# METHOD FOR GENERATING INTERPRETABLE CODE FOR STORAGE IN A DEVICE HAVING LIMITED STORAGE

## BACKGROUND OF THE INVENTION

[0001] Java source code files (*.java* files) are compiled into *.class* files by the Java compiler. These *.class* files may be read into memory by the Java Virtual Machine (VM) in an in-memory format that is suitable for interpretation by the VM. The *.class* files are then linked with other *.class* files that have been similarly read. References between *.class* files are resolved symbolically, using character strings. These character strings appear in both the *.class* file that contains the definition and in the *.class* file that references the definition. Therefore, the presence of references between *.class* files may increase the size of the *.class* files.

[0002] Java *.class* files may be archived (and optionally compressed) into a *.jar* file. However, *.jar* files are not directly interpretable by the Java VM, and the *.class* files must be extracted (and decompressed, if applicable) from the *.jar* file (and read into memory) in order for them to be linked, resolved and interpreted by the Java VM.

[0003] Although *.jar* files comprising archived and compressed *.class* files are smaller than the *.class* files themselves (and are therefore more suitable for transmission between communication devices), storage space for the extracted (and decompressed, if applicable) *.class* files needs to be available in the environment where the application is to be executed, so that the Java VM may access the *.class* files. Consequently, a solution involving *.jar* files may not represent a savings in storage space. In fact, a solution involving *.jar* files may require extra storage space, for both the *.jar* files and the extracted *.class* files. The extracted *.class* files need not be retained once they have been loaded into memory, linked and resolved. However, both the *.jar* file and the in-memory representations of the *.class* files must be retained. In an environment having limited storage, where storage space is at a premium, it may therefore be preferable to store only the *.class* files and not to use a solution involving *.jar* files. However, as explained above, the size of *.class* files increases when the number of references between *.class* files increases.

[0004] Therefore, it would be beneficial to generate directly interpretable files that are of a smaller size than *.class* files, while providing a solution for references between *.class* files.

## SUMMARY OF THE INVENTION

[0004a] In one aspect, an embodiment of the invention provides a device comprising: a memory unit including executable software; a plurality of class files stored in the memory unit; and, a computing unit connected to the memory unit, the computing unit being able to execute a Java Virtual Machine. The computing unit executes the executable software for generating a number of cod files from the plurality of class files by combining elements from the plurality of class files without duplication of entries for reducing storage space, wherein the number of cod files is less than the number of class files. A given cod file includes a constant pool created by combining constant pool entries from two or more of the class files without duplication of entries; a byte codes and information structure created by combining byte codes and information structure entries from the two or more of the class files without duplication of entries; and, a fixup table for providing information to the Java Virtual Machine for resolving at least one entry in the given cod file at link time.

[0004b] In another aspect, an embodiment of the invention provides a method for generating a number of cod files from a plurality of class files by combining elements from the plurality of class files such that the number of cod files is less than the number of class files without duplication of entries for reducing storage space. For a given cod file, the method comprises:

identifying class files with common entries in at least one of the constant pool and the byte codes and information structure;

generating a constant pool for the given cod file by combining constant pool entries from the class files with common entries without duplication;

generating the byte codes and information structure for the given cod file by combining byte codes and information structure entries from the class files with common entries without duplication; and,

generating a fixup table for providing information to a Java Virtual Machine for resolving at least one entry in the given cod file at link time.

[0004c] In yet another aspect, an embodiment of the invention provides an article storing executable software that when executed by a computing unit generates a number of cod files from a plurality of class files by combining elements from the plurality of class files without duplication of entries for reducing storage space. The executable software comprises code for generating a given cod file to include a constant pool created by combining constant pool entries from two or more of the class files without duplication of entries; a byte codes and information structure created by combining byte codes and information structure entries from the two or more of the class files without duplication of entries; and, a fixup table for providing information to a Java Virtual Machine for resolving at least one component of the given cod file at link time.

[0004d] In each of these aspects, the given cod file may further include a sibling list for listing other related cod files to define a sibling group and the fixup table of the given cod file includes indices to the other related cod files specified in the sibling group.

