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REDUCED COMPUTATION IN JOINT DETECTION

BACKGROUND

The invention generally relates to wireless communication systems. In particular, the invention relates to joint detection of multiple user signals in a wireless communication system.

Figure 1 is an illustration of a wireless communication system 10. The communication system 10 has base stations 12_1 to 12_5 which communicate with user equipments (UEs) 14_1 to 14_3 . Each base station 12_1 has an associated operational area where it communicates with UEs 14_1 to 14_3 in its operational area.

In some communication systems, such as code division multiple access (CDMA) and time division duplex using code division multiple access (TDD/CDMA), multiple communications are sent over the same frequency spectrum. These communications are typically differentiated by their chip code sequences. To more efficiently use the frequency spectrum, TDD/CDMA communication systems use repeating frames divided into time slots for communication. A communication sent in such a system will have one or multiple associated chip codes and time slots assigned to it based on the communication's bandwidth.

Since multiple communications may be sent in the same frequency spectrum and at the same time, a receiver in such a system must distinguish between the multiple communications. One approach to detecting such signals is single user detection. In single user detection, a receiver detects only the communication from a desired transmitter using a code associated with the desired transmitter, and treats signals of other transmitters as interference.

In some situations, it is desirable to be able to detect multiple communications simultaneously in order to improve performance. Detecting multiple communications simultaneously is referred to as joint detection. Some joint detectors use Cholesky decomposition to perform a minimum mean square error (MMSE)



detection and zero-forcing block equalizers (ZF-BLEs). These detectors have a high complexity requiring extensive receiver resources.

Karimi et al., "A Novel and Efficient Solution to Block-based Joint-detection using Approximate Cholesky Factorization", discloses a data detection approach. The channel codes and the channel impulse responses are modeled as a channel response matrix, A . The received data is modeled using a zero forcing solution. In solving the model; and inversion of a system response matrix $A^H A$ is required. The solution is achieved using an approximate Cholesky factor.

Accordingly, it is desirable to have alternate approaches to joint detection.

SUMMARY

A plurality of transmitted data signals are received at a receiver. The receiver measures a channel response associated with the transmitted data signals. A system response is determined. The system response is expanded to be piecewise orthogonal. The received data signals data is retrieved based on in part the expanded system response.

BRIEF DESCRIPTION OF THE DRAWING(S)

Figure 1 is a wireless communication system.

Figure 2 is a simplified transmitter and a receiver using joint detection.

Figure 3 is an illustration of a communication burst.

Figure 4 is an illustration of reduced computation joint detection.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT(S)

Figure 2 illustrates a simplified transmitter 26 and receiver 28 using joint detection in a TDD/CDMA communication system. In a typical system, a transmitter 26 is in each UE 14₁ to 14₃ and multiple transmitting circuits 26 sending multiple communications are in each base station 12₁ to 12₅. A base station 12₁ will typically require at least one transmitting circuit 26 for each actively communicating UE 14₁ to 14₃. The joint detection receiver 28 may be at a base station 12₁, UEs 14₁ to 14₃

or both. The joint detection receiver 28 receives communications from multiple transmitters 26 or transmitting circuits 26.

Each transmitter 26 sends data over a wireless communication channel 30. A data generator 32 in the transmitter 26 generates data to be communicated over a reference channel to a receiver 28. Reference data is assigned to one or multiple codes and/or time slots based on the communications bandwidth requirements. A spreading and training sequence insertion device 34 spreads the reference channel data and makes the spread reference data time-multiplexed with a training sequence in the appropriate assigned time slots and codes. The resulting sequence is referred to as a communication burst. The communication burst is modulated by a modulator 36 to radio frequency. An antenna 38 radiates the RF signal through the wireless radio channel 30 to an antenna 40 of the receiver 28. The type of modulation used for the transmitted communication can be any of those known to those skilled in the art, such as direct phase shift keying (DPSK) or quadrature phase shift keying (QPSK).

