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# DATA TRANSFER BETWEEN A DATA STORAGE SUBSYSTEM AND HOST SYSTEM

## Technical Field of the Invention

This invention relates to the field of data transfer between data storage subsystems and connected host data processing systems.

## Background of the Invention

Data storage subsystems used in data processing systems commonly comprise a device controller connected to one or more storage devices on which customer data is retained. An important requirement of such subsystems is the ability to store large amounts of data which is at the same time accessible by the host data processing system at high speed. There is an ongoing effort to develop subsystems which can transfer the stored data at ever increasing speeds.

A number of interfaces have been developed to connect the host to the device controller and also the controller to the devices. One such interface is the Small Computer Systems Interface which has been adopted as a standard by the American National Standards Institute. Data Storage subsystems are known which employ a SCSI bus to connect the controller to the device.

A number of techniques for transferring data from the subsystem to the host have been described. In one known technique, for every data transfer operation, the data is checked for errors before it is transmitted to the device. Each sector of data is read from the device and held in a data buffer where it is checked for errors before it is sent up to the attached host data processing system. This technique has the advantage in that only good data is transferred to host. However there is an inherent performance penalty in the increased time taken to transfer all the data.

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In the IBM System/370, another technique is employed whereby the data is sent directly from the device to the host without being checked. This sending of 'Raw' data has a performance benefit over the previously described method in that there is no need to hold a sector in the data buffer before transmission. However, a disadvantage of this technique as employed in the System/370 arises when there is an error in the transmitted data. When this happens, it is necessary to retransmit a large portion of the data originally requested by the host.

A number of performance enhancing techniques have been developed to speed up transmission of data between device and host. One of these is split transfer which reduces the latency caused when the required data is not under the head when the data is requested. US patent 4 494 157 describes the use of split transfer to transfer data in two bursts to a buffer. Both bursts of data are received in the buffer before being sent up to the host system.

## Disclosure of the Invention

The invention seeks to achieve high speed data transfer between a data storage subsystem and attached host data processing system.

Accordingly the invention provides a method of transferring data from a data storage subsystem to an attached host data processing system, the storage system comprising a device controller connected for communication to a direct access storage device on which the data is stored, the method comprising sending a data transfer command from the system to the device controller to initiate data transfer, the command specifying a sequence of blocks of data to be transferred and the start address in host memory to which the first block of the data sequence is to be transferred; and transferring the sequence of data from the subsystem to the system in response to the command, the controller being capable of respecifying the start address in host memory.

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The controller is thus able to respecify the start address in host memory within the range specified in the command. In effect the controller has random access to host memory. This is a facility useful in a number of ways.

In a preferred method, the controller requests the transfer of said sequence of data from the device in a non-sequential order and uses the start address specified by the host to calculate an amended start address in host memory to which the first block of the non-sequential data is directed, the controller transferring the data in non-sequential order to the host, the first block of the data being directed to the amended start address.

Thus in split read transfer, the controller is able to respecify the address in host to which the second burst of data should be sent. In this way, it is not necessary to wait for the second burst of data to be received by the controller before the first burst can be sent to the host. This provides a performance benefit over prior systems which received the data 'out of order' from the device but then had to transfer the data to the host in the correct order.

Alternatively, the controller requests transfer of data from the device in the same order as specified in the command, the controller sending the first block of said sequence of data to the address specified by the host. This is a 'normal' read operation in which data is transferred in ascending block order up to host memory.

In a preferred method, on detection of an error in a portion of the data sequence, the controller requests the device to retransmit said portion and uses the address specified in the command from the host to calculate the address in host memory to which the retry portion is to be directed.

Thus if an error occurs during data transfer, the controller retransmits a portion of the data up to host. The random access to host memory avoids the need to resend the whole of the data specified in the original command.

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In another aspect of the invention there is provided a data storage subsystem for storing data accessible by a connected host data processing system, the subsystem comprising a device controller connected for communication to a direct access storage device wherein in operation data is transferred from the device via the controller to memory in the data processing system, characterised in that the controller is capable of specifying the start address in host memory to which the data is transferred.

A preferred embodiment of the invention will now be described with reference to the accompanying drawings. It should be noted that not all the features of the following description are necessary in order to put the claimed invention into effect.

Brief Description of the Drawings

Figure 1 is a block diagram of the major functional units of a data storage subsystem according to the present invention;

Figure 2 is a block diagram showing the main components of the adapter of Fig.1;

Figure 3 shows the structure of the adapter link chip of Fig.2;

Figure 4 is a block diagram showing the main components of the controller of Fig.1;

Figure 5 is a block diagram showing the structure of the controller link chip of Fig.4;

Figure 6 is a block diagram showing the communication between tasks defined in the controller microprocessor;

Figure 7 is a block diagram showing the inbound and outbound serial links.

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Detailed Description of the Invention

A data storage subsystem is described which is suitable for connection to a data processing system and provides large amounts of storage which may be accessed by the host system at high speed. The main functional units of the subsystem, shown in Fig 1, are (i) Host Adapter, (ii) Device Controller & (iii) Direct Access Storage Device (DASD). The functional units are interconnected by point to point, full duplex serial links. Figure 1 shows a basic configuration of the subsystem wherein one adapter 10 is connected via a dedicated serial link 15 to one controller 20 which is in turn connected by four serial links 25-28 to four DASDs 30. The bulk of the following description will relate to this basic configuration. However the architecture of the subsystem (described in detail below) is designed to allow each adapter to be connected to up to four controllers, each of which may be connected to up to four devices. In a preferred embodiment of the subsystem the host adapter is housed in the host system and is connected via a serial link to a housing (eg rack mounted drawer or free standing unit) which comprises one controller and four DASDs with associated power supply and cooling system (not shown).

Firstly a brief description of the main functions of the adapter, controller and DASD is given.

1. ADAPTER

The adapter is essentially a general purpose multiplexer that connects the host system to the controllers through the serial links. The adapter may be designed to attach to the host system through a variety of existing interfaces e.g. IBM's Micro Channel architecture. (Micro Channel is a trademark of International Business Machines Corporation)

The principal functions of the adapter are as follows:

1) The adapter fetches SCSI (Small Computer System Interface) commands by Direct Memory Access (DMA) from system memory and forwards them to the controller over the serial links.

2) The adapter manages a pool of DMA channels and allocates them on request to the controllers for the transfer of read/write data.

3) The adapter fetches packets of write data from host memory by DMA and transmits them on the serial links to the controllers.

4) The adapter receives packets of read data from the serial links and stores them in system memory by DMA.

5) The adapter assembles the ending status for each command and presents it to the system. Good status can be presented for up to 4 devices at a time.

6) The adapter provides the means to abort a previous SCSI command

2. CONTROLLER

The controller implements the SCSI command set (the members of which relevant to this description are defined elsewhere) for the attached DASD. Its principal functions are as follows:

1) The controller maintains a command queue for each DASD.

2) The controller has a data buffer (shared between the DASDs) to prefetch write data from the host system, to correct read data and to prefetch read data from the DASD.

3) The controller generates SCSI status.

3. DASD

The principal functions of the DASD are as follows:

1) The DASD seeks to the specified cylinder and head.

2) The DASD searches for the starting Logical Block Address (LBA) supplied by the controller and then it reads or writes the specified number of blocks, seeking to the next track as necessary. If any defective blocks are found the DASD skips over them automatically.

3) The DASD generates and checks the ECC bytes that are appended to each block. The ECC hardware is contained in the DASD to allow the controller to support a range of DASD that may have different ECC algorithms.

If the DASD detects a data error then the controller requests the DASD to supply the error pattern and displacement. The controller will then correct the data in its buffer and restart the transfer to the adapter.

4) The DASD has a recording channel for read/write data. The write data is encoded, serialised and fed to the head. The read signal from the head is detected, deserialised and decoded.

SERIAL LINKS

The serial link provides point to point communication between two nodes of the subsystem i.e. between adapter & controller and between controller & DASD.

The unit of data transfer is a packet. The format of a packet is shown below and comprises a control field, an address field, a variable length data field and a CRC field.



Packets may be multiplexed on the serial link to perform several commands on different DASDs simultaneously.

Full duplex data transmission supports the transfer of read data and write data simultaneously. In effect each serial link comprises two links providing data transfer in two opposite directions. This is shown in Fig 7 in which each node has an inbound link on which it receives incoming data and messages and an outbound link over which it transmits data and messages.

The link has a simple protocol. Each node may transmit a packet on its outbound link subject to pacing responses and acknowledgements received from the remote node on its inbound link.

Packets on the serial link may be classified into two types.

Message packets originate from a software process in one node and are addressed to a process in the destination node. (A description of the processes is given below in the sections on Controller and Adapter operations). Message packets are typically used for commands and status (the different types of messages sent in a message packet are described in greater detail below).

Data packets originate from a DMA channel in one node and are addressed to a DMA channel in the destination node. Data packets normally contain read/write data.

Each packet contains an address field which indicates the source and/or destination of the packet. For further details of the operation of the serial links, reference should be made to pending UK applications 9026338.5 and 9026336.9.

**MESSAGES**

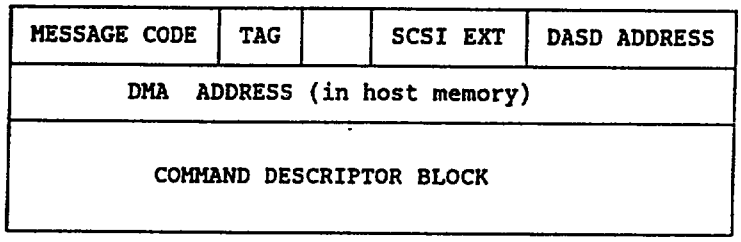
Messages are packets on the serial link that are addressed to a process in the destination node. The first data byte in the packet (that is the

first byte in the data field of the packet) identifies the message. The subsequent bytes are the parameters.

Most messages carry a TAG as a parameter. This allows the messages to be associated with the corresponding command sent from the adapter.

Adapter to controller messages

SCSI_COMMAND




This message transfers a SCSI Command Descriptor Block to the controller queue.

DASD ADDRESS identifies the target storage device that is to execute the command;

SCSI EXT provides an extension or modification beyond the function provided by the SCSI command set as described in ANSI specification 'Small Computer Systems Interface/2' : X3T9.2/86-109. This portion of the message is set to enable Split Write to the DASD or to enable Split Read on the Adapter to Controller link;

DMA Address is the start address in the system memory of the data area for the SCSI command;

COMMAND DESCRIPTOR BLOCK (CDB) is the command descriptor block for the SCSI command. The CDB comprises one of the commands of the SCSI command set.

READY_FOR_READ

| MESSAGE CODE | TAG | LINK ADDRESS |
|---|---|---|

This message is sent by the adapter to the controller in response to a DATA_READY message. The Tag identifies the command with which the particular READY_FOR_READ message is associated. The link address identifies the DMA channel allocated in the adapter for this read operation.

ABORT

| MESSAGE CODE | TAG 1 | TAG 2 |
|---|---|---|

This message is generated by the Adapter when executing the ABORT SCSI operation fetched from the host. TAG 1 identifies the mailbox containing the ABORT_SCSI Command operation; TAG 2 identifies the command to be aborted. The message causes the controller to terminate execution of the command if it is in progress or to remove the command from its queue if execution has not begun.