## BRIEF DESCRIPTION OF THE DRAWINGS

[0005] The subject matter regarded as the invention is particularly pointed out and distinctly claimed in the concluding portion of the specification. The invention, however, both as to organization and method of operation, together with objects, features and advantages thereof, may best be understood by reference to the following detailed description when read with the accompanied drawings in which:

[0006] FIGS. 1A and 1B are simplified prior art illustrations of six exemplary *.class* files;

[0007] FIGS. 2A and 2B are simplified illustrations of exemplary *.cod* files, according to some embodiments of the present invention;

[0008] FIG. 3 is a flowchart illustration of a method for generating *.cod* files, according to some embodiments of the present invention; and

[0009] FIG. 4 is a simplified block-diagram illustration of a device having a computing unit and directly addressable memory, according to some embodiments of the present invention.

[0010] It will be appreciated that for simplicity and clarity of illustration, elements shown in the figures have not necessarily been drawn to scale. For example, the dimensions of some of the elements may be exaggerated relative to other elements for clarity. Further, where considered appropriate, reference numerals may be repeated among the figures to indicate corresponding or analogous elements.

## DETAILED DESCRIPTION OF THE INVENTION

[0011] In the following detailed description, numerous specific details are set forth in order to provide a thorough understanding of the invention. However it will be understood by those of ordinary skill in the art that the present invention may be practiced without these specific details. In other instances, well-known methods, procedures and components have not been described in detail so as not to obscure the present invention.

[0012] Java source code files (*.java* files) are compiled into *.class* files by the Java compiler. These *.class* files, once read into memory, may be linked, resolved and interpreted by the Java Virtual Machine (VM). FIGS. 1A and 1B are simplified prior art illustrations of six exemplary *.class* files – *Aa.class* file 100A, *Bb.class* file 100B, *Cc.class* file 100C, *Dd.class* file 100D, *Ee.class* file 100E, and *Ff.class* file 100F, collectively referred to as file 100.

[0013] The structure of the *.class* file is well-documented and will not be described here in detail. File 100 is illustrated in a simplified format for clarity of explanation, however it will be understood by persons of ordinary skill in the art that this simplified format is not an exact representation of the structure of file 100. Each file 100 comprises a constant pool 102. Constant pool 102 comprises indexed *cp_info* entries, some of which comprise character strings indicating the names of the class and parent class, the names of methods, the type of methods, the names of fields, the type of fields, etc. referred to within the *ClassFile* structure. Each file 100 also comprises "byte codes and information structures" 104 regarding the class properties, the methods, fields and attributes of the class, and their types. These structures may point to entries in constant pool 102 using ordinal index numbers.

[0014] In *Aa.class* file 100A, Method G includes a reference to another method, Method K, found in class Bb. In fact, Method G is defined in the *.class* file as a *method_info* structure comprising, among other things, indices of entries in constant pool 102A. However, for clarity of description, the definition of Method G and its reference to Method K are shown in FIG. 1A in a simplified format. Therefore, the character strings "K" and "Bb" are included in constant pool 102A. The inclusion of "K" and "Bb" in constant pool 102 A actually appear as part of several *cp_info* structures. However, for clarity of description, only the character strings "K" and

"Bb" are shown in FIG. 1A. Since the argument of Method K is of type integer, the *BaseType* character "I" is included in constant pool 102A. The symbolic reference to Method K by Method G in the byte codes and information structures 104A comprises the indices of the relevant *cp_info* structures in constant pool 102A.

[0015] Similarly, in *Bb.class* file 100B, Method K is defined in class Bb, so constant pool 102B includes all the strings "Bb", "K" and the *BaseType* character "I" (representing the 'int' type).

[0016] FIGS. 2A and 2B are simplified illustrations of exemplary *.cod* files, the generation of which will be described hereinbelow with respect to FIG. 3. *.cod* files are directly interpretable files that comprise the information of *.class* files in a manner that requires less storage space than the *.class* files themselves. *.cod* files may be "solo" *.cod* files or "sibling" *.cod* files, as will be explained in greater detail hereinbelow.

[0017] In the example shown in FIG. 2B, solo *.cod* file 220 is named "EeFf.cod", to signify that it results from combining the files *Ee.class* and *Ff.class*, although the scope of the present invention is not limited in this respect. Solo *EeFf.cod* file 220 comprises a constant pool 222, one or more fixup tables 228, and "byte codes and information structures" 224 regarding the class properties, the methods, fields and attributes of the class, and their types.

[0018] FIG. 3 is a simplified flowchart illustration of a method for the generation of *.cod* files from a collection of *.java* files or *.class* files. The method may be performed by executable software on a general-purpose computer, the software being provided, for example, as part of a software developer's kit (SDK).