A typical communication burst 16 has a midamble 20, a guard period 18 and two data bursts 22, 24, as shown in Figure 3. The midamble 20 separates the two data bursts 22, 24 and the guard period 18 separates the communication bursts to allow for the difference in arrival times of bursts transmitted from different transmitters. The two data bursts 22, 24 contain the communication burst's data and are typically the same symbol length.

The antenna 40 of the receiver 28 receives various radio frequency signals. The received signals are demodulated by a demodulator 42 to produce a baseband signal. The baseband signal is processed, such as by a channel estimation device 44 and a joint detection device 46, in the time slots and with the appropriate codes assigned to the communication bursts of the corresponding transmitters 26. The channel estimation device 44 uses the training sequence component in the baseband signal to provide channel information, such as channel impulse responses. The channel information is used by the joint detection device 46 to estimate the transmitted data of the received communication bursts as soft symbols.

The joint detection device 46 uses the channel information provided by the channel estimation device 44 and the known spreading codes used by the transmitters 26 to estimate the data of the various received communication bursts. Although joint detection is described in conjunction with a TDD/CDMA communication system, the same approach is applicable to other communication systems, such as CDMA.

One approach to joint detection in a particular time slot in a TDD/CDMA communication system is illustrated in Figure 4. A number of communication bursts are superimposed on each other in the particular time slot, such as K communication bursts. The K bursts may be from K different transmitters. If certain transmitters are using multiple codes in the particular time slot, the K bursts may be from less than K transmitters.

Each data burst 22, 24 of the communication burst 16 has a predefined number of transmitted symbols, such as N_s . Each symbol is transmitted using a predetermined number of chips of the spreading code, which is the spreading factor (SF). In a typical TDD communication system, each base station 12₁ to 12₅ has an associated scrambling code mixed with its communicated data. The scrambling code distinguishes the base stations from one another. Typically, the scrambling code does not affect the spreading factor. Although the terms spreading code and factor are used hereafter, for systems using scrambling codes, the spreading code for the following is the combined scrambling and spreading codes. Each data burst 22, 24 has $N_s \times SF$ chips.

The joint detection device 46 estimates the value that each data burst symbol was originally transmitted. Equation 1 is used to determine the unknown transmitted symbols.

$$\underline{r} = A \underline{d} + \underline{n}$$

Equation 1

In Equation 1, the known received combined chips, $\underline{r}$, is a product of the system response, A, and the unknown transmitted symbols, $\underline{d}$. The term, $\underline{n}$, represents the noise in the wireless radio channel.

For K data bursts, the number of data burst symbols to be recovered is $N_s \times K$. For analysis purposes, the unknown data burst symbols are arranged into a column matrix, $\underline{d}$. The $\underline{d}$ matrix has column blocks, $\underline{d}_1$ to $\underline{d}_{N_s}$, of unknown data symbols. Each data symbol block, $\underline{d}_i$, has the i^{th} unknown transmitted data symbol in each of the K data bursts. As a result, each column block, $\underline{d}_i$, has K unknown transmitted symbols stacked on top of each other. The blocks are also stacked in a column on top of each other, such that $\underline{d}_1$ is on top of $\underline{d}_2$ and so on.

The joint detection device 46 receives a value for each chip as received. Each received chip is a composite of all K communication bursts. For analysis purposes, the composite chips are arranged into a column matrix, $\underline{r}$. The matrix $\underline{r}$ has a value of each composite chip, totaling $N_s \times \text{SF}$ chips.

A is the system response matrix. The system response matrix, A, is formed by convolving the impulse responses with each communication burst chip code. The convolved result is rearranged to form the system response matrix, A (step 48).

The joint detection device 46 receives the channel impulse response, $\underline{h}_i$, for each i^{th} one of the K communication bursts from the channel estimation device 44. Each $\underline{h}_i$ has a chip length of W. The joint detection device convolves the channel impulse responses with the known spreading codes of the K communication bursts to determine the symbol responses, $\underline{s}_1$ to $\underline{s}_K$, of the K communication bursts. A common support sub-block, S, which is common to all of the symbol responses is of length $K \times (\text{SF} + W - 1)$.