RESET

| MESSAGE CODE | TAG | Type | DASD ADDRESS |
|---|---|---|---|

This message is sent by the Adapter to the Controller to reset selected resources within the controller or DASD.

Controller to Adapter messages

READY_FOR WRITE

| MESSAGE CODE | TAG | LINK ADDRESS |
|---|---|---|
| DMA START ADDRESS (in host memory) | | |
| DMA LENGTH (bytes) | | |

This message instructs the adapter to transfer data from the host at the DMA start address for the DMA length. LINK ADDRESS identifies the DMA channel in the controller to which the data packets are to be addressed. The Tag identifies the command with which the data is associated.

DATA_READY

| MESSAGE CODE | TAG | |
|---|---|---|
| DMA START ADDRESS (in host memory) | | |
| DATA LENGTH | | |

This message instructs the adapter to allocate a DMA channel to this tag if it has not already done so and prime it for a transfer into host memory beginning at the specified start address and for the specified length. The Adapter responds with a READY_FOR_READ message telling the controller to which DMA channel the data packets should be addressed.

STATUS

| MESSAGE CODE | TAG | STATUS |
|---|---|---|

This message carries the SCSI status generated on completion of the command identified by the tag.

The major components of the controller and adapter will now be described.

ADAPTER

HARDWARE:

The main components of the Adapter hardware are shown in Fig 2. The core of the adapter is the microprocessor chip (MPC) 110 which contains a high performance controller controlling the transfer of messages and data between the host system and the attached controller. There are also two identical adapter link chips (ALC) 120, each of which provides an interface to the Microchannel, two serial links and 16 DMA channels (not shown). Each serial link has four 128 byte packet buffers for commands, data and status in transit to and from the controllers. The interface between the MPC and the ALCs is an I/O bus 115.

Adapter Link Chip (ALC)

The major components of the ALC are shown in Fig 3.

Data RAM:

The Data Ram 121 contains the following areas:

1) Eight 128 byte packet buffers (four for each serial link: two outbound and two inbound);

2) Sixteen 8 byte DMA registers; one DMA register is provided for each DMA channel. The registers are used in DMA data transfer.

3) A 32 byte message buffer for outbound serial link messages: this buffer is used by the high performance microprocessor to build READY_FOR_READ messages destined for the controller. By setting the appropriate hardware, this buffer can then be transmitted to the controller.

4) Host-interface registers; Three interface registers are involved in the process of passing mailboxes to the adapter. Details of mailboxes and adapter operations are given later in the description.

Mailbox Pointer register: 4 Byte register that can be read and written by the system. It is initialised by the system to point to the first mailbox in the chain. The system is only allowed to write to this register when the current Tag Register is equal to the Last Tag register or immediately after the adapter has been reset.

Current Tag Register: This is a 1 byte register that can only be read by the system. It allows the system to view the progress of the adapter. It is cleared by reset of the adapter and immediately after the host writes the Last Tag register. The adapter stores each mailbox tag in the Current Tag register after it completes the processing of a mailbox.

Last Tag Register: 1 Byte register that can be read and written by the system. It is written by the system when it adds some mailboxes to the queue. It indicates the tag which is in the last mailbox. By this means the adapter knows when it has reached the end of the list. When the register is written, the adapter is interrupted.

5) A 32 byte DMA packet buffer for fetching all mailboxes from host: SCSI_COMMAND, ABORT and RESET messages to the controller are sent directly from the DMA buffer. (READY_FOR READ message is built in the message buffer and sent over outbound link to the controller) The DMA buffer can be used to read or write from host memory under DMA control.

The data RAM is time multiplexed between the serial links, the Microchannel or inter-link transfers and the high performance microprocessor.

Status RAM:

The packet buffers each require a packet status register (PSR). These are held in the Status RAM 122 and are 16 bits wide. Packet buffers and

associated Packet Status registers are shown in Figure 7. Each register contains two fields:

DESTINATION - for outbound data packets this field contains a value which will be copied into the address field of the outgoing packet when the contents of the corresponding packet buffer are transmitted by the link. This value may be loaded automatically by hardware when the packet is being fetched from the packet buffer, in preparation for transmission. For inbound packets, this field contains an address extracted from the address field of the incoming packet. This value is written into the PSR by the inbound link FSM and its value is used to determine the subsequent routing of the packet.

BYTE COUNT - for outbound packets this contains a value which indicates the number of bytes which have been placed in the corresponding packet buffer. When the link transmits the packet, this value has to be copied into a byte counter (part of the link hardware) which is decremented as each data byte is sent. The value in the PSR is preserved in case the packet has to be retransmitted due to an error in transmission. For inbound packets, this field contains a value which indicates the number of data bytes which were received in the incoming packet.

Microchannel Interface:

The Microchannel interfaces with host memory and employs the Data RAM host-interface registers defined above.

DMA Channels:

There are sixteen DMA channels (not shown), numbered 0-15 in each ALC. They are employed to DMA data between host memory and the DMA packet buffer.

CONTROLLER

HARDWARE

Figure 4 shows the main functional components of the controller. The core of the controller is the MPC chip 210 that contains the high performance controller (HPC) and a DMA controller which controls the transfer of data to and from the data buffer 220.

A DMA bus 225 connects the DMA controller to two Controller Link Chips 230.

Data Buffer 220 : All data between the adapter(s) and DASDs passes through the data buffer. The buffer is also used for storing read ahead data in case it is subsequently requested by the system (see section on READ AHEAD later in the description). Sixteen DMA channels (numbered 0-15) are provided for transfer of data between the link packet buffers and the data buffer. There are two channels per device (DA) link and 4 channels per SA link.

The data buffer consists of an array of DRAM modules. Data in this buffer is stored with ECC to ensure data integrity. The data buffer is allocated in seven 32K byte segments per DASD. If more than one task is executing on the DASD, a different segment is allocated for each task.

The high performance controller controls the interfaces to the controller via a series of external registers implemented in the Controller Link Chip. The I/O bus 226 is used by the microprocessor to access these registers.

Static RAM 240: used for programme execution.

EPROM 250 : stores the microcode employed in the operation of high performance controller. The structure and operation of the microcode is described in more detail below.

Controller Link Chip (CLC)

The main functional areas contained within a CLC are shown in Figure 5. They are:

1) Two DASD serial interfaces (DA0 and DA1)

2) One adapter serial interface (SA)

3) Link Packet buffers 240 and associated packet status registers 242

Packet buffers are employed to hold incoming and outgoing data. In order to perform back to back link transfers an A/B buffer implementation is used. This allows the DMA logic to fill (or empty) buffer B whilst buffer A is being used by the link and vice versa. The full duplex nature of the link means that both inbound and outbound links need individual sets of packet buffers. Since the controller link chip contains three serial interfaces, this means that a total of 12 packet buffers are required to service the serial links. Additional packet buffers in which microcode can build outgoing Messages are also implemented (one per link). This allows the high performance microprocessor to construct messages without having to withdraw one of the A/B packet buffers from servicing DMA transfer hence preventing ongoing data transfers from being adversely affected.

Each of the three links in the CLC is provided with five packet buffers classified as outbound, inbound or message.

OUTBOUND Each link is provided with two A/B outbound packet buffers which are serviced by the DMA hardware. These buffers are filled with data obtained from the data buffer and sent to the DASD (DA links) or to the adapter (SA link).

INBOUND Each link is also provided with two A/B inbound packet buffers. Incoming packets (from the adapter or the DASD) are stored in these buffers and they are serviced either by DMA hardware or Spinnaker, depending on the content of the address field in the incoming packet.

MESSAGE Each link interface is also provided with a message packet buffer which is used by the microprocessor to build outbound messages to send to the adapter or the DASD.

4) DMA Interface logic: this transfers data from the packet buffer to the controller data buffer under the supervision of the DMA controller.

DMA operations

Data is transferred between the adapter and the controller buffer and between the device and controller buffer by DMA. The controller microprocessor includes a DMA controller which coordinates transfer between the controller link chips and the shared data buffer.

The Controller Link Chip incorporates a DMA interface to transfer data between the packet buffers in the CLC and the controller data buffer. The DMA interface is supervised by the DMA controller contained within the microprocessor chip. The MP chip contains logic to arbitrate between DMA requests.

DMA transfers are preceded by an arbitration phase during which the CLC chips are allowed to signal requests for DMA channels which require servicing. The DMA controller issues a grant to one of these requests, after which point the CLC can initiate transfer. Each controller link chip can use 8 DMA channels to service data transfers. These DMA channels are assigned as follows;

Channel 0,1 : DA link no 0 (DA0 in figure 5)
Channel 2,3 : DA link no 1 (DA1 in figure 5)
Channels 4-7: Adapter Link (SA in figure 5)

(the second CLC uses channels 8 to 15 in the same order as above)

This arrangement allows up to 2 DMA channels to service each DASD and allows the adapter link to be serviced by up to 4 DMA channels. These

channels can be used concurrently to exploit the packet multiplexing feature of the serial link. The device links are given priority over the adapter link. The DMA bus allows a data transfer rate of up to 40 MBytes/s. A 32 byte (DA) transfer from packet buffer to data buffer takes approximately 1.2 microseconds. A 128 byte (SA) transfer takes 3.6 microseconds.

During arbitration, each CLC raises requests on the DMA bus as follows:

STORES (writes to the buffer) A DMA request will be raised on receipt of an inbound link packet, provided that the address field of the packet indicates that the data was destined for a DMA channel.

FETCHES (reads from the buffer) A DMA request will be raised for a DMA channel if either one or both of its associated link packet buffers are empty.

DMA Store operation: DMA store operations are used to empty inbound packet buffers.

DMA Fetch operation: DMA fetch operations are used to fill outbound packet buffers. Each transfer will normally be for a complete packet buffer.

Command Descriptor Queues

The CDB's (plus ABORT and RESET) are stored in Command Descriptor Queue Entries under the control of the SA RECEIVE MESSAGE process when they are first received.

FREE queue: Initially all CDQE's are free. The FREE queue is a queue only in that one member points to the next allowing all free CDQEs to be found. No significance in the order of entries.

NEW COMMAND queue: when a new command arrives, the SA task copies it into the CDQE at the head of the free queue. The CDQE is removed from the free queue and added to the New Command queue. Abort messages are also put in a CDQE and added to the New Command queue.

DEVICE COMMAND Queue: There are four device command queues - one for each device. Each queue is serviced by its own Command process. When the QUEUE MANAGER process finds a new command in a CDQE in the New Command queue, it finds the device to which it is addressed and transfers the CDQE to the corresponding Device Command queue.

MESSAGE Queue: when command processing completes, its CDQE is loaded by the relevant Command process with a SCSI status message and transferred from the Device command queue. The Message queue is thus a queue of requests to the SA Transmit Message process.

The same queues are used for commands received via different Adapter links.

Queue Linkage: The head and tail pointers for the Free, New Command and Message queues are fields in the Message Control Block (MCB). The head and tail pointers for the device command queues are in the corresponding Device Control Block (DCB). The head addresses the first member of the queue. The tail addresses the last member.

Each CDQE contains a 'Next' pointer. Although a CDQE may be in one of several queues it can never be in more than one at a time.

The operation of the adapter and controller will now be described.