[0019] A list of *.java* files or *.class* files is received (step 300). If the input is *.java* files, then a Java compiler is applied to the *.java* files to produce corresponding *.class* files (step 302). Alternatively, although this is not shown in FIG. 3, one or more *.jar* files could be received, and information regarding the archived *.class* files could be extracted therefrom.

[0020] The executable software identifies common entries in the constant pools of the *.class* files (step 304). For example, in FIG. 1B, the parent class name character string "java/lang/Object", the method name character string "T" and the *BaseType* character

"C" (representing the 'char' type) are common to *Ee.class* file 100E and to *Ff.class* file 100F.

[0021] The executable software identifies cross-references between classes, methods, and fields in the *.class* files (step 306). For example, in FIG. 1B, Method Ee.T is referenced by Method Ff.U.

[0022] The executable software then generates the *.cod* file by combining elements of the *.class* files (step 308). Unlike standard compression algorithms that compress data regardless of its content, the combination of elements of the *.class* files into the *.cod* file is performed differently based upon the type of element and its content. For example, the *.cod* file comprises a constant pool whose entries include the information stored in the constant pools of the *.class* files, but without duplications of redundant entries. In the generated *.cod* file, the constant pool contains only a single copy of the common entries identified in step 304. For example, in FIG. 2B, constant pool 222 of solo *EeFf.cod* file comprises only one instance of character strings "java/lang/Object" and "T" and one instance of *BaseType* character "C".

[0023] The executable software uses hard offsets in the generated solo *.cod* file for cross-references between classes and methods defined in the *.class* files being combined. For example, in FIG. 2B, the reference to Method Ee.T by Method Ff.U in the "byte codes and information structures" 224 comprises a hard offset, $HO_{Ee.T}$, specifying the location of the definition of Method Ee.T within solo *EeFf.cod* file 220. This hard offset does not need to be resolved or put into context by the Java VM at link time.

[0024] The above-described use of a hard offset in a *.cod* file may be contrasted with the use of offsets in Windows™ .DLL files. References in Windows™ .DLL files may be in terms of a symbolic name or an ordinal number. The ordinal number is used as an index into a table of offsets. Embodiments of the present invention use the hard offsets directly in the *.cod* file, giving a more compact representation. Moreover, Windows™ .DLL files are used for compiled code, while the *.cod* files are used for interpreted Java.

[0025] The executable software uses symbolic references in the generated solo *.cod* file for cross-references between the classes received in step 300 and other classes. In the present example, class Ee and class Ff extend java.lang.Object. Therefore,

constant pools 102E, 102F and 222 each comprise a single instance of class name "java/lang/Object" so that definitions of the classes may refer to this parent class. In solo *EeFf.cod* file 220, constant pool 222 comprises the string "java/lang/Object", and the reference to the java.lang.Object class in the definitions of the classes in "byte codes and information structures" 224 is a symbolic reference (using the index of the string in constant pool 222) that needs to be resolved by the Java VM at link time using information stored in fixup table 228.

[0026] The executable software may perform additional actions that have not been illustrated in FIG. 3. For example, other space-saving techniques may be used to further reduce the size of the solo *.cod* files.

[0027] It will be appreciated by persons of ordinary skill in the art that when either of the source code files *Ee.java* and *Ff.java* are modified, it is necessary to invoke the executable software to generate a new *EeFf.cod* file, or alternatively to invoke the executable software on the modified *.java* or *.class* files along with one or more additional *.java* or *.class* files to generate a new *EeFfxxx.cod* file, where "x" indicates the additional classes.

[0028] It will also be appreciated by persons of ordinary skill in the art that if it were possible to generate a single solo *.cod* file for all the *.class* files to be used in a particular application, then the Java VM would not need to resolve symbolic references between classes, methods and fields, since all such references would appear in the *.cod* file as hard offsets. Moreover, such a *.cod* file might be significantly smaller than the total size of the *.class* files, since duplications of information in the *.class* files would be eliminated in the *.cod* file. Such a single *.cod* file would also be smaller than the multiple solo *.cod* files due to the reduction in duplicated information.