The A matrix is arranged to have N_s blocks, B_1 to B_{N_s} . Each block has all of the symbol responses, $\underline{s}_1$ to $\underline{s}_K$, arranged to be multiplied with the corresponding unknown data block in the $\underline{d}$ matrix, $\underline{d}_1$ to $\underline{d}_{N_s}$. For example, $\underline{d}_1$ is multiplied with B_1 . The symbol responses, $\underline{s}_1$ to $\underline{s}_K$, form a column in each block matrix, B_i , with the rest of the block being padded with zeros. In the first block, B_1 , the symbol response row starts at the first row. In the second block, the symbol response row is SF rows lower in the block and so on. As a result, each block has a width of K and a height of $N_s \times \text{SF}$. Equation 2 illustrates an A block matrix showing the block partitions.

$$A = \begin{bmatrix} \underline{s}_1 & \underline{s}_2 & \cdots & \underline{s}_K & 0 & 0 & \cdots & 0 & 0 & 0 & 0 & \cdots \\ 0 & 0 & 0 & 0 & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ & & & & 0 & 0 & 0 & 0 & & & & \\ & & & & \underline{s}_1 & \underline{s}_2 & \cdots & \underline{s}_K & & & & \\ \vdots & \vdots & \vdots & \vdots & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \cdots \\ & & & & \vdots & \vdots & \vdots & \vdots & & & & \\ & & & & & & & & \underline{s}_1 & \underline{s}_2 & \cdots & \underline{s}_K \\ & & & & & & & & 0 & 0 & 0 & 0 \\ & & & & & & & & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \cdots \end{bmatrix} = [B_1 \ B_2 \ \cdots \ B_{N_s}]$$

Equation 2

The $\underline{n}$ matrix has a noise value corresponding to each received combined chip, totaling $N_s \times SF$ chips. For analysis purposes, the $\underline{n}$ matrix is implicit in the received combined chip matrix, $\underline{r}$.

5 Using the block notation, Equation 1 can be rewritten as Equation 3.

$$\underline{r} = [B_1 \ B_2 \ B_3 \ \cdots \ B_{N_s}] \times \begin{bmatrix} \underline{d}_1 \\ \underline{d}_2 \\ \underline{d}_3 \\ \vdots \\ \underline{d}_{N_s} \end{bmatrix} + \underline{n} = \sum_{i=1}^{N_s} B_i \underline{d}_i + \underline{n}$$

Equation 3

Using a noisy version of the $\underline{r}$ matrix, the value for each unknown symbol can be determined by solving the equation. However, a brute force approach to solving Equation 1 requires extensive processing.

10 To reduce the processing, the system response matrix, A , is repartitioned. Each block, B_i , is divided into N_s blocks having a width of K and a height of SF .

These new blocks are referred to as A_1 to A_L and 0. L is the length of the common support S , as divided by the height of the new blocks, A_1 to A_L , per Equation 4.

$$L = \left\lceil \frac{SF + W - 1}{SF} \right\rceil$$

Equation 4

5 Blocks A_1 to A_L are determined by the supports, $\underline{s}_1$ to $\underline{s}_K$, and the common support, S . A 0 block is a block having all zeros. A repartitioned matrix for a system having a W of 57, SF of 16 and an L of 5 is shown in Equation 5.

$$A = \begin{bmatrix} A_1 & 0 & 0 & 0 & 0 & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ A_2 & A_1 & 0 & 0 & 0 & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ A_3 & A_2 & A_1 & 0 & 0 & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ A_4 & A_3 & A_2 & A_1 & 0 & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ A_5 & A_4 & A_3 & A_2 & A_1 & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & A_5 & A_4 & A_3 & A_2 & A_1 & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & 0 & A_5 & A_4 & A_3 & A_2 & A_1 & \ddots & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & A_5 & A_4 & A_3 & A_2 & \ddots & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & A_5 & A_4 & A_3 & \ddots & A_1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & A_5 & A_4 & \ddots & A_2 & A_1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & A_5 & \ddots & A_3 & A_2 & A_1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & \ddots & A_4 & A_3 & A_2 & A_1 \end{bmatrix}$$