**ADAPTER**

OPERATION

Tasks defined within the high speed microprocessor control the operation of the adapter. Tasks are initiated by interrupts, which may be from a

hardware event or via a software interrupt from another task. Software interrupts are the means where one task can set an interrupt to another. The communication between the different tasks is shown in Figure 6.

Tasks

Status: the status task is responsible for managing status to be presented to the host system. Status is passed to this task from one of the other tasks and may be directly presented by writing it to the hardware.

Link: there is one link task to handle the 4 serial links to the controllers. This task is responsible for interpreting any messages received from the controllers(s) and taking appropriate action.

Mailbox: this manages the mailbox interface from the host system. It is responsible for receiving each mailbox from the system. If the mailbox is SEND_SCSI Command, it will be passed to the link task for transmission to the appropriate controller.

(There are other tasks defined but are not described as not being directly relevant to this description).

When instructed by the host processor, the adapter fetches commands from host memory and forwards them immediately to the appropriate controller for execution. The mechanism for fetching commands is dependent on the architecture of the host system and will vary accordingly. In this description, the following mechanism is used on the Microchannel.

The host system initiates subsystem operations by means of Mailboxes which are built in host memory. Each mailbox contains a unique tag which identifies that particular command. For example, when the host wishes to initiate an operation in the subsystem, it will build the operation in the next available mailbox and write the Last Tag Register. Writing the Last Tag Register interrupts the MAILBOX task in the adapter microprocessor. The MAILBOX task instructs the adapter hardware to DMA

the mailbox from host memory into the 32 byte DMA buffer to a designated one of the adapter link chips (master chip) to which all messages from the host are directed.

Once in the DMA buffer, the MAILBOX task decodes the mailbox to determine the type of operation defined by the mailbox contents and if it finds that it is a SEND_SCSI command, the mailbox is converted to a SCSI_COMMAND message for transmission to the appropriate controller. The SCSI_COMMAND message is sent from the 32 byte DMA buffer over the link in the data field of a message packet. The address field of the packet contains the address of the destination which in this case is the controller microprocessor. If the command is destined for a controller which is not serviced by the master chip, it is copied into the DMA buffer in the other ALC and sent over the serial link.

The host defines a number of different operations as well as the SEND_SCSI command, many of which are acted on by the adapter and do not require transmission to the controller. However, two operations namely ABORT_SCSI COMMAND and RESET are passed on to the appropriate controller in the form of ABORT and RESET messages. Details of the format of these messages are given in the list of Adapter-Controller messages elsewhere in the description.

ABORT_SCSI COMMAND and RESET operations are handled by the adapter in essentially the same way as the SEND_SCSI COMMAND operation. The MAILBOX task decodes the mailbox and sends the ABORT or RESET message from the DMA buffer to the appropriate controller. Again, depending on which controller is addressed, it may be necessary to copy the message into the DMA buffer of the second ALC.

The adapter starts a timer for each command that it issues. This serves to detect lost commands or a hung controller without burdening the host system with a large number of timers. The Adapter Idle Task periodically updates the timer and checks that the operation has not timed out.

CONTROLLER

OPERATION

Operation of the controller is achieved by means of tasks defined in microcode within the Spinnaker microprocessor. There are 8 tasks defined in this processor. These include:

4 Device (DA) tasks which manage the interface with the respective device.

An SA task which manages the interface with the adapter/host.

A Command Control Task which is the task with overall control. New SCSI commands are passed to it from SA task. It queues them, decodes them and sends instructions to SA task and appropriate device task. SA and DA tasks perform the data transfer.

The controller uses the above tasks but the SA task and Command Control task are extended via the concept of subtasks. The controller has a number of processes which are implemented as an independent task or as a subtask within a task. Subtasks are run under the control of a subtask scheduler.

Fig 6 is a block diagram showing the communication between the different processes.

Control blocks are used in this communication, one process enters information into control block and posts another process which accesses the information in the control block. Control Blocks are passed between processes.

Controller processes

SA RECEIVE MESSAGE (SARXMSG)

This process deals with all messages sent from the adapter to the controller i.e. SCSI_COMMAND, ABORT, RESET & READY_FOR READ. The format

of these messages can be found elsewhere in this description). A message packet from the adapter is received in the inbound packet buffer of the CLC. The contents of the address field of the incoming packet identify it as a message and it is serviced by the high performance controller. If the message is a new COMMAND, ABORT or RESET the SA Receive message process copies it into the Command Descriptor Queue Entry (CDQE) at the head of the free queue. The CDQE is then enqueued to the Queue Manager process. If the message is READY_FOR READ it is passed to the appropriate SA XFER process (ie the process associated with the device which contains the data to be read).

QUEUE MANAGER

This process is a subtask of the Command Control Task and services interrupts from the SA RECEIVE MESSAGE process. Normally, the message is a SCSI_COMMAND message. It may also be an ABORT or RESET. The SA RECEIVE MESSAGE Process has copied the message into the CDQE and transferred the CDQE from the free queue to the 'New Command Queue' and then posted this process which then transfers the commands from the 'new command' queue to the device specific queue and posts the appropriate command process. The QUEUE MANAGER process carries some limited processing of the message in order to determine which COMMAND process should be posted.

COMMAND

The Command process (one of four subtasks of the command control task) processes the SCSI commands on the Device Command queues. There are four instances of the command process running in parallel, one for each of the four supported devices. Normally, each process handles commands addressed to its device.

Each Command process takes a command off its queue and :

ensures that it is valid to execute at this time
invokes a routine which handles the individual commands which:

validates the command

Instructs the SA XFER (SA) and DEVICE (DA) process

Suspends until the SA and DA have completed

Returns SCSI status

Normally posts the SA TRANSMIT MESSAGE process to send SCSI status to the adapter.

This procedure is repeated for each command in the Device Queue until the queue is exhausted, then it suspends and is resumed when a new command is added to the queue by the QUEUE MANAGER process.

DEVICE

One process per device (implemented as separate tasks). The DA process handles the following requests from the COMMAND process:

1) Read:

the appropriate routine is called to process Read command requested by the COMMAND process. The routine issues a SEEK order to the appropriate DASD, initialises the DMA i.e. allocates a DMA channel, calculates the buffer size available and if there is space available sets up and issues a READ order to the DASD which initiates transfer of data from the DASD to the data buffer in the controller. The DMA address is passed to the DASD in the READ order and is used in the address field of the incoming data packets to identify the destination of the data. When

2) Write:

the routine issues a SEEK order if not already issued

3) Extend Current Operation

4) Stop Event

plus other more specialised commands e.g. Format ID's.

A COMMAND process communicates with a DEVICE process by posting one of these three events

NEWREQ: signals to a device process that a new request is to be started
EXTREQ: this event is used to extend a request
STOP: signals to device process to stop all work currently in process

On receiving a request from the COMMAND process, the DEVICE process will initiate the appropriate action to the DASD by sending the appropriate DASD order over the serial link.

DASD orders are low level read/write orders that are generated by the controller. The following orders are provided to allow the controller to read and write data. Each of the orders defined below are sent to the DASD over the serial link in the data field of a packet. All 'order' packets are addressed to the microprocessor in the DASD for execution or for distribution to other components of the DASD.

STOP_AND_SEEK order

| ORDER CODE | CYLINDER | HEAD |
|---|---|---|

This order instructs the DASD to terminate read-ahead (if active) and seek to a specified cylinder and head. Also, for a write command, the separate seek order allows the controller to initiate the seek as soon as it decodes the command and without waiting to receive the write data from the adapter. If the STOP order is issued before the DASD has completed a read or extended read order (status packet not returned), the DASD terminates the read operation immediately, starts a seek to the cylinder and head specified in the STOP_AND_SEEK order and returns a

status packet for the terminated READ order. No Status packet is sent for the STOP order.

This order is also sent to the DASD on receipt of an ABORT_SCSI command message from the adapter. In this case no seek operation is initiated.

READ order

| ORDER CODE | ADDRESS | CYLINDER | HEAD | LBA | COUNT |
|---|---|---|---|---|---|

This order instructs the DASD to search for a particular Logical Block Address (LBA) and read a specified number of blocks. The parameters are as follows:

The physical cylinder and head for verifying the seek, the logical block address and a count of the number of blocks to be read. In addition, the Address field contains the byte which is to be placed in the address field of any data packets which are returned as a result of this order.

The DASD sends the requested data to the controller and checks the ECC bytes at the end of each block. If the DASD encounters any blocks that are marked as defective then it skips over them automatically. Finally the DASD returns status to indicate whether any errors were detected.

CONDITIONAL_READ order

This order has the same format as the READ order and invokes a seek operation to the selected cylinder and head address. The sector corresponding to the LBA contained in the order is located and as many records as are specified in the count field are read from the disk. The address field contains the byte that is placed in the Address field of any data packets which are returned as a result of this order. Note that the CONDITIONAL_READ order is issued by the controller only when the amount of read data requested by the host is larger than a selected

amount. If the amount of data requested is small then the use of CONDITIONAL_READ is not warranted.

This order enables split read operation, more details of which are given below.

WRITE order

This order instructs the DASD to search for a particular LBA and write a specified number of blocks. The parameters are the same as for a READ order (see above) except that the controller also supplies the data to be written. In addition there is no address field.

CONDITIONAL_WRITE order

The Conditional Write order has the same format as the WRITE order.

EXTEND order

| ORDER CODE | LBA | COUNT |
|---|---|---|

This order extends the operation of a preceding Read, Conditional_Read, Write or Conditional_Write order. The 'count' parameter specifies the number of individual sectors required to be read or written after the current order completes.

The LBA field defines the address of the first block to be read or written. This will be one more than the LBA of the last block to be read or written by the previous order if contiguous reading or writing is required. If the LBA field is not the first block after the last LBA of the previous order, the blocks in between these LBA definitions are skipped over and not read or written.

To be effective, the DASD must receive an EXTEND order before the previous order has completed. The controller uses EXTEND to perform a

back to back write and to continue read ahead. These operations are described in more detail below.

When the entire operation is complete, the appropriate event will be posted back to the Command process.

SA XFER

Per device process which transfers data between host and a read or write buffer in the controller. The COMMAND process can issue the following commands to the SA XFER process:

Send Read Data
Get Write Data
Stop Current Transfer

The parameters required to carry out the data transfer are passed from the COMMAND process to the SA XFER process in a control block.

