Watermark Generator, Watermark Decoder, Method for Providing a Watermark Signal, Method for Providing Binary Message Data in Dependence on a Watermarked Signal and a Computer Program using Improved Synchronization Concept

Description

Technical Field

Embodiments according to the invention are related to a watermark generator for providing a watermark signal in dependence on binary message data. Further embodiments according to the invention relate to a watermark decoder for providing binary message data in dependence on a watermarked signal. Further embodiments according to the invention are related to a method for providing a watermark signal in dependence on binary message data. Further embodiments according to the invention are related to a method for providing binary message data in dependence on a watermarked signal. Further embodiments are related to corresponding computer programs.

Some embodiments according to the invention are related to a robust low complexity audio watermarking system.

Background of the Invention

In many technical applications, it is desired to include an extra information into an information or signal representing useful data or "main data" like, for example, an audio signal, a video signal, graphics, a measurement quantity and so on. In many cases, it is desired to include the extra information such that the extra information is bound to the main data (for example, audio data, video data, still image data, measurement data, text data, and so on) in a way that it is not perceivable by a user of said data. Also, in some cases it is desirable to include the extra data such that the extra data are not easily removable from the main data (e.g. audio data, video data, still image data, measurement data, and so on).

This is particularly true in applications in which it is desirable to implement a digital rights management. However, it is sometimes simply desired to add substantially unperceivable side information to the useful data. For example, in some cases it is desirable to add side

information to audio data, such that the side information provides an information about the source of the audio data, the content of the audio data, rights related to the audio data and so on.

For embedding extra data into useful data or "main data", a concept called "watermarking" may be used. Watermarking concepts have been discussed in the literature for many different kinds of useful data, like audio data, still image data, video data, text data, and so on.

In the following, some references will be given in which watermarking concepts are discussed. However, the reader's attention is also drawn to the wide field of textbook literature and publications related to the watermarking for further details.

DE 196 40 814 C2 describes a coding method for introducing a non-audible data signal into an audio signal and a method for decoding a data signal, which is included in an audio signal in a non-audible form. The coding method for introducing a non-audible data signal into an audio signal comprises converting the audio signal into the spectral domain. The coding method also comprises determining the masking threshold of the audio signal and the provision of a pseudo noise signal. The coding method also comprises providing the data signal and multiplying the pseudo noise signal with the data signal, in order to obtain a frequency-spread data signal. The coding method also comprises weighting the spread data signal with the masking threshold and overlapping the audio signal and the weighted data signal.

In addition, WO 93/07689 describes a method and apparatus for automatically identifying a program broadcast by a radio station or by a television channel, or recorded on a medium, by adding an inaudible encoded message to the sound signal of the program, the message identifying the broadcasting channel or station, the program and/or the exact date. In an embodiment discussed in said document, the sound signal is transmitted via an analog-to-digital converter to a data processor enabling frequency components to be split up, and enabling the energy in some of the frequency components to be altered in a predetermined manner to form an encoded identification message. The output from the data processor is connected by a digital-to-analog converter to an audio output for broadcasting or recording the sound signal. In another embodiment discussed in said document, an analog bandpass is employed to separate a band of frequencies from the sound signal so that energy in the separated band may be thus altered to encode the sound signal.

US 5, 450,490 describes apparatus and methods for including a code having at least one code frequency component in an audio signal. The abilities of various frequency components in the audio signal to mask the code frequency component to human hearing are evaluated and based on these evaluations an amplitude is assigned to the code frequency component. Methods and apparatus for detecting a code in an encoded audio signal are also described. A code frequency component in the encoded audio signal is detected based on an expected code amplitude or on a noise amplitude within a range of audio frequencies including the frequency of the code component.

WO 94/11989 describes a method and apparatus for encoding/decoding broadcast or recorded segments and monitoring audience exposure thereto. Methods and apparatus for encoding and decoding information in broadcasts or recorded segment signals are described. In an embodiment described in the document, an audience monitoring system encodes identification information in the audio signal portion of a broadcast or a recorded segment using spread spectrum encoding. The monitoring device receives an acoustically reproduced version of the broadcast or recorded signal via a microphone, decodes the identification information from the audio signal portion despite significant ambient noise and stores this information, automatically providing a diary for the audience member, which is later uploaded to a centralized facility. A separate monitoring device decodes additional information from the broadcast signal, which is matched with the audience diary information at the central facility. This monitor may simultaneously send data to the centralized facility using a dial-up telephone line, and receives data from the centralized facility through a signal encoded using a spread spectrum technique and modulated with a broadcast signal from a third party.