Equation 5

10 To reduce the complexity of the matrix, a piecewise orthogonalization approach is used. Any of the blocks B_i for i being L or greater is non-orthogonal to any of the preceding L blocks and orthogonal to any blocks preceding by more than L . Each 0 in the repartitioned A matrix is an all zero block. As a result to use a piecewise orthogonalization, the A matrix is expanded (step 50).

The A matrix is expanded by padding $L-1$ zero blocks to the right of each block of the A matrix and shifting each row in the A matrix by its row number less

one. To illustrate for the A1 block in row 2 of Figure 2, four (L-1) zeros are inserted between A2 and A1 in row 2. Additionally, block A1 (as well as A2) is shifted to the right by one column (row 2 -1). As a result, Equation 5 after expansion would become Equation 6.

$$A_{\text{exp}} = \begin{bmatrix} A_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & A_2 & 0 & 0 & 0 & A_1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & A_3 & 0 & 0 & 0 & A_2 & 0 & 0 & 0 & A_1 \\ 0 & 0 & 0 & A_4 & 0 & 0 & 0 & A_3 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & A_5 & 0 & 0 & 0 & A_4 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & A_5 & 0 & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

5

Equation 6

To accommodate the expanded A matrix, the $\underline{d}$ matrix must also be expanded, $\underline{d}_{\text{exp}}$. Each block, $\underline{d}_1$ to $\underline{d}_{Ns}$, is expanded to a new block, $\underline{d}_{\text{exp}1}$ to $\underline{d}_{\text{exp}Ns}$. Each expanded block, $\underline{d}_{\text{exp}1}$ to $\underline{d}_{\text{exp}Ns}$, is formed by repeating the original block L times. For example for $\underline{d}_{\text{exp}1}$, a first block row would be created having L versions of $\underline{d}_1$, stacked one below the other.

10

As a result, Equation 1 can be rewritten as Equation 7.

$$\underline{r} = A_{\text{exp}} \cdot \underline{d}_{\text{exp}} + \underline{n}$$

$$= \begin{bmatrix} B_{\text{exp}1} & B_{\text{exp}2} & B_{\text{exp}3} & \cdots & B_{\text{exp}N_s} \end{bmatrix} \times \begin{bmatrix} \underline{d_{\text{exp}1}} \\ \underline{d_{\text{exp}2}} \\ \underline{d_{\text{exp}3}} \\ \vdots \\ \underline{d_{\text{exp}N_s}} \end{bmatrix} + \underline{n} = \sum_{i=1}^{N_s} B_{\text{exp}i} \underline{d_{\text{exp}i}} + \underline{n},$$

Equation 7

Equation 7 can be rewritten to partition each $B_{\text{exp}i}$ orthogonally in L partitions, $U_j^{(i)}$, $j = 1$ to L , as in Equation 8.

$$\begin{aligned} \underline{r} &= A_{\text{exp}} \cdot \underline{d_{\text{exp}}} + \underline{n} \\ &= \sum_{i=1}^{N_s} \begin{bmatrix} U_1^{(i)} & U_2^{(i)} & \cdots & U_L^{(i)} \end{bmatrix} \times \begin{bmatrix} \underline{d_i} \\ \underline{d_i} \\ \underline{d_i} \\ \vdots \\ \underline{d_i} \end{bmatrix} + \underline{n} = \sum_{i=1}^{N_s} \sum_{j=1}^L U_j^{(i)} \underline{d_i} = \sum_{i=1}^{N_s} B_i \underline{d_i} + \underline{n} \end{aligned}$$

Equation 8

5 To reduce computational complexity, a QR decomposition of the A_{exp} matrix is performed (step 52). Equation 9 illustrates the QR decomposition of A_{exp} .

$$A_{\text{exp}} = Q_{\text{exp}} R_{\text{exp}}$$

Equation 9

10 Due to the orthogonal partitioning of A_{exp} , the QR decomposition of A_{exp} is less complex. The resulting Q_{exp} and R_{exp} matrices are periodic with an initial transient extending over L blocks. Accordingly, Q_{exp} and R_{exp} can be determined by calculating the initial transient and one period of the periodic portion. Furthermore,

the periodic portion of the matrices is effectively determined by orthogonalizing A_1 to A_L . One approach to QR decomposition is a Gramm-Schmidt orthogonalization.