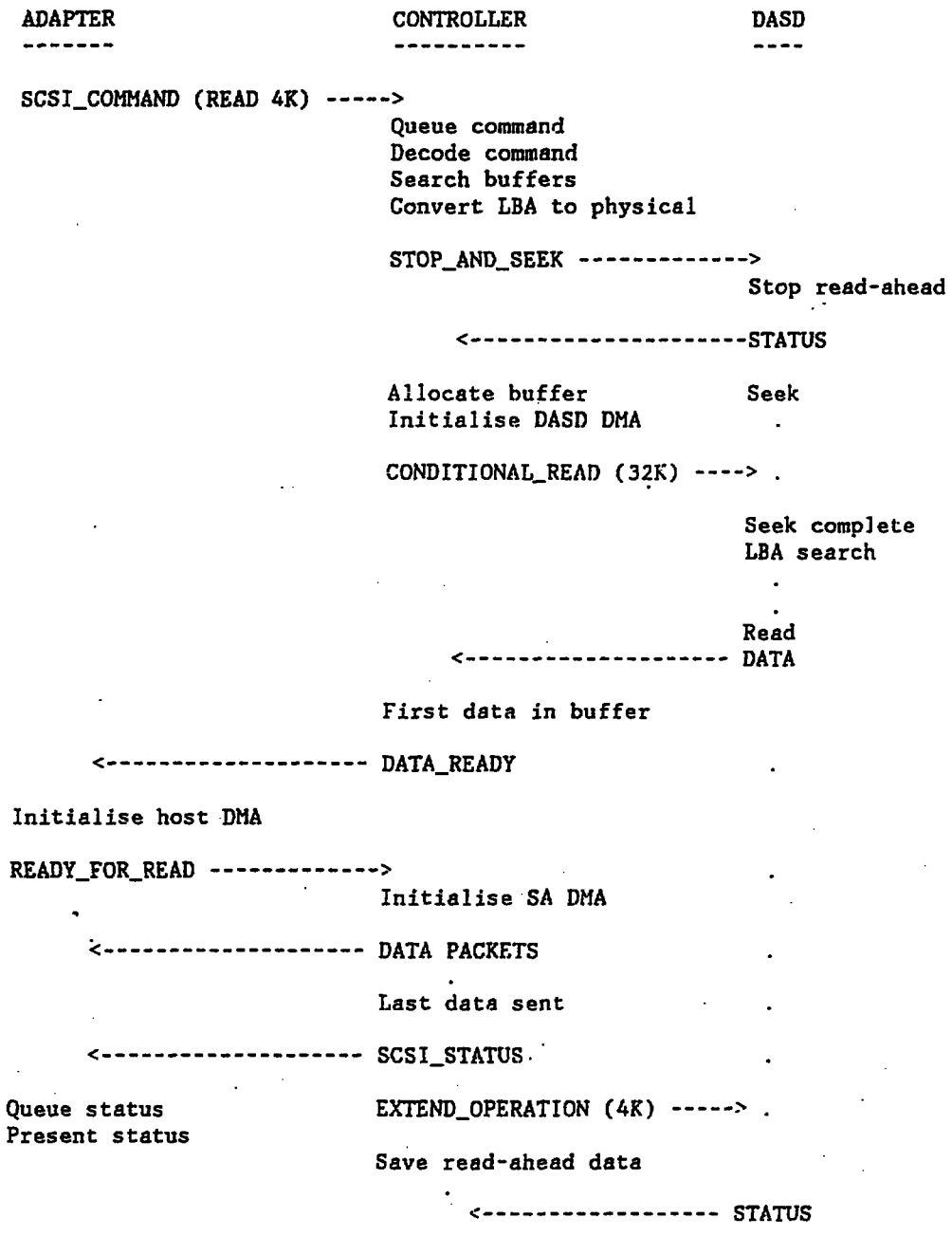
SA TRANSMIT (SATXMSG)

This process transmits messages (READY_FOR_WRITE, DATA_READY & STATUS) to the adapter on behalf of other processes. READY_FOR_WRITE & DATA_READY are passed to this process from the SA XFER process and STATUS messages are passed from the COMMAND process.

EXAMPLE READ AND WRITE OPERATIONS

Examples of typical Read and Write operations beginning with receipt of the command from host to presentation of completion status to the host will now be described.

**EXAMPLE READ OPERATION**



READ:

1. The Adapter sends a SCSI_COMMAND message including a READ operation in the CDB to the controller. The message includes the address of the DASD from which the read data is to be transferred and the address in host memory to which the data is to be sent.

2. The Controller processes the command as described above and passes control to the device task which sends a STOP_AND_SEEK order to the DASD. This terminates any currently active Read Ahead operation. (If no read ahead is currently active, this order is not sent).

3. The DASD returns STATUS to the controller indicating the status of the terminated Read Ahead operation and then begins the seek to the specified head and cylinder.

4. The controller device task allocates a 32K segment of the data buffer for the read data that is to be transferred. The device task also allocates a DMA channel over which the data is to be transferred to the controller data buffer.

5. In this example the device task then sends a CONDITIONAL_READ order to the DASD including the address of the allocated DMA channel, the data start address and the number of blocks to be transferred. As mentioned previously if the amount of requested data is small a 'normal' READ order rather than a CONDITIONAL_READ order will be sent.

6. On receipt of the CONDITIONAL_READ order the DASD searches for the LBA and begins transfer of the data. The address field of the data packet(s) contains the address of the DMA channel.

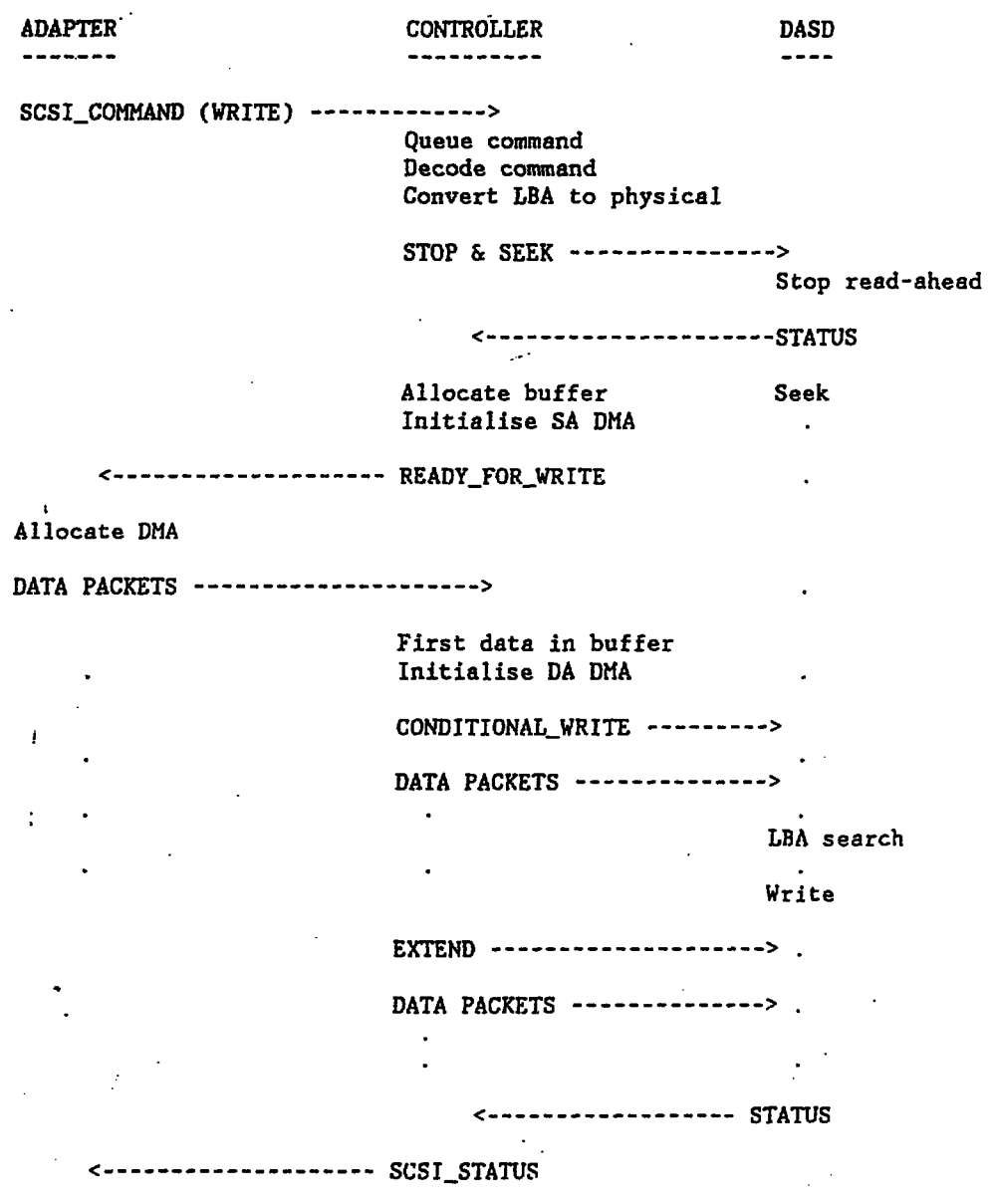
7. The read data is transferred to the controller via the serial link and into the allocated space in the controller data buffer.

8. The Controller sends a DATA_READY message to the adapter which causes the host to initialise the host DMA channel via which the read data is to be transferred between controller and host memory. Although the above flow diagram shows transmission of the DATA_READY message after receipt of the first data into the buffer, it will be more commonly the case that the message is sent before any data is received in the buffer. The purpose of this message is to initialise host DMA ie to prepare the host for receipt of data.

9. The adapter sends a READY_FOR_READ message to the controller in response to the DATA_READY message. The READY_FOR_READ message, which identifies the DMA channel initialised in the host, is received by the controller SA RECEIVE message process and passed to the SA XFER process which initialises the SA DMA ie allocates a DMA channel over which data will be transferred from the data buffer.

10. As data is received in the data buffer from the DASD it is transferred over the serial link into host memory. When the last data is sent, the controller returns SCSI_STATUS to the adapter which queues the status and presents it to the host.

**EXAMPLE WRITE OPERATION**



WRITE:

1. The adapter sends a SCSI_COMMAND message defining a write operation to the controller.

2. The controller processes the command (as described earlier) and control is passed to the DEVICE process from the COMMAND process. In this example, the DEVICE process sends a STOP_AND_SEEK order to the DASD (read ahead operation is currently active).

3. The DASD stops read ahead and sends status to the controller indicating the status of the read ahead operation just terminated. The DASD begins seeking to the cylinder and head specified in the STOP_AND SEEK order.

4.The controller SA XFER process allocates space in the data buffer and initialises the SA DMA channel to which the packets of write data from the host are to be addressed.

5. The SA XFER process posts the SA TRANSMIT process which transmits a READY_FOR_WRITE message to the adapter. This message identifies the DMA channel initialised in the last step.

6. The adapter LINK task allocates a DMA channel which is to be employed to transfer data to host memory and then begins transmission of packets of write data to the controller.

7. The controller DEVICE task initialises a DMA channel to be used for transfer of data between the data buffer and the DASD and in this example sends a CONDITIONAL_WRITE order to the DASD identifying the LBA and amount of write data which the DASD is to expect. The DASD begins the LBA search.

8. As the write data is received over one DMA channel to the buffer it is transferred out of the buffer over a second DMA channel and via the serial link to the DASD.

READ AHEAD

Read Ahead is a function provided by the controller to improve the performance of a set of READ commands which together constitute a long sequential read. This is achieved by continuing reading from the DASD into the controller's buffer in anticipation of the next read as follows:

A SCSI READ command causes the controller to instruct the DASD to transfer the required number of sectors into the 32K segment of the controller's data buffer that has been allocated to this read. In the described embodiment, the controller will normally request 32K of data (i.e. the amount of data to fill the allocated buffer space) even when the command from host requests a lesser amount. As the data requested by the host arrives from the DASD it is transferred to the host. When the transfer of the required data up to host is complete, status is generated by the controller. Meanwhile the transfer of the read ahead data continues.