WO 95/27349 describes apparatus and methods for including codes in audio signals and decoding. An apparatus and methods for including a code having at least one code frequency component in an audio signal are described. The abilities of various frequency components in the audio signal to mask the code frequency component to human hearing are evaluated, and based on these evaluations, an amplitude is assigned to the code frequency components. Methods and apparatus for detecting a code in an encoded audio signal are also described. A code frequency component in the encoded audio signal is detected based on an expected code amplitude or on a noise amplitude within a range of audio frequencies including the frequency of the code component.

However, in the known watermarking systems, it is sometimes difficult or time consuming to acquire appropriate synchronization at the decoder side. Also, in some cases, a resource

effort for transmitting synchronization information from a watermark generator to a watermark decoder is comparatively high. Also, in some cases, reliability problems arise.

In view of this situation, it is the object of the present invention to create a synchronization concept for use in the transmission of a watermark, which allows for a good tradeoff between resource efficiency, reliability and time required to acquire the synchronization.

Summary of the Invention

The objective is achieved by a watermark generator according to claim 1, a watermark decoder according to claim 10, a method for providing a watermark signal according to claim 17, a method for providing binary message data according to claim 18 and a computer program according to claim 19.

An embodiment according to the invention creates a watermark generator for providing a watermark signal in dependence on binary message data. The watermark generator comprises an information spreader configured to spread an information unit to a plurality of time-frequency-domain values, to obtain a spread information representation. The watermark generator also comprises a synchronization inserter configured to multiplicatively combine the spread information representation with a synchronization sequence to obtain a combined information-synchronization representation. The watermark generator also comprises a watermark signal provider configured to provide the watermark signal on the basis of the combined information-synchronization representation.

It is the key idea of the present invention that a synchronization information, which allows for a fast synchronization recovery and still provides for a good tradeoff between reliability and resource efficiency can be obtained, at the side of a watermark generator, by a multiplicatively combining a spread information representation, which represents one or more information units (for example bits) of the binary message data, with a synchronization sequence, such that a combined information-synchronization representation can be obtained. The combined information-synchronization representation, on the basis of which the watermark signal is provided, carries both message data information and synchronization information simultaneously and at identical frequencies or within identical frequency bands.

Also, by multiplicatively combining the spread information representation with the synchronization sequence, it can be ensured (for example by an appropriate choice of the magnitudes of the values of the synchronization sequence and/or the spread information

representation) that the intensity of the values of the combined information-synchronization representation is well-defined and predictable even without consideration of the data contents. Accordingly, the concept of multiplicatively combining the spread information representation with the synchronization sequence facilitates the possibility to set the intensity of the combined information-synchronization information to a maximum value which is allowable without perceivably distorting the content of a useful signal, in which the watermark signal is introduced.

To summarize, the described concept to multiplicatively combine the spread information representation with a synchronization sequence, to obtain a combined information-synchronization representation allows to use a continuous synchronization, which is detectable rapidly, because there is no need for a time slicing, which is detectable with high reliability, because the synchronization sequence can be spread over a wide frequency range, and which is well-controlled with respect to the energy, because the multiplicative combination can be conducted such that the energy of a value of the combined information-synchronization representation is independent from the data content of the spread information representation.

In a preferred embodiment, the information spreader is configured to frequency-spread a bit of the binary message data in accordance with a bit spread sequence, such that, in the spread information representation, a set of time-frequency-domain values associated with a plurality of different frequencies (or frequency bands) is associated to the bit. In this case, the synchronization sequence comprises a plurality of synchronization spread sequences, the synchronization spread sequences defining a plurality of different synchronization patterns spread in frequency. This embodiment allows to apply different synchronization spread sequences to different time portions of the spread information representation, such that different synchronization spread sequences are multiplicatively combined with spread representations of different bits of the binary message data.

In a preferred embodiment, the synchronization inserter is configured to multiplicatively combine different time portions of the spread information representation, which different time portions represent different bits of the binary message data, with different synchronization spread sequences, such that the different synchronization spread sequences are combined with different bit spread sequences or with differently scaled versions of a given spread sequence.