[0029] However, unlike Windows™ .DLL files, which are relatively unlimited in size, there are sometimes significant limitations on the size of a single *.cod* file. For example, FIG. 4 shows a device 400 comprising a computing unit 402 able to execute a Java VM and a directly addressable memory 404 having one or more *.cod* files 406 stored therein. Although the scope of the present invention is not limited in this respect, computing unit 402 may be a general-purpose microprocessor or any computing unit capable of executing a Java VM, and device 400 may be a digital

camera, a handheld video recorder, a cellular telephone, a personal digital assistant, a handheld computer, an MP3 player, a CD player, a handheld gaming device, a pager, a two-way pager, etc. Although the scope of the present invention is not limited in this respect, directly addressable memory 404 may be not-OR (NOR) flash memory whose physical layout imposes an upper limit of 64 kiloBytes for the size of each *.cod* file stored therein.

[0030] If an application were to be stored exclusively as solo *.cod* files in a storage medium that imposes a limit on the size of individual *.cod* files, then each *.cod* file would likely comprise many symbolic references to account for the cross-references between classes in that *.cod* file and classes in other solo *.cod* files. Since symbolic references require more storage space in the *.cod* file than hard offsets, the overall size of the solo *.cod* representation of the application may be quite large.

[0031] As an alternative to the exclusive use of solo *.cod* files in the representation of an application, sibling *.cod* files may be used when the *.cod* files are to be stored in a storage medium that imposes a limit on the size of individual *.cod* files. The software developer may group together *.java* files or *.class* files into sibling groups. This grouping may be based on the software developer's knowledge of the intricacy of cross-references between the classes represented in the files. A *.cod* file may be generated from these *.java* files or *.class* files, and if, after a portion of the *.java* files or *.class* files have been combined into a *.cod* file, the size of the *.cod* file would exceed a predetermined limit if another *.java* or *.class* file were to be combined into the *.cod* file, then one or more sibling *.cod* files are created for the remaining *.java* files or *.class* files in the sibling group. Classes are not split across *.cod* file boundaries.

[0032] It will be appreciated by persons of ordinary skill in the art that for a given set of *.class* files, packing the *.class* files into two sibling *.cod* files will be a more compact representation than packing the *.class* files into three sibling *.cod* files. Therefore, it is desirable to pack the *.class* files in the group into as few *.cod* files as possible, while maintaining the constraint of individual *.cod* file size. The problem of how to divide a set of cross-referencing classes among sibling *.cod* files is similar to the well-known "packing problem" in the art of software engineering. Various

techniques to minimize the overall size of the *.cod* representation of an application may be used.

[0033] FIG. 2A is a simplified illustration of two exemplary sibling *.cod* files according to some embodiments of the present invention. A sibling *.cod* file 200 is named "AaBb.cod" to signify that it results from combining the files *Aa.class* and *Bb.class*, although the scope of the present invention is not limited in this respect. Similarly, a sibling *.cod* file 210 is named "CcDd.cod" to signify that it results from combining the files *Cc.class* and *Dd.class*.

[0034] Each sibling *.cod* file comprises a list of its siblings. For example, a sibling list 206 of *AaBb.cod* file 200 comprises "CcDd" to indicate that *CcDd.cod* file 210 is a sibling of *AaBb.cod* file 200. Similarly, a sibling list 216 of *CcDd.cod* file 210 comprises "AaBb" to indicate that *AaBb.cod* file 200 is a sibling of *CcDd.cod* file 210.

[0035] Sibling *AaBb.cod* file 200 also comprises a constant pool 202, "byte codes and information structures" 204, and one or more fixup tables 208. Similarly, sibling *CcDd.cod* file 210 also comprises a constant pool 212, "byte codes and information structures" 214, and one or more fixup tables 218. The fixup tables 208, 218 may include indications of the places in the *.cod* file where resolution work is to be done by the Java VM at link time. The fixup tables may include pointers to the name of another *.cod* file, if necessary to the name of the class containing the symbol, if necessary to the name of the method or field within that class being referenced, and to method or field typing information.

[0036] FIG. 3 is a simplified flowchart illustration of a method for the generation of *.cod* files from a collection of *.java* files or *.class* files. The method may be performed by executable software on a general-purpose computer, the software being provided, for example, as part of a software developer's kit (SDK).

[0037] A list of *.java* files or *.class* files is received (step 300). If the input is *.java* files, then a Java compiler is applied to the *.java* files to produce corresponding *.class* files (step 302). Alternatively, although this is not shown in FIG. 3, one or more *.jar* files could be received, and information regarding the archived *.class* files could be extracted therefrom.