If the READ order to the DASD specified a larger amount of data than that requested by the host, the extra data will be stored in the controller buffer. As the host makes buffer space available, the data transfer is extended by an EXTEND order sent by the controller to the DASD. On receipt of the EXTEND order, the DASD transfers the sectors specified in the EXTEND order to the new end of the buffer.

As an example, if the READ command sent by the adapter requests 4k of data (8 sectors) and the controller requests the device for 32K, then as soon as the required 4k of data is transferred up to host there will be

4K of space in the buffer. The controller sends an EXTEND order to the device requesting 4K of new data to fill the buffer.

On receipt of the next READ command from the adapter, the controller examines the buffer to see whether it already has the required data (or is about to have). If so it transfers that to the host. If not, it instructs the DASD for the new read, terminating any active read ahead. The Read Ahead continues until the DASD is reinstructed for a new transfer or when the Read Ahead buffer fills.

BACK TO BACK WRITES

Back to back writes are consecutive write commands writing consecutive blocks. The first block of a subsequent write immediately follows the last block of the previous write. Special support for back to back writes in the controller and the DASD allow the writes to be effected without the DASD taking the revolution between the commands which would otherwise be required.

When the command specific routine for the WRITE command reaches a point where it could extend its currently active device transfer to include a following consecutive write, it 'looks over its shoulder' and checks whether the following CDQE contains such a write. If it does, it issues an Extend Request to the device task which sends an EXTEND order to the DASD. If the DASD receives the EXTEND order before the preceding write is finished, it simply extends the count of the current write by the amount specified by the Count parameter in the EXTEND order.

If the EXTEND order reached the DASD too late, the DASD presents status after the first write in the normal way and rejects the EXTEND order with a 'Missed' status. The controller then generates a normal write. EXTEND orders are themselves candidates for extension.

Example:

Host sends SCSI write command for sectors 0 - 7
Controller issues WRITE order for sectors 0 - 7
Host sends SCSI write command for sectors 8 - 15
Controller detects back-to-back write in queue
Controller issues EXTEND order for 8 sectors

If EXTEND order arrives before the end of the WRITE, the DASD extends its current block count by 8 and does not return status after sector 7. Instead it returns status after sector 15.

If the EXTEND order arrives too late, the DASD generates normal status for the WRITE order and tells controller that EXTEND order was received too late. The Controller then reissues a WRITE order for sectors 8-15. Sectors 8 - 15 are then written on the next revolution.

Note that the implementation of back to back writes using the EXTEND order employs the packet multiplexing feature of the serial link i.e. the controller has to be able to send the EXTEND order over the Serial link at the same time as it is transmitting the write data to the DASD.

SPLIT READ

Split Read is a performance enhancement gained by starting reading from the DASD at the first sector, within the range of the current READ operation, which appears under the head, as opposed to the first sector of the range. For example, a read for sectors 4,5,6...15,16 might be read in the order 6,7,8,...15,16,4,5 if the head happened to arrive too late to read sector 4 but early enough to read sector 6. Thus, the transfer from the DASD would complete 11 sectors (i.e 16 minus 5) earlier than would otherwise be the case.

The optimisation is achieved by the controller instructing the DASD to read the first sector. If the DASD then finds that the sector is more than a certain time away (1 millisecond for example), it aborts the read and returns the current LBA to the controller. The controller then makes

the decision whether to reissue the same read, or to split it into a 'tail' read which transfers from the current sector to the end, followed by a 'head' read which transfers from the first sector to the start of the tail.

To support this, the CONDITIONAL READ order to the DASD, which is aborted if there is a long delay to the first sector, is available. The data arrives in the controller buffer, in the above example, 11 sectors earlier than it would if the read had been done in the normal order. However if the controller then has to rearrange the data to be sent to the host in the normal order, then the performance benefit of split read operations is greatly reduced.

The technique described herein realises more of the potential of 'split reads' by avoiding the need to reorder the data in the controller buffer. This is achieved by defining the message interface between the controller and the adapter to allow the controller the necessary control over the host DMA addresses. In effect the controller has random access to host memory. The DATA_READY message sent from the controller to the adapter when data is placed into the controller data buffer specifies the address in host memory to which the data should be directed. When a split read operation is in progress, the controller calculates an amended start address from the address sent in the initial command from the host and sends this to the host in the DATA_READY message. For the example given above, the controller is told the start address for sector 4 and is thus able to calculate the amended address in host where the first sector of the split read data (in this case sector 6) should be directed. If sectors 6 to 16 are already in the data buffer, then the controller will issue a single DATA_READY message to initiate transfer of sectors 6 to 16. The transfer of sectors 6 to 16 will complete. When the controller has sectors 1 to 5 in its buffer, it will then issue a DATA_READY message to the controller indicating readiness to transfer sectors 1 to 5 and the address in host memory to which sector 1 is to be directed.

RAW READ

The controller's ability to randomly access areas of the host memory ie control the host address also facilitates another performance benefit. During any read operation when the DASD is in the process of transferring multiple sectors of data to the controller buffer, the DASD doesn't wait to the end of a sector (512 bytes) to check the ECC before sending the sector up to the host. In the system described, the DASD doesn't have sufficient buffering to store a sector's worth of data. The data is read serially from the disk and compiled into 128 byte packets for transmission to the the controller. At the end of the sector, the device checks the ECC and appends 6 ECC bytes to the end of the data. If all the data in the sector was 'good', no error will be indicated and the controller will continue transmission to host. It should be noted that in this system, most of the sector of data will have already been transferred from the controller to the host before the device sends the last packet of data to the controller and therefore before the device is aware of any errors contained in the data. Thus the data already in the host will be 'raw' data ie there was no checking before transfer to the host.

For most data transfers, the raw data will also be good data i.e. the data is transferred from the device. Transfer of data directly to host without checking provides a performance benefit over prior systems in which each sector is checked any of the sector is transmitted to the host. The time taken for multiple sectors of data to be received by the host is reduced in the direct case by approximately the time taken for one sector of data to travel from device to host. In high performance systems such as described herein, this benefit can reduce the overall overhead by a significant amount.

As described, for most data transfers the sending of raw data as opposed to checked data results in a performance benefit. To be traded against this is the reduced performance when there is an error in the transferred data. In one prior system, if there is an error in the raw

data, the whole of the data must be retransmitted. In the present invention wherein the controller has random access to host memory, there is no need to retransmit all the data.

If, at the end of a sector, the DASD detects the presence of an error in that data, the ECC bytes indicate this to the controller which is alerted. When it has completed sending the current packet of data to the host, the controller ceases transmission and sends a DATA_RETRY message to the host. The controller requests retransmission of the sector which contained the error. The DASD retransmits the requested sector in 128 byte packets to the controller which stores the data in its buffer. If the retransmitted block of data is good, the controller passes the data up to the host.

The controller sends a DATA_RETRY message to the adapter which specifies the address in host memory to which the retry data is to be directed. The DATA_RETRY message instructs the adapter to set up a new DMA channel over which the data is to be transferred into host memory. The adapter responds to the DATA_RETRY message with a READY_FOR_READ message. The DATA_RETRY message indicates the amount of data which is being retried.

If the retry data sent to the controller is still in error, the controller will re-request the data a number of times. If, after a predetermined number of attempts, the controller still hasn't received good data, it will attempt to correct the error in the data held in its buffer. To this end, it sends 'Operate on ECC' order to the device, which causes the device to calculate which bytes were in error and also the correct data for those bytes. The number of bytes which may be corrected is implementation dependent: in the described system the number is two. The device sends the correction information to the controller which corrects the data held in its buffer. The controller now sends the block of data including the corrected bytes to the host.

Claims

1. A method of transferring data from a data storage subsystem to an attached host data processing system, the storage system comprising a device controller connected for communication to a direct access storage device on which the data is stored, the method comprising

sending a data transfer command from the system to the device controller to initiate data transfer, the command specifying a sequence of blocks of data to be transferred and the start address in host memory to which the first block of the data sequence is to be transferred; and

transferring the sequence of data from the subsystem to the system in response to the command, the controller being capable of respecifying the start address in host memory.

2. A method of transferring data as claimed in claim 1;

in which the controller requests the transfer of said sequence of data from the device in a non-sequential order, the controller using the start address specified by the host to calculate an amended start address in host memory to which the first block of the non-sequential data is directed, the controller transferring the data in non-sequential order to the host, the first block of the data being directed to the amended start address.

3. A method of transferring data as claimed in claim 1;

the controller requesting transfer of data from the device in the same order as specified in the command, the controller sending the first block of said sequence of data to the address specified by the host.

4. A method as claimed in claim 3, wherein on detection of an error in a portion of the data sequence, the controller requests the device to retransmit said portion and uses the address specified in the command

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from the host to calculate the address in host memory to which the retransmitted portion is to be directed.

5. A method as claimed in claim 4, the controller requesting retransmission of said portion until either the portion is received in host without error or retransmission has been attempted a predetermined number of times whereupon the controller requests correction data from the device, corrects the errors in the portion and uses the address specified in the command from the host to calculate the address in host memory to which the corrected portion is to be directed.

6. A data storage subsystem for storing data accessible by a connected host data processing system, the subsystem comprising a device controller connected for communication to a direct access storage device wherein in operation data is transferred from the device via the controller to memory in the data processing system, characterised in that the controller is capable of specifying the start address in host memory to which the data is transferred.

7. A data storage subsystem as claimed in claim 6, the controller including a data buffer via which the data is transferred from the device to host memory.

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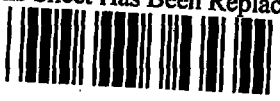
# DATA TRANSFER BETWEEN A DATA STORAGE SUBSYSTEM AND HOST SYSTEM

## ABSTRACT

Described is a technique for transferring data from a data storage system to a connected host data processing system. The subsystem comprises a device controller connected to one or more direct access storage devices e.g. disk drives. The system issues data transfer commands to the subsystem to iniciate transfer of data between system and device(s). Read/write data is transferred directly from device to host via a buffer in the controller. For a read operation, the read command specifies the data to be transferred and the start address in host memory to which it should be sent. The device controller is capable of amending the start address to which it actually sends the data. This provides a performance benefit for split data transfers. In addition, if an error occurs during a read operation, the controller can specify the host address to which the replacement data should be sent.

Fig. 1.