To orthogonalize A_{exp} as in Equation 6, $B_{\text{exp}1}$ is orthogonalized by independently orthogonalizing each of its orthogonal partitions, $\{U_j^{(i)}\}, j = 1 \cdots L$.

5 Each $\{A_j\}, j = 1 \cdots L$ is independently orthogonalized, and the set is zero-padded appropriately. $\{Q_j\}$ are the orthonormal sets obtained by orthogonalizing $\{U_j^{(i)}\}$.

To determine $B_{\text{exp}2}$, its $U_1^{(2)}$ needs to be orthogonalized with respect to only Q_2 of $B_{\text{exp}1}$ formed previously. $U_2^{(2)}, U_3^{(2)}$ and $U_4^{(2)}$ only need to be orthogonalized with respect to only Q_3, Q_4 and Q_5 , respectively. $U_5^{(2)}$ needs to be orthogonalized to all
10 previous Q s and its orthogonalized result is simply a shifted version of Q_5 obtained from orthogonalizing $B_{\text{exp}1}$.

As the orthogonalizing continues, beyond the initial transient, there emerges a periodicity which can be summarized as follows. The result of orthogonalizing $B_{\text{exp}i}, i \geq 6$ can be obtained simply by a periodic extension of the result of
15 orthogonalizing $B_{\text{exp}5}$.

The orthogonalization of $B_{\text{exp}5}$, is accomplished as follows. Its Q_5 is obtained by orthogonalizing A_5 , and then zero padding. Its Q_4 is obtained by orthogonalizing the support of Q_5 and A_4 , $[\text{sup}(Q_5) A_4]$, and then zero padding. Since $\text{sup}(Q_5)$ is already an orthogonal set, only A_4 needs to be orthogonalized with respect to $\text{sup}(Q_5)$ and itself. Its Q_3 is obtained by orthogonalizing $[\text{sup}(Q_5) \text{sup}(Q_4) A_3]$ and then zero
20 padding. Its Q_2 is obtained by orthogonalizing $[\text{sup}(Q_5) \text{sup}(Q_4) \text{sup}(Q_3) A_2]$ and then zero padding. Its Q_1 is obtained by orthogonalizing $[\text{sup}(Q_5) \text{sup}(Q_4) \text{sup}(Q_3) \text{sup}(Q_2) A_1]$ and then zero padding. Apart from the initial transient, the entire A_{exp} can be efficiently orthogonalized, by just orthogonalizing A_p per Equation 10.

$$A_p = [A_5 \ A_4 \ A_3 \ A_2 \ A_1]$$

Equation 10

By effectively orthogonalizing the periodic portion of A_{exp} by using only A_p , computational efficiency is achieved. Using a more compact notation, Q_i^s , for $\text{sup}(Q_i)$, this orthogonalization of A_p results in the orthonormal matrix, Q_p , of Equation 11.

$$Q_p = [Q_5^s \quad Q_4^s \quad Q_3^s \quad Q_2^s \quad Q_1^s]$$

Equation 11

The periodic part of Q_{exp} is per Equation 12.

$$\text{PeriodicPartof}Q_{\text{exp}} = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \\ Q_1^s & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & Q_2^s & 0 & 0 & 0 & Q_1^s & 0 & 0 & 0 & 0 & \\ 0 & 0 & Q_3^s & 0 & 0 & 0 & Q_2^s & 0 & 0 & 0 & \\ 0 & 0 & 0 & Q_4^s & 0 & 0 & 0 & Q_3^s & 0 & 0 & \\ 0 & 0 & 0 & 0 & Q_5^s & 0 & 0 & 0 & Q_4^s & 0 & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & Q_5^s & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \end{bmatrix}$$