[0038] The executable software identifies common entries in the constant pools of the *.class* files (step 304).

[0039] The executable software identifies cross-references between classes, methods, and fields in the *.class* files (step 306).

[0040] The executable software then generates sibling *.cod* files by combining elements of the *.class* files (step 308). In the generated sibling *.cod* file, the constant pool contains only a single copy of the common entries identified in step 304. For example, in FIG. 2A, constant pool 202 of sibling *AaBb.cod* file 200 comprises only one instance of character strings "Bb" and "K" and one instance of *BaseType* character "I". In another example, constant pool 212 of sibling *CcDd.cod* file 210 comprises only one instance of character strings "Cc" and "P" and one instance of *BaseType* character "I".

[0041] The executable software uses hard offsets in the generated sibling *.cod* file for cross-references between classes whose *.cod* files are in the same sibling group. The executable software uses symbolic references in the generated sibling *.cod* file for cross-references between classes whose *.cod* files belong to different sibling groups, and between classes having solo *.cod* files.

[0042] This is better understood by way of examples:

1) As shown in FIG. 1A, Method Aa.G comprises a reference to Method Bb.K. Therefore, as shown in FIG. 2A, this reference appears in the definition of Method Aa.G in "byte codes and information structures" 204 of sibling *AaBb.cod* file 200 as a hard offset, $HO_{Bb.K}$, specifying the location of the definition of Method Bb.K within sibling *AaBb.cod* file 200. This hard offset does not need to be resolved or put into context by the Java VM at link time.

2) Similarly, as shown in FIG. 1B, Method Dd.Q comprises a reference to Method Cc.P. Therefore, as shown in FIG. 2A, this reference appears in the definition of Method Dd.Q in "byte codes and information structures" 214 of sibling *CcDd.cod* file 210 as a hard offset, $HO_{Cc.P}$, specifying the location of the definition of Method Cc.P within sibling *CcDd.cod* file 210. This hard offset does not need to be resolved or put into context by the Java VM at link time.

3) As shown in FIG. 1A, Method Bb.L comprises a reference to Method Cc.N. Therefore, as shown in FIG. 2A, this reference appears in the definition of Method

Bb.L in "byte codes and information structures" 204 of sibling *AaBb.cod* file 200 as a hard offset, $HO_{Cc.N}$, specifying the location of the definition of Method Cc.N within sibling *CcDd.cod* file 210. This hard offset needs to be placed into context by the Java VM at link time using information stored in fixup table 208. For example, fixup table 208 may comprise a pointer to the name of sibling *CcDd.cod* file along with the location of the reference in Method Bb.L to Method Cc.N so that the Java VM knows which sibling *.cod* file contains Method Cc.N at the specified hard offset.

4) As shown in FIGS. 1A and 1B, Method Cc.M comprises a reference to Method Ff.U. The *.cod* file for class Ff is outside of the sibling group of *CcDd.cod* file 210. Therefore, as shown in FIG. 2A, this reference appears in the definition of Method Cc.M in "byte codes and information structures" 214 of sibling *CcDd.cod* file 210 as a symbolic reference to Method Ff.U and its argument(s). This symbolic reference needs to be resolved by the Java VM at link time using information stored in fixup table 218. For example, fixup table 218 may comprise a pointer to the name of the *.cod* file for the class where Method Ff.U is defined (e.g. "EeFf.cod"), a pointer to the name of the class where Method Ff.U is defined (e.g. "Ff"), a pointer to the name of the method (e.g. "U") and a pointer to the name of the method typing information (e.g. "I").

[0043] It will be appreciated by persons of ordinary skill in the art that when either of the source code files *Cc.java* and *Dd.java* are modified, it is necessary to invoke the executable software to generate a new *CcDd.cod* file, or alternatively to invoke the executable software on the modified *.java* or *.class* files along with one or more additional *.java* or *.class* files to generate a new *CcDdxxx.cod* file, where "*x*" indicates the additional classes. If a new *CcDd.cod* file or a new *CcDdxxx.cod* file has been generated, then the hard offset $HO_{Cc.N}$ appearing in resolved *AaBb.cod* file 200 no longer accurately describes the location of the byte codes and other information for Method N in the new *CcDd.cod* file or a new *CcDdxxx.cod* file.