In a preferred embodiment, the synchronization inserter is configured to multiplicatively combine the spread information representation with a periodic sequence of synchronization

spread sequences. In a first case, a periodicity of the sequence of synchronization spread sequences is equal to a number of bits of a binary message (for example of a block of bits provided as an output information of a block-based channel coder). Accordingly, a unique identification of a temporal alignment of a message can be made. In an alternative case, a periodicity of the sequence of synchronization spread sequences is shorter than the number of bits of the binary message. In this case, an additional mechanism may be used in order to uniquely obtain a message synchronization.

In a preferred embodiment, the synchronization spread sequences are chosen to be orthogonal. This improves the quality of a synchronization acquisition at the decoder side.

In another preferred embodiment, the synchronization inserter is configured to multiplicatively combine each time portion of the spread information representation with at least one synchronization sequence, such that each time portion of the combined information-synchronization representation comprises a plurality of values which are dependent on a bit of the binary message data and which also represent a synchronization information. Accordingly, there is a continuous synchronization information, which accelerates the acquisition of the synchronization at the decoder side.

In a preferred embodiment, the information spreader is configured to selectively spread a given bit onto a first bit representation, which is a positive multiple of a bit spread sequence, or onto a second bit representation, which is a negative multiple of the bit spread sequence, in dependence on the value of the given bit.

An embodiment according to the invention creates a watermark decoder for providing binary message data in dependence on a watermarked signal. The watermark decoder comprises a time-frequency-domain representation provider configured to provide a time-frequency-domain representation of the watermarked signal, for example an analyzer or a filterbank configured to analyze the watermarked signal. The watermark decoder also comprises a synchronization determinator configured to find temporal alignment of watermark information within the time-frequency-domain representation of the watermarked signal. The synchronization determinator is configured to compute a first sum of result values of an element-wise multiplication between a first subset of values of the time-frequency-domain representation, values of a first synchronization despreading sequence and values of a bit despreading sequence, and to compute a second sum of result values of an element-wise multiplication between a second subset of values of the time-frequency-domain representation, values of a second synchronization despreading sequence and values of the bit despreading sequence. The synchronization determinator is

also configured to sum absolute values of the first sum of result values and of the second sum of result values, to obtain a likelihood measure representing a likelihood that the first and second subsets of values of the time-frequency-domain representation carry a synchronization information, which is in temporal alignment with the synchronization spreading sequences. The synchronization determinator is also configured to determine the synchronization in dependence on the likelihood measure.

This embodiment according to the invention is based on the finding that the synchronization, i.e. the information regarding the temporal alignment of the watermark information within the time-frequency-domain representation of the watermarked signal, can be found in a watermarked signal, in which a spread information representation is multiplicatively combined with a synchronization sequence, independent on the binary message data determining the spread information representation, by computing the first sum of result values of an element-wise multiplication between a first subset of values of the time-frequency-domain representation, values of a first synchronization despreading sequence and values of a bit despreading sequence, computing a second sum of result values of an element-wise multiplication between a second subset of values of the time-frequency-domain representation, values of a second synchronization despreading sequence and values of the bit despreading sequence, and by summing absolute values of the first sum of result values and of the second sum of result values. By taking the sum of the absolute values of the first and second sum of result values, a so-called non-coherent correlation can be computed, which results in a large correlation value (or even a correlation peak) if a synchronization is found, independent of the data of the spread information representation, which is multiplicatively combined with the synchronization sequence to obtain the combined information-synchronization representation, which is included in the watermarked signal.

In other words, the formation of the absolute value of the first sum of result values and of the second sum of result values may render a correlation, which comprises the element-wise multiplication between the first subset of values of the time-frequency-domain representation, values of the first synchronization despreading sequence and values of the bit despreading sequence, and also the element-wise multiplication between the second subset of values of the time-frequency-domain representation, values of the second synchronization despreading sequence and values of the bit despreading sequence, data independent. Accordingly, a reliable and data-independent detection of the synchronization can be obtained.