Equation 12

To constructing the upper triangular matrix $R_{\text{exp}}, \langle A_i \rangle_j$, is a block of size $K \times$

10 K representing the projections of each column of A_i onto all the columns of Q_j^s . For example, the first column of $\langle A_4 \rangle_5$ represents the projections of the first column of A_4 on each of the K columns of Q_5^s . Similarly, $\langle A_4 \rangle_4$ represents the projections of the first column of A_4 on each of the K columns of Q_4^s . However, this block will be

upper triangular, because the k^{th} column of A_4 belongs to the space spanned by the orthonormal vectors of Q_5^S and the first k vectors of Q_4^S . This block is also orthogonal to subsequent vectors in Q_4^S , leading to an upper triangular $\langle A_4 \rangle_4$. Any $\langle A_i \rangle_j$ with $i = j$ will be upper triangular. To orthogonalize other blocks, the following results.

The first block of $B_{\text{exp}5}$, viz., $U_1^{(5)}$ is formed by a linear combination of $\{Q_j^S\}$, $j = 1 \cdots 5$, with coefficients given by $\langle A_1 \rangle_j$, $j = 1 \cdots 5$. The second block, $U_2^{(5)}$, is formed by a linear combination of $\{Q_j^S\}$, $j = 2 \cdots 5$, with coefficients given by $\langle A_2 \rangle_j$, $j = 2 \cdots 5$. The third block, $U_3^{(5)}$, is formed by a linear combination of $\{Q_j^S\}$, $j = 3 \cdots 5$, with coefficients given by $\langle A_2 \rangle_j$, $j = 3 \cdots 5$. The fourth block, $U_4^{(5)}$, is formed by a linear combination of $\{Q_j^S\}$, $j = 4, 5$, with coefficients given by $\langle A_2 \rangle_j$, $j = 4, 5$. The fifth block, $U_5^{(5)}$, is formed by $Q_5^S \times \langle A_5 \rangle_5$.

Accordingly, the coefficients in the expansion of subsequent $B_{\text{exp}i}$, $i \geq 6$ are simply periodic extensions of the above. Since the R_{exp} entries are computed during the orthogonalization of A_{exp} , no additional computations are needed to construct R_{exp} . Disregarding the initial transient, the remainder of R_{exp} is periodic, and two periods of it are shown in Equation 13.

$$R_{\text{exp}} = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \\ 0 & & & & & & & & & & \\ \langle A_1 \rangle_5 & & & & & & & & & & \\ 0 & & & & & & & & & & \\ 0 & & & & & & & & & & \\ \langle A_1 \rangle_4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ 0 & \langle A_2 \rangle_5 & 0 & 0 & 0 & \langle A_1 \rangle_5 & 0 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ \langle A_1 \rangle_3 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ 0 & \langle A_2 \rangle_4 & 0 & 0 & 0 & \langle A_1 \rangle_4 & 0 & 0 & 0 & & \\ 0 & 0 & \langle A_3 \rangle_5 & 0 & 0 & 0 & \langle A_2 \rangle_5 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ \langle A_1 \rangle_2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ 0 & 0 & \langle A_2 \rangle_3 & 0 & 0 & \langle A_1 \rangle_3 & 0 & 0 & 0 & & \\ 0 & 0 & 0 & \langle A_3 \rangle_4 & 0 & 0 & \langle A_2 \rangle_4 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & \langle A_4 \rangle_5 & 0 & 0 & 0 & \langle A_3 \rangle_5 & 0 & \dots \\ \langle A_1 \rangle_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & & \\ 0 & \langle A_2 \rangle_2 & 0 & 0 & 0 & \langle A_1 \rangle_2 & 0 & 0 & 0 & & \\ 0 & 0 & \langle A_3 \rangle_3 & 0 & 0 & 0 & \langle A_2 \rangle_3 & 0 & 0 & & \\ 0 & 0 & 0 & \langle A_4 \rangle_4 & 0 & 0 & 0 & \langle A_3 \rangle_4 & 0 & & \\ 0 & 0 & 0 & 0 & \langle A_5 \rangle_5 & 0 & 0 & 0 & \langle A_4 \rangle_5 & & \\ 0 & 0 & 0 & 0 & 0 & \langle A_1 \rangle_1 & 0 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & \langle A_2 \rangle_2 & 0 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & \langle A_3 \rangle_3 & 0 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \langle A_4 \rangle_4 & & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \langle A_5 \rangle_5 & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \langle A_1 \rangle_1 \end{bmatrix}$$