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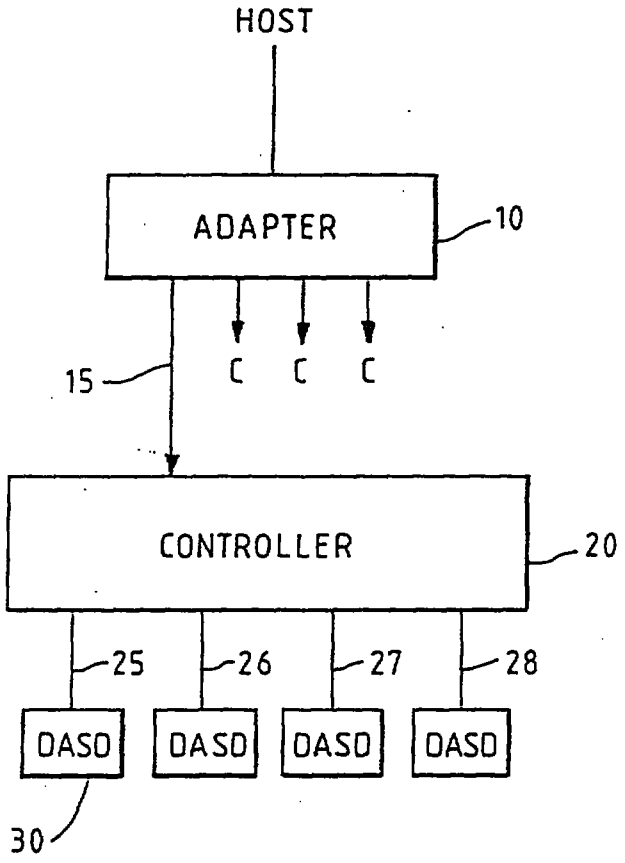

HOST
ADAPTER
10
C C C
15
CONTROLLER
20
25
26
27
28
DASD
DASD
DASD
DASD
30

FIG 1

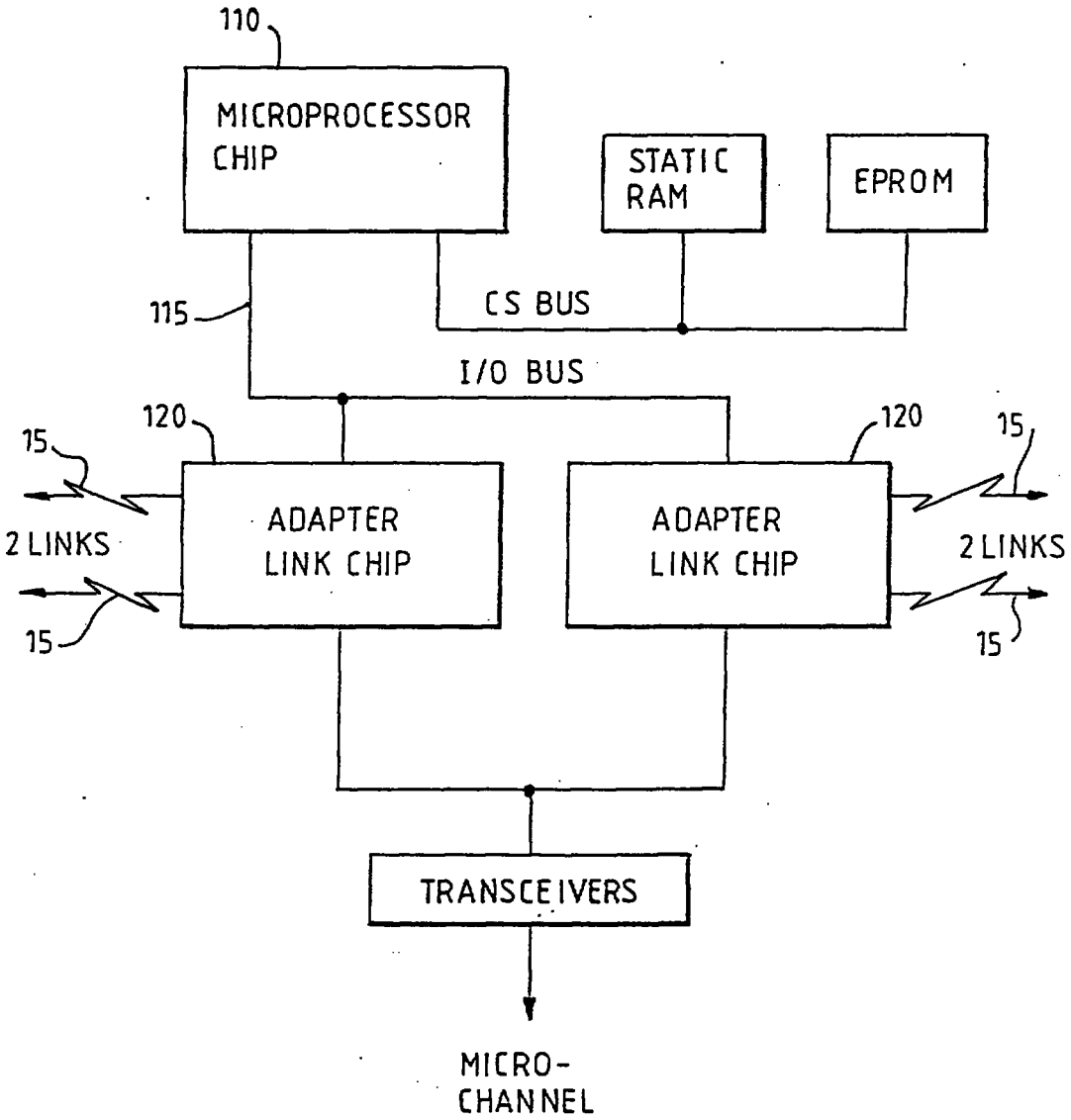

110
MICROPROCESSOR CHIP
STATIC RAM
EPROM
115
CS BUS
I/O BUS
120
120
15
15
2 LINKS
ADAPTER LINK CHIP
ADAPTER LINK CHIP
2 LINKS
15
15
TRANSCEIVERS
MICRO-CHANNEL

FIG. 2

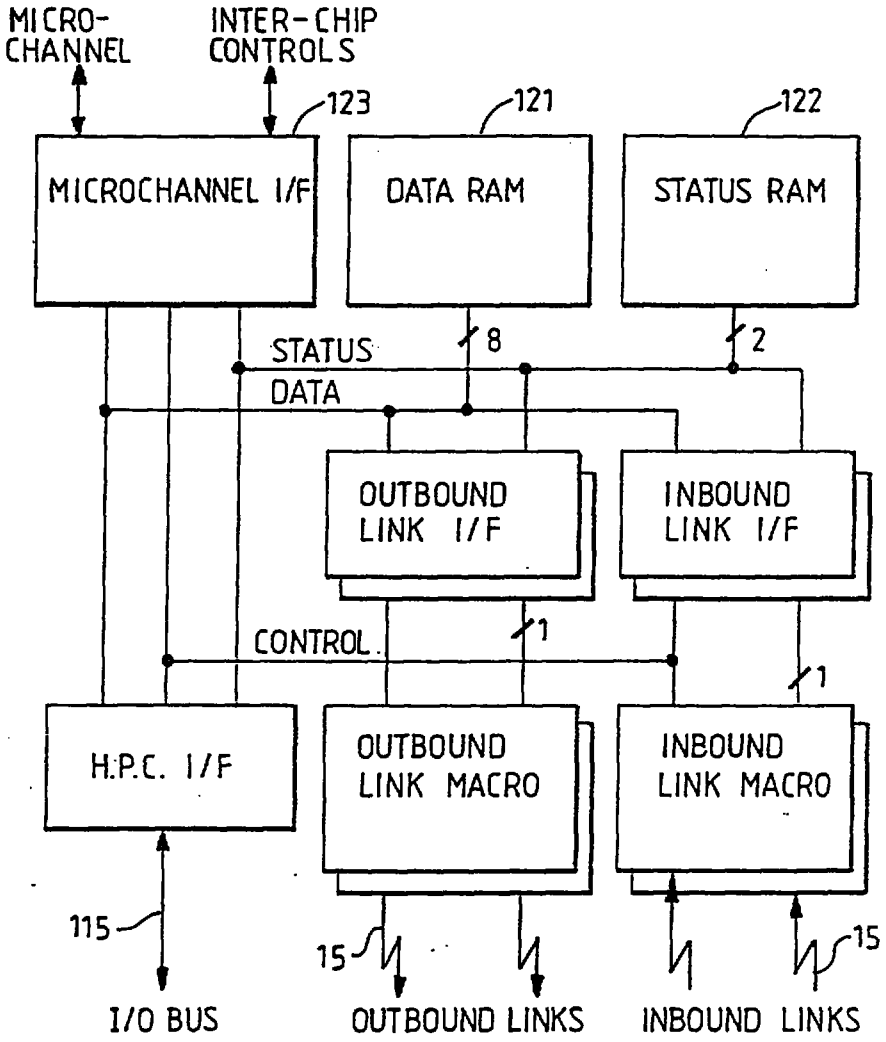

MICRO-CHANNEL
INTER-CHIP CONTROLS
123
121
122
MICROCHANNEL I/F
DATA RAM
STATUS RAM
STATUS
DATA
8
2
OUTBOUND LINK I/F
INBOUND LINK I/F
CONTROL
1
1
H.P.C. I/F
OUTBOUND LINK MACRO
INBOUND LINK MACRO
115
15
15
I/O BUS
OUTBOUND LINKS
INBOUND LINKS

FIG. 3

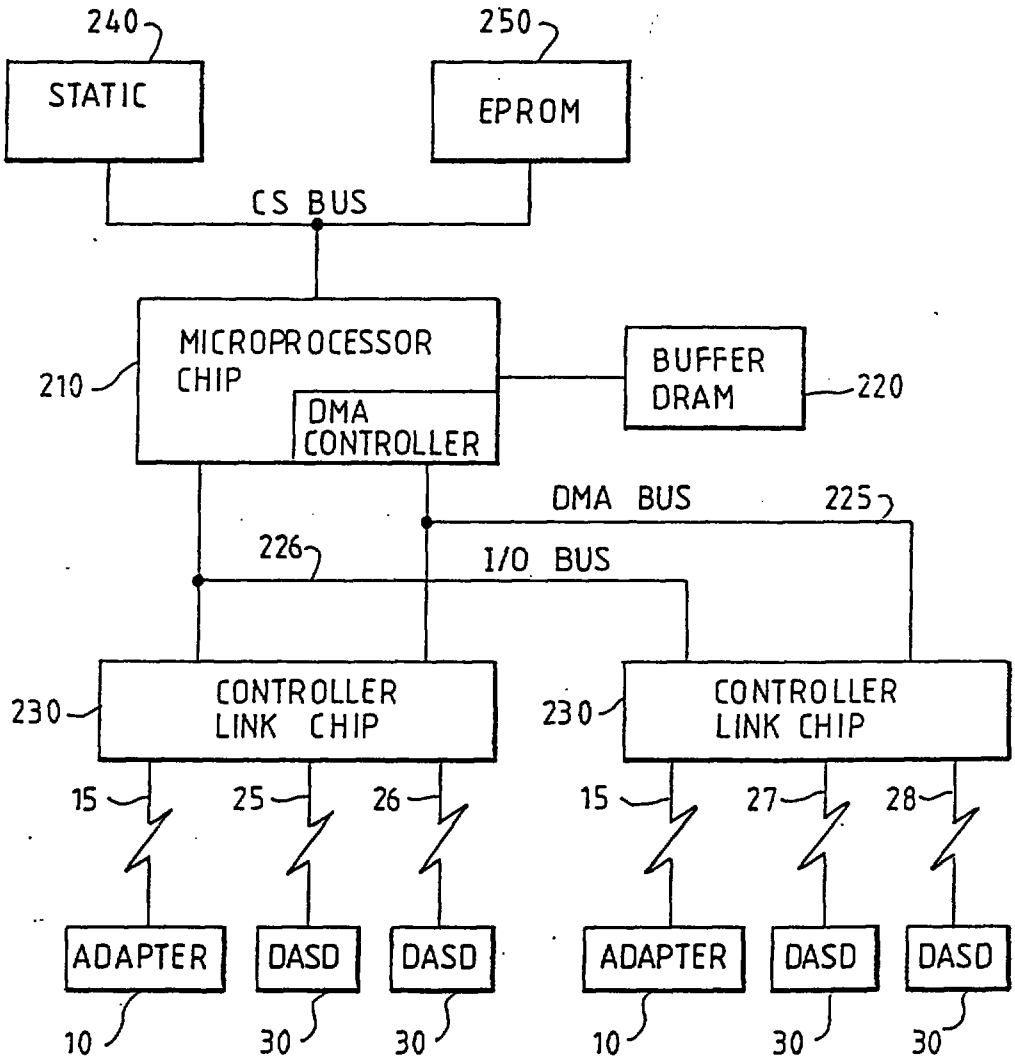

240
STATIC
250
EPROM
CS BUS
MICROPROCESSOR CHIP
DMA CONTROLLER
210
BUFFER DRAM
220
DMA BUS
225
226
I/O BUS
230
CONTROLLER LINK CHIP
230
CONTROLLER LINK CHIP
15
25
26
15
27
28
ADAPTER
DASD
DASD
ADAPTER
DASD
DASD
10
30
30
10
30
30