[0044] Accordingly, when a *.cod* file is updated, all of its sibling *.cod* files are updated simultaneously. Although the scope of the present invention is not limited in this respect, this is generally done by taking the entire set of *.class* files and regenerating all the output sibling .cod files. Since some of the classes have changed,

the division of classes among the newly generated sibling *.cod* files may be different from that in the original sibling *.cod* files.

[0045] While certain features of the invention have been illustrated and described herein, many modifications, substitutions, changes, and equivalents will now occur to those of ordinary skill in the art. It is, therefore, to be understood that the appended claims are intended to cover all such modifications and changes as fall within the true spirit of the invention.

Claims:

1. A device comprising:

a memory unit;

a computing unit connected to the memory unit, the computing unit being able to execute a Java Virtual Machine; and

one or more files stored in the memory unit, the one or more files directly interpretable by the Java Virtual Machine, the one or more files being generated from a plurality of class files for reducing storage space, wherein a given file of the one or more files comprises:

a constant pool created by combining constant pool entries from two or more of the plurality of class files without duplication of entries;

a byte codes and information structure created by combining byte codes and information structure entries from the two or more of the plurality of class files; and,

a fixup table for providing information to the Java Virtual Machine for resolving at least one entry in the given file at link time.

2. The device of claim 1, wherein the memory unit includes executable software and the computing unit is arranged to execute the executable software to generate the one or more files from the plurality of class files by combining elements from the plurality of class files without duplication of entries for reducing storage space, wherein the number of files generated from the plurality of class files is less than the number of the plurality of class files.

3. The device of claim 1 or claim 2, wherein the information in the fixup table comprises the location of data needed for resolving a symbolic reference in the given file.

4. The device of any one of claims 1 to 3, wherein there are at least two files from the number of files generated from the plurality of class files defined as sibling files in a common sibling group, each of the sibling files comprises a sibling list for listing other sibling files in the common sibling group, wherein cross-references

between the sibling files in the common sibling group are indicated using hard offsets and references to files that are not part of the common sibling group are indicated using symbolic references.

5. The device of claim 4, wherein the information in the fixup table of a given sibling file comprises the location of data of a cross-referenced sibling file to place one of the hard offsets that corresponds to the cross-referenced sibling file into context at link time.

6. The device of any preceding claim, wherein for the given file, the byte codes and information structure includes a second hard offset for cross-referencing a method included in the given file that was previously symbolically referenced.

7. The device of claim 4 or claim 6 wherein at least one of the hard offsets does not need to be resolved or put into context by the Java Virtual Machine at link time.

8. A method for generating one or more files from a plurality of class files, the method comprising:

identifying class files with common entries in constant pools of the class files;

generating a constant pool for a given generated file by combining constant pool entries from the class files with common entries without duplication;

generating the byte codes and information structure for the given generated file by combining byte codes and information structure entries from the class files; and,

generating a fixup table for providing information to a Java Virtual Machine for resolving at least one entry in the given generated file at link time.

9. The method of claim 8, wherein the method further comprises generating the one or more files from a plurality of class files by combining elements from the plurality of class files without duplication of entries for reducing storage space,

wherein the number of files generated from the plurality of class files is less than the number of the plurality of class files.

10. The method of claim 8 or claim 9, wherein the method further includes providing location of data needed for resolving a symbolic reference in the given generated file as the information in the fixup table.

11. The method of any one of claims 8 to 10, wherein the method further includes identifying cross-references between the class files and defining at least two files from the number of files generated from the plurality of class files as sibling files in a common sibling group, each of the sibling files comprising a sibling list for listing other sibling files in the common sibling group, indicating cross-references between the sibling files in the common sibling group using hard offsets, and indicating references to files that are not part of the common sibling group using symbolic references.

12. The method of claim 11, wherein the method further includes providing the location of data of a cross-referenced sibling file as the information in the fixup table of a given sibling file to place one of the hard offsets that corresponds to the cross-referenced sibling file into context at link time.

13. The method of any one of claims 8 to 12, wherein the method further includes, for the given generated file, providing a second hard offset in the byte codes and information structure for cross-referencing a method now included in the given generated file that was previously symbolically referenced.

14. The method of claim 11 or claim 13, wherein the method further includes providing at least one of the hard offsets in a manner that does not need to be resolved or put into context by the Java Virtual Machine at link time.