Equation 13

The least squares approach to solving Q_{exp} and R_{exp} is shown in Equation 14.

$$Q_{\text{exp}} \cdot R_{\text{exp}} \cdot \underline{d_{\text{exp}}} = \underline{r}$$

Equation 14

By pre-multiplying both sides of Equation 14 by the transpose of Q_{exp} , Q_{exp}^T , and using $Q_{\text{exp}}^T \cdot Q_{\text{exp}} = I_{LKN_s}$, Equation 14 becomes Equation 15.

$$R_{\text{exp}} \cdot \underline{d_{\text{exp}}} = Q_{\text{exp}}^T \underline{r}$$

Equation 15

Equation 15 represents a triangular system whose solution also solves the LS problem of Equation 14.

Due to the expansion, the number of unknowns is increased by a factor of L. Since the unknowns are repeated by a factor of L, to reduce the complexity, the repeated unknowns can be collected to collapse the system. R_{exp} is collapsed using L coefficient blocks, CF_1 to CF_L , each having a width and a height of K. For a system having an L of 5, CF_1 to CF_5 can be determined as in Equation 16.

$$CF_1 = \langle A_1 \rangle_1 + \langle A_2 \rangle_2 + \langle A_3 \rangle_3 + \langle A_4 \rangle_4 + \langle A_5 \rangle_5$$

$$CF_2 = \langle A_1 \rangle_2 + \langle A_2 \rangle_3 + \langle A_3 \rangle_4 + \langle A_4 \rangle_5$$

$$CF_3 = \langle A_1 \rangle_3 + \langle A_2 \rangle_4 + \langle A_3 \rangle_5$$

$$CF_4 = \langle A_1 \rangle_4 + \langle A_2 \rangle_5$$

$$CF_5 = \langle A_1 \rangle_5$$

Equation 16

Collapsing R_{exp} using the coefficient blocks produces a Cholesky like factor, $\hat{G}$ (step 54). By performing analogous operations on the right hand side of Equation 15 results in a banded upper triangular system of height and width of K x N_s as in Equation 17.

$$\begin{bmatrix} Tr_1 & Tr_2 & Tr_3 & Tr_4 & CF_5 & 0 & 0 & 0 & 0 & 0 & \dots \\ 0 & Tr_1 & Tr_2 & CF_3 & CF_4 & CF_5 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & CF_1 & CF_2 & CF_3 & CF_4 & CF_5 & 0 & 0 & 0 & \dots \\ 0 & 0 & 0 & CF_1 & CF_2 & CF_3 & CF_4 & CF_5 & 0 & 0 & \dots \\ 0 & 0 & 0 & 0 & CF_1 & CF_2 & CF_3 & CF_4 & CF_5 & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \ddots & 0 \end{bmatrix} \times \begin{bmatrix} \underline{d_1} \\ \underline{d_2} \\ \underline{d_3} \\ \vdots \\ \underline{d_{N_s}} \end{bmatrix} = \underline{\hat{r}}$$

Equation 17

Tr_1 to Tr_4 are the transient terms and $\hat{r}$. By solving the upper triangle via back substitution, Equation 17 can be solved to determine $\underline{d}$ (step 56). As a result, the transmitted data symbols of the K data bursts is determined.