FIG. 4

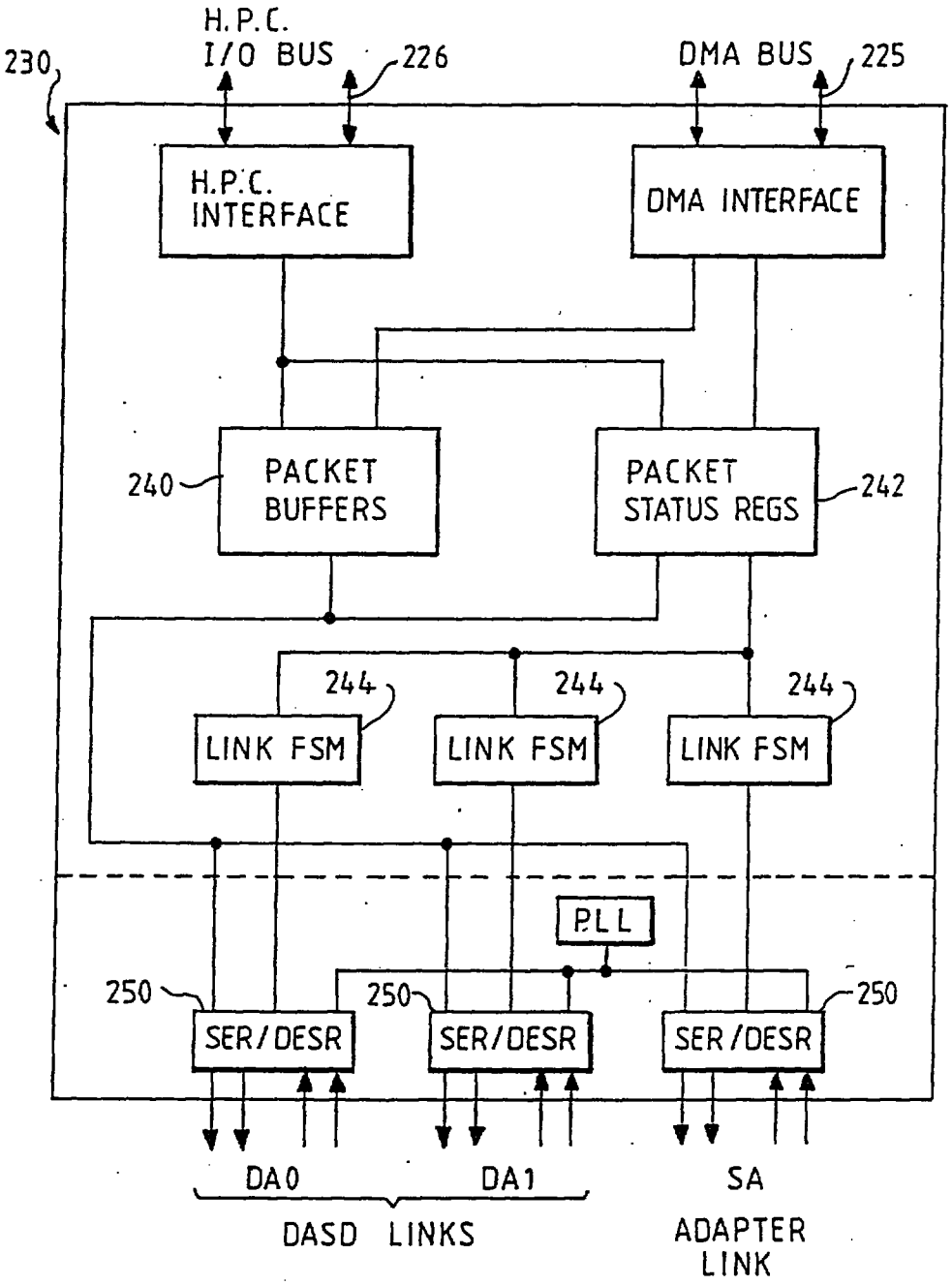

H.P.C.
I/O BUS
226
230
DMA BUS
225
H.P.C.
INTERFACE
DMA INTERFACE
240
PACKET
BUFFERS
PACKET
STATUS REGS
242
244
244
244
LINK FSM
LINK FSM
LINK FSM
PLL
250
250
250
SER/DESR
SER/DESR
SER/DESR
DA0
DA1
SA
DASD LINKS
ADAPTER
LINK

FIG. 5

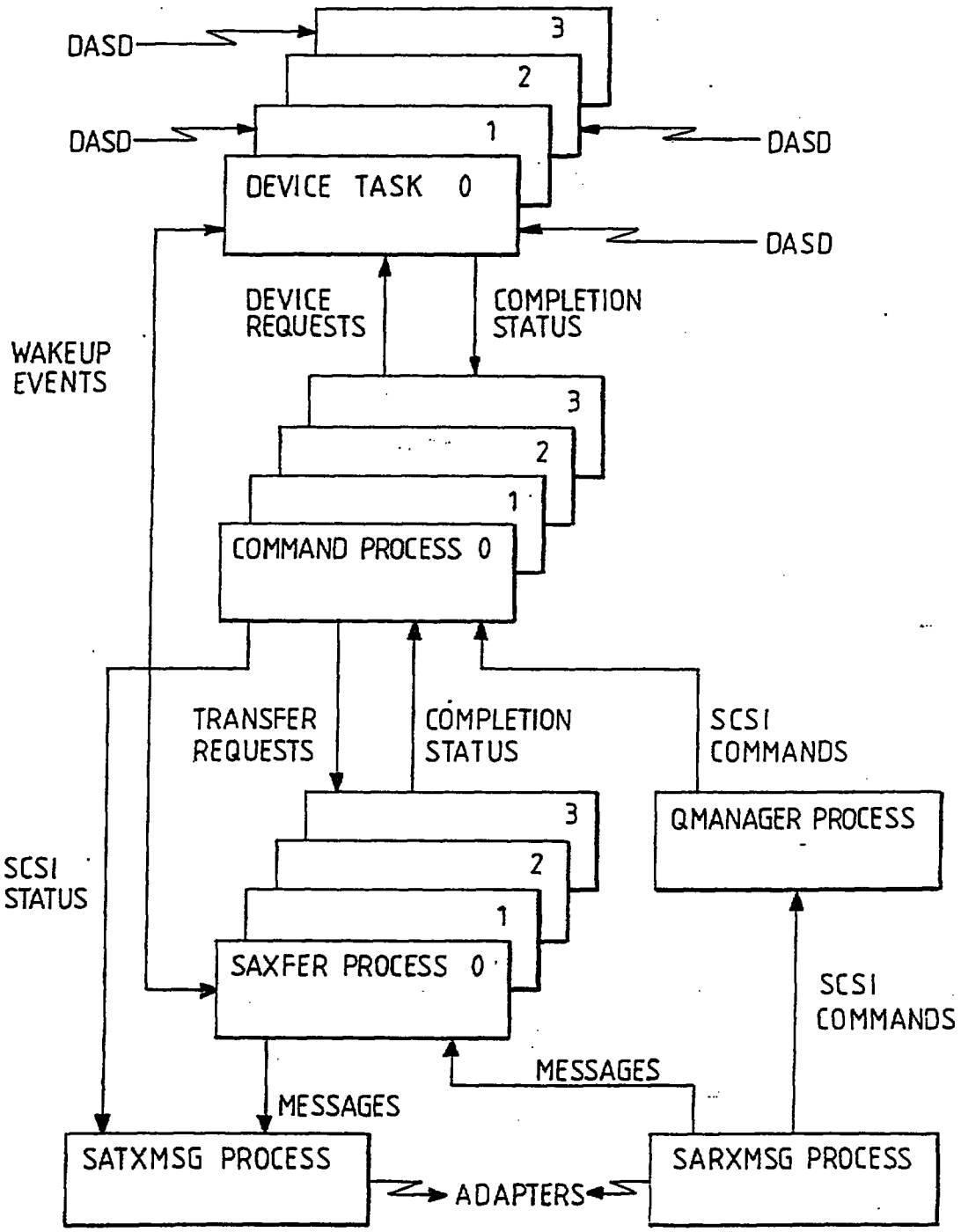
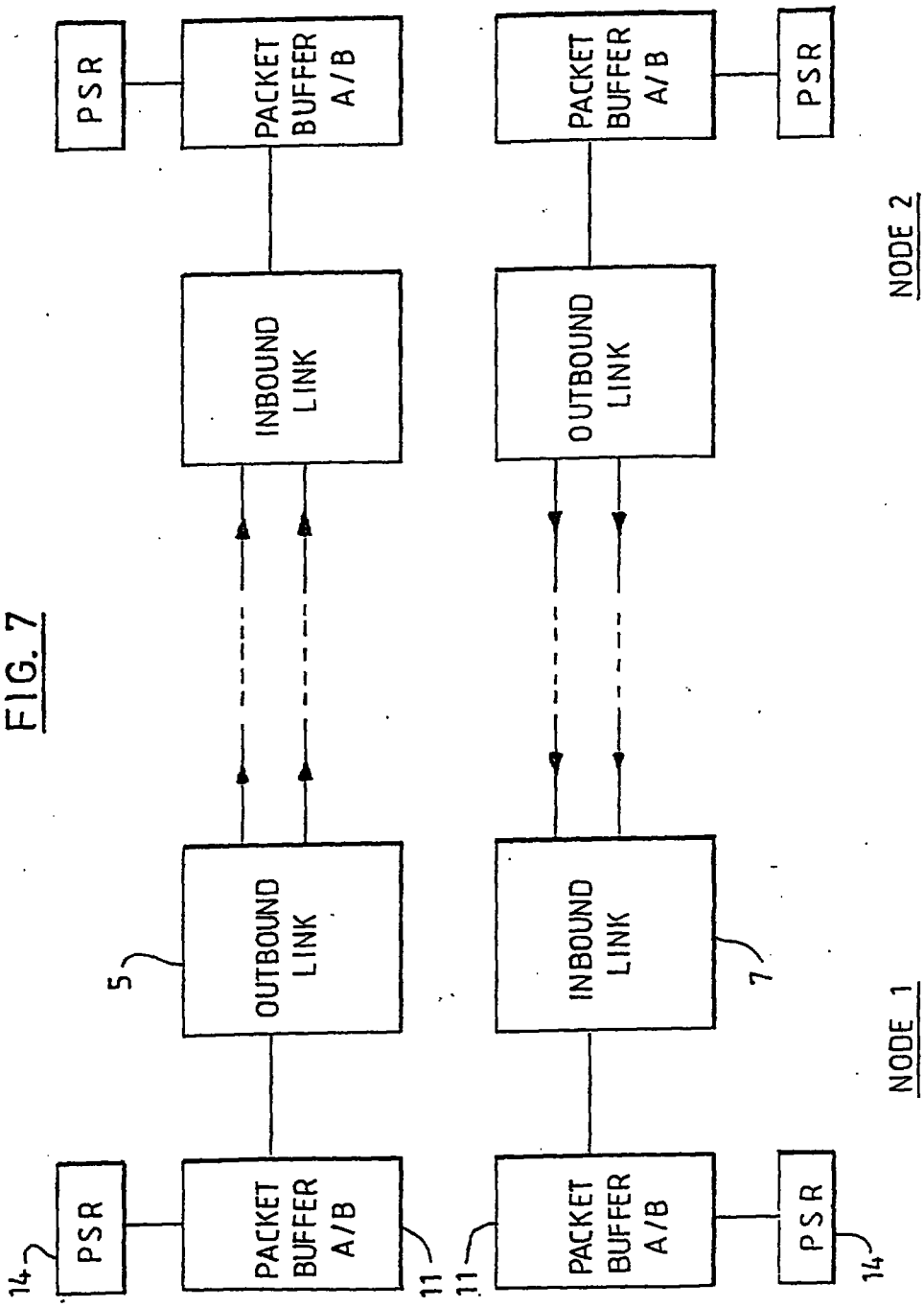

