ai} to said plurality of receiver apparatuses as objects of broadcast

28 communication; and

29 a step of calculating an encryption-decryption key K_A used in

30 broadcast communication, using:

$$31 \quad K_A = g_A^{rr'} (\in G_A).$$

- 1 24. A storing medium which stores a program for making an
- 2 information processing apparatus execute a process in receiver
- 3 apparatuses which receive an encryption-decryption key distributed
- 4 from sender apparatuses, being supported by a trusted center apparatus,
- 5 having said encryption-decryption key used in secure broadcast

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6 communication in common among senders and receivers, wherein:
 7 said program makes the information processing apparatus
 8 execute:
 9 a step of calculating an encryption-decryption key K_A used for
 10 broadcast communication, using an expression:

$$11 \quad K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A)s_x(\sigma_x)} (\in G_A)$$

12 from a receiver registration key $r_x(\sigma_x, A)$ received from said sender
 13 apparatus used by a sender A, key distribution data z_{Ai} , and a receiver
 14 secret key σ_x , $s_x(\sigma_x)$ distributed from the trusted center, the
 15 encryption-decryption key K_A being generated by said sender apparatus
 16 used by the sender A, using:

$$17 \quad K_A = g_A^{r'} (\in G_A).$$

1 25. A key distribution system for having an encryption-decryption
 2 key used for secure broadcast communication in common among senders
 3 and receivers, wherein:

4 said system comprises a trusted center apparatus, a plurality
 5 of sender apparatuses, and a plurality of receiver apparatuses;

6 said trusted center apparatus comprises:
 7 means for producing:

$$8 \quad e_i \in Z \quad (1 \leq i \leq m)$$

9 as a trusted center secret key,

$$10 \quad t_A \in Z$$

11 as a sender registration key of a sender A, and, as receiver secret keys of
 12 a receiver x, $\sigma_x \in S_{km}$ and

$$13 \quad s_x(\sigma_x) \in Z$$

14 (here, when $\sigma, \sigma' \in S_{km}$ for a set S_{km} defined as:

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15 $S_{km} = \{\sigma \mid \text{one to one map } \sigma : A = \{1, 2, \Lambda, k\} \\ \rightarrow B = \{1, 2, \Lambda, m\}, 0 < k < m\}$,

16 it is expressed by:

17 $\sigma \sim \sigma' \stackrel{\text{def}}{\Leftrightarrow} \sigma(A) = \sigma'(A)$

18 with “ $\sim$ ” becoming an equivalence relation on S_{km} , leading to the
19 following expression:

20 $\tilde{S}_{km} = S_{km} / \sim$

21), and of distributing the receiver secret key σ_x , $s_x(\sigma_x)$ to said receiver
22 apparatus used by a receiver x;

23 means for calculating:

24 $s_x(\sigma_x, A) = t_A s_x(\sigma_x) \sum_{i=1}^k e_{\sigma_x(i)}$

25 from the key distributor secret key e_i , the sender registration key t_A ,
26 and the receiver secret keys $s_x(\sigma_x)$, σ_x , and sending the receiver
27 registration data $s_x(\sigma_x, A)$ to said sender apparatus used by the sender
28 A; and

29 means for calculating key distribution data y_{Ai} :

30 $y_{Ai} = g_A^{u_i^q} \bmod N_A \quad (1 \leq i \leq m)$

31 from g_A received from the sender A, the trusted center secret key e_i , and
32 the sender registration key t_A , and sending the key distribution data y_{Ai}
33 to said sender apparatus used by the sender A;

34 said sender apparatus used by the sender A comprises:

35 means for producing:

36 $P_A, Q_A : \text{prime numbers}$
 $L_A = \text{lcm}(\text{ord}_{P_A}(g_A), \text{ord}_{Q_A}(g_A))$
 $g_A \in \mathbb{Z}, 0 < g_A < N_A$
 $r, r' \in \mathbb{Z}, 0 < r, r' < L_A$

37 as sender secret keys of the sender A,





38 $N_A (= P_A Q_A)$
 39 as a sender public key of the sender A, and sending the sender secret
 40 key g_A to said trusted center apparatus;
 41 means for generating r, r' , calculating receiver registration key
 42 $r_x(\sigma_x, A)$ satisfying:

43
$$r_x(\sigma_x, A) s_x(\sigma_x, A) \equiv r' \pmod{L_A}$$

 44 from receiver registration data $s_x(\sigma_x, A)$ received from the trusted center,
 45 the sender secret key L_A , and the integer r' , and sending the receiver
 46 registration key $r_x(\sigma_x, A)$ to said receiver apparatus used by the receiver
 47 x; and

48 means for calculating key distribution data defined as:

49
$$z_{Ai} = y_{Ai}^r \pmod{N_A} \quad (1 \leq i \leq m)$$

 50 from key distribution data y_{Ai} received from the trusted center, and the
 51 integer r , and broadcasting the key distribution data z_{Ai} to said
 52 plurality of receiver apparatuses as objects of broadcast
 53 communication;

54 and,

55 said receiver apparatus used by the receiver x comprise:

56 means for calculating an encryption-decryption key K_A for
 57 broadcast communication, using:

58
$$K_A = \left(\prod_{i=1}^k z_{A\sigma_x(i)} \right)^{r_x(\sigma_x, A) s_x(\sigma_x)} \pmod{N_A}$$

59 from the receiver registration key $r_x(\sigma_x, A)$ received from said sender
 60 apparatus used by the sender A, the key distribution data z_{Ai} , and the
 61 receiver secret keys $\sigma_x, s_x(\sigma_x)$ distributed from the trusted center, said
 62 encryption-decryption key K_A being generated by said sender apparatus
 63 used by the sender A, using:

64
$$K_A = g_A^{r'} \pmod{N_A}.$$

This Sheet Has Been Replaced



THIS SHEET HAS BEEN
REPLACED