- 5 Using the piecewise orthogonalization and QR decomposition, the complexity of solving the least square problem when compared with a banded Cholesky decomposition is reduced by a factor of 6.5.

* * *

CLAIMS

What is claimed is:

1. A method for use at a receiver of retrieving data from a plurality of data signals transmitted in a communication system, the method comprising: receiving the plurality of transmitted data signals at the receiver and measuring a channel response associated with the transmitted data signals; determining a system response based on in part the channel response (48); the method characterized by:
 - expanding the system response to be piecewise orthogonal (50); and
 - retrieving data from the received data signals based on in part the expanded system response (52, 54, 56).
2. The method of claim 1 wherein each of the transmitted data signals has an associated chip code and the system response is determined by convolving the associated chip codes with the channel response.
3. The method of claim 1 wherein the system response is a system response matrix, further comprising dividing the system response matrix into blocks of columns prior to the expanding.
4. The method of claim 3 wherein the expanding is by padding zeros in the column blocks to such that each block is orthogonal.
5. The method of claim 3 wherein SF is a spreading factor associated with the data signals and W is a chip length associated with the channel response and the blocks contain L columns, where L is
$$\left[L = \frac{SF + W - 1}{SF} \right].$$
6. The method of claim 1 further comprising QR decomposing the expanded system response (52).

7. The method of claim 6 wherein the QR decomposing is by Gram-Schmitt orthogonalization.

8. The method of claim 3 further comprising replacing elements of a column block with sub-blocks prior to the expanding.

9. The method of claim 8 wherein in the sub-blocks have a same width as the column blocks and have a height associated with a length of a symbol response and a support of the symbol response.

10. The method of claim 8 further comprising QR decomposing the system response matrix and orthogonalizing a periodic portion of a Q matrix is by orthogonalizing the sub-blocks.

11. The method of claim 10 further comprising orthogonalizing a periodic portion of an R matrix based on in part the orthogonalized sub-blocks associated with the Q matrix.

12. The method of claim 6 wherein the QR decomposing results in a Q matrix and an R matrix, each having an initial transient and a periodic portion.

13. The method of claim 12 further comprising collapsing the R matrix by replacing elements in the R matrix with coefficient blocks.

14. The method of claim 13 wherein the R matrix with coefficient blocks is a Cholesky-like factor (54).

15. A joint detection device (46) for use in a receiver receiving a plurality of data signals, the joint detection device (46) comprising: means (46) for

determining a system response based on in part a measured channel response; the joint detection device characterized by:

- 5 means (46) for expanding the system response to be piecewise orthogonal; and
 means (46) for retrieving data from the received data signals based on in part the expanded system response.

16. A receiver for receiving a plurality of transmitted data signals in a communication system, the receiver comprising: an antenna (40) for receiving the transmitted data signals; a channel estimation device for determining a channel response for each received data signal (44); the joint detection device characterized by:

- 5 a joint detection device (46) having an input configured to receive the channel responses and the received data signals for determining a system response based on in part the channel signals, expanding the system response to be piecewise orthogonal, and retrieving data from the received data signals based on in part the expanded system response.
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17. The receiver of claim 16 for use in a time division duplex using code division multiple access communication system.

18. The receiver of claim 17 wherein each of the transmitted data signals has an associate code and is transmitted in a shared frequency spectrum and the system response is determined by convolving the associated chip codes with the channel response.

19. The receiver of claim 17 wherein the channel estimation device measures the channel response using a received training sequence associated with the data signals.

20. The receiver of claim 16 wherein the system response is a system response matrix, further comprising dividing the system response matrix into blocks of columns prior to the expanding.

21. The receiver of claim 20 wherein the expanding is by padding zeros in the column blocks such that each column block is orthogonal.

22. The receiver of claim 20 wherein SF is a spreading factor associated with the data signals and W is a chip length associated with the channel response and the blocks contain L columns, where L is $\left\lceil L = \frac{SF + W - 1}{SF} \right\rceil$.