DASD
DASD
DASD
DASD
3
2
1
DEVICE TASK 0
DEVICE REQUESTS
COMPLETION STATUS
WAKEUP EVENTS
3
2
1
COMMAND PROCESS 0
TRANSFER REQUESTS
COMPLETION STATUS
SCSI COMMANDS
QMANAGER PROCESS
SCSI STATUS
3
2
1
SAXFER PROCESS 0
SCSI COMMANDS
MESSAGES
MESSAGES
SATXMSG PROCESS
ADAPTERS
SARXMSG PROCESS

FIG. 6

FIG. 7


PSR
PACKET BUFFER A/B
INBOUND LINK
OUTBOUND LINK
NODE 2
NODE 1
5
7
11
14

# DATA TRANSFER BETWEEN A DATA STORAGE SUBSYSTEM AND HOST SYSTEM

## ABSTRACT

Described is a technique for transferring data from a data storage system to a connected host data processing system. The subsystem comprises a device controller connected to one or more direct access storage devices e.g. disk drives. The system issues data transfer commands to the subsystem to iniciate transfer of data between system and device(s). Read/write data is transferred directly from device to host via a buffer in the controller. For a read operation, the read command specifies the data to be transferred and the start address in host memory to which it should be sent. The device controller is capable of amending the start address to which it actually sends the data. This provides a performance benefit for split data transfers. In addition, if an error occurs during a read operation, the controller can specify the host address to which the replacement data should be sent.

Fig. 1.

REPLACED SHEETS SG
☐ Description : paged .... to paged ....
☐ Claim(s) : paged .... to paged ....
☐ Drawing(s) : paged .... to paged ....
☐ Abstract : paged 42 to paged ....

REPLACEMENT SHEETS SG
☐ Description : paged .... to paged ....
☐ Claim(s) : paged .... to paged ....
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In the IBM System/370, another technique is employed whereby the data is sent directly from the device to the host without being checked. This sending of 'Raw' data has a performance benefit over the previously described method in that there is no need to hold a sector in the data buffer before transmission. However, a disadvantage of this technique as employed in the System/370 arises when there is an error in the transmitted data. When this happens, it is necessary to retransmit a large portion of the data originally requested by the host.

A number of performance enhancing techniques have been developed to speed up transmission of data between device and host. One of these is split transfer which reduces the latency caused when the required data is not under the head when the data is requested. US patent 4 494 157 describes the use of split transfer to transfer data in two bursts to a buffer. Both bursts of data are received in the buffer before being sent up to the host system.

## SUMMARY OF THE INVENTION

It is an object of the present invention to ameliorate one or more disadvantages of the prior art.

According to one aspect of the present invention there is provided a method of transferring data from a data storage subsystem to an attached host data processing system having a host memory, said host memory having a plurality of addresses, the data storage subsystem comprising a device controller connected for communication to a direct access storage device on which the data is stored, the method comprising the steps of: sending a data transfer command from the host data processing system to the device controller of the data storage subsystem to initiate data transfer, the data transfer

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command specifying a sequence of blocks of data to be transferred, said sequence of blocks further comprising a first set of blocks and a second set of blocks, said sequence of blocks in sequential order when a first set of blocks of data is followed by a second set of blocks of data, and said data transfer command specifying a start address in the host memory to which the first block of the first set of blocks of data in the sequence of blocks of data is to be transferred; and transferring the sequence of blocks of data from the data storage subsystem to the host data processing system in response to the data transfer command as a first set of blocks and a separate second set of blocks, the device controller of the data storage subsystem capable of calculating and specifying an amended start address in host memory to which the second set of blocks is to be transferred, said amended start address being different from the start address specified in the data transfer command.

According to another aspect of the present invention there is provided a data storage subsystem for storing data accessible by a connected host data processing system having memory, the subsystem comprising; a device controller connected for communication to a direct access storage device which transfers blocks of data from the direct access storage device via the controller to memory in the host data processing system responsive to a command from the host data processing system, said command from the host specifying a start address in the memory of the host data processing system; said device controller further including means for sending the blocks of data from the direct access storage device to the memory in the host data processing system in other than

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normal sequential order, said device controller capable of calculating and specifying an amended start address in host memory to which a second portion of data is initially transferred, said amended start address originating in the device controller rather than the host data processing system, and said amended start address different than the start address specified by the host data processing system; means for sending a first portion of data to the address specified in the command from the host data processing system so that after sending the first portion of data to the memory of the host data processing system, the data appears sequentially in the memory of the host processing system as if it had been sent in sequential order.

Brief Description of the Drawings

Figure 1 is a block diagram of the major functional units of a data storage subsystem according to the present invention;

Figure 2 is a block diagram showing the main components of the adapter of Fig. 1;

Figure 3 shows the structure of the adapter link chip of Fig.2;

Figure 4 is a block diagram showing the main components of the controller of Fig.1;

Figure 5 is a block diagram showing the structure of the controller link chip of Fig.4;

Figure 6 is a block diagram showing the communication between tasks defined in the controller microprocessor;

Figure 7 is a block diagram showing the inbound and outbound serial links.

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## CLAIMS

1. A method of transferring data from a data storage subsystem to an attached host data processing system having a host memory, said host memory having a plurality of addresses, the data storage subsystem comprising a device controller connected for communication to a direct access storage device on which the data is stored, the method comprising the steps of:

sending a data transfer command from the host data processing system to the device controller of the data storage subsystem to initiate data transfer, the data transfer command specifying a sequence of blocks of data to be transferred, said sequence of blocks further comprising a first set of blocks and a second set of blocks, said sequence of blocks in sequential order when a first set of blocks of data is followed by a second set of blocks of data, and said data transfer command specifying a start address in the host memory to which the first block of the first set of blocks of data in the sequence of blocks of data is to be transferred; and

transferring the sequence of blocks of data from the data storage subsystem to the host data processing system in response to the data transfer command as a first set of blocks and a separate second set of blocks, the device controller of the data storage subsystem capable of calculating and specifying an amended start address in host memory to which the second set of blocks is to be transferred, said amended start address being different from the start address specified in the data transfer command.

2. The method of transferring blocks of data as claimed in claim 1 wherein the step of transferring the sequence of data from the data storage subsystem includes the steps of:

the device controller requesting the transfer of said sequence of data from the device controller to the host in a non-sequential order,

the device controller using the start address for the first block of the first set of blocks as specified by the host and the length of the first set of blocks, to calculate an amended start address in host memory to which the first block of the second set of blocks of data is directed.

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the device controller transferring the second set of blocks of data before the first set of blocks of data to the host, and

the first block of the second set of blocks of the data directed to the amended start address.

3. The method of transferring data as claimed in claim 1 wherein the step of transferring the sequence of data from the data storage subsystem further includes the steps of:

the device controller requesting transfer of data from the device in the same order as specified in the command, and

the device controller sending the first set of blocks of said sequence of data to the address specified by the host.

4. The method of transferring data as claimed in claim 3, further comprising the steps of:

detecting an error in the second set of blocks of data in the sequence of blocks of data,

the device controller using the start address for the first block of the first set of blocks as specified by the host and the length of the first set of blocks, to calculate an amended start address in host memory to which the first block of the second set of blocks of data is directed, and

the device controller retransmitting the second set of blocks from the data storage subsystem to the amended start address in the host memory.

5. The method of data transfer as claimed in claim 4 wherein the step of the device controller retransmitting the second set of blocks to the amended start address is repeated until the second set of blocks is received in host without error.

6. The method of data transfer as claimed in claim 4 wherein the step of the device controller retransmitting the second set of blocks to the amended start address further comprises the steps of:

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the device controller repeating the retransmission request a predetermined number of times without receiving a correct second set of blocks;

the device controller requesting correction data from a disk drive;

the device controller correcting the errors in the second set of blocks.

7. A data storage subsystem for storing data accessible by a connected host data processing system having memory, the subsystem comprising;

a device controller connected for communication to a direct access storage device which transfers blocks of data from the direct access storagx e device via the controller to memory in the host data processing system responsive to a command from the host data processing system, said command from the host specifying a start address in the memory of the host data processing system;

said device controller further including means for sending the blocks of data from the direct access storage device to the memory in the host data processing system in other than normal sequential order, said device controller capable of calculating and specifying an amended start address in host memory to which a second portion of data is initially transferred, said amended start address originating in the device controller rather than the host data processing system, and said amended start address different than the start address specified by the host data processing system;

means for sending a first portion of data to the address specified in the command from the host data processing system so that after sending the first portion of data to the memory of the host data processing system, the data appears sequentially in the memory of the host processing system as if it had been sent in sequential order.

8. The data storage subsystem as claimed in claim 7, the device controller further including a data buffer via which the blocks of data are transferred from the direct access storage device to the memory of the host data processing system.

# DATA TRANSFER BETWEEN A DATA STORAGE SUBSYSTEM AND HOST SYSTEM

## ABSTRACT

Described is a technique for transferring data from a data storage system to a connected host data processing system. The subsystem comprises a device controller connected to one or more direct access storage devices e.g. disk drives. The system issues data transfer commands to the subsystem to iniciate transfer of data between system and device(s). Read/write data is transferred directly from device to host via a buffer in the controller. For a read operation, the read command specifies the data to be transferred and the start address in host memory to which it should be sent. The device controller is capable of amending the start address to which it actually sends the data. This provides a performance benefit for split data transfers. In addition, if an error occurs during a read operation, the controller can specify the host address to which the replacement data should be sent.

Fig. 1.

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