In a preferred embodiment, the first subset of values of the time-frequency-domain representation comprises values of the time-frequency-domain representation associated with a first time portion of the time-frequency-domain representation, and the second subset of values of the time-frequency-domain representation comprises values of the time-frequency-domain representation associated with a second time portion of the time-frequency-domain representation. Also, preferably, the first synchronization despreading sequence and the second synchronization despreading sequence are orthogonal. This embodiment allows for a clear detection of a synchronization using a synchronization pattern which comprises orthogonal synchronization spread sequences at different time portions.

In a preferred embodiment, the synchronization determinator is configured to receive, as the time-frequency-domain representation, a hard bit information or a soft bit information. In this case, the synchronization determinator is configured to apply the first synchronization despreading sequence to the first subset of values, to reverse a first synchronization embedding information, and to apply the second synchronization despreading sequence to the second subset of values, to reverse a second synchronization embedding operation. In addition, the synchronization determinator is configured to apply the bit despreading sequence to reverse a bit spreading operation. Accordingly, the synchronization determinator is preferably configured to reverse the both the bit spreading operation and the synchronization embedding operation, which are performed at the encoder side.

In a preferred embodiment, the synchronization determinator is configured to compute likelihood measures for a plurality of positional choices of the first and second subset with reference to the time-frequency-domain representation, to obtain a sequence of likelihood measures associated with the different positional choices of the first and second subset. The synchronization determinator is configured to correlate the sequence of likelihood measures with a reference sequence, in order to determine a synchronization hit, or to filter the sequence of likelihood measures with a matched filter adapted to an expected sequence of likelihood values, in order to determine a synchronization hit. By using a correlation of the sequence of likelihood measures, or by filtering the sequence of likelihood measures with a matched filter, an improved accuracy of the synchronization and/or an improved reliability of the detection of a synchronization hit can be obtained.

Some embodiments according to the invention also create a method for providing a watermark signal in dependence on binary message data and a method for providing binary

message data in dependence on a watermarked signal. Said methods are based on the same findings as the apparatus discussed before.

Some embodiments according to the invention comprise a computer program for performing the inventive methods.

Brief Description of the Figures

Embodiments according to the invention will subsequently be described taking reference to the enclosed figures, in which:

Fig. 1 shows a block schematic diagram of a watermark inserter according to an embodiment of the invention;

Fig. 2 shows a block-schematic diagram of a watermark decoder, according to an embodiment of the invention;

Fig. 3 shows a detailed block-schematic diagram of a watermark generator, according to an embodiment of the invention;

Fig. 4 shows a detailed block-schematic diagram of a modulator, for use in an embodiment of the invention;

Fig. 5 shows a detailed block-schematic diagram of a psychoacoustical processing module, for use in an embodiment of the invention;

Fig. 6 shows a block-schematic diagram of a psychoacoustical model processor, for use in an embodiment of the invention;

Fig. 7 shows a graphical representation of a power spectrum of an audio signal output by block 801 over frequency;

Fig. 8 shows a graphical representation of a power spectrum of an audio signal output by block 802 over frequency;

Fig. 9 shows a block-schematic diagram of an amplitude calculation;

Fig. 10a shows a block schematic diagram of a modulator;

Fig. 10b shows a graphical representation of the location of coefficients on the time-frequency claim;

Figs. 11a and 11b show a block-schematic diagrams of implementation alternatives of the synchronization module;

Fig. 12a shows a graphical representation of the problem of finding the temporal alignment of a watermark;

Fig. 12b shows a graphical representation of the problem of identifying the message start;

Fig. 12c shows a graphical representation of a temporal alignment of synchronization sequences in a full message synchronization mode;

Fig. 12d shows a graphical representation of the temporal alignment of the synchronization sequences in a partial message synchronization mode;

Fig. 12e shows a graphical representation of input data of the synchronization module;

Fig. 12f shows a graphical representation of a concept of identifying a synchronization hit;

Fig. 12g shows a block-schematic diagram of a synchronization signature correlator;

Fig. 13a shows a graphical representation of an example for a temporal despreading;

Fig. 13b shows a graphical representation of an example for an element-wise multiplication between bits and spreading sequences;

Fig. 13c shows a graphical representation of an output of the synchronization signature correlator after temporal averaging;

Fig. 13d shows a graphical representation of an output of the synchronization signature correlator filtered with the auto-correlation function of the synchronization signature;

Fig. 14 shows a block-schematic diagram of a watermark extractor, according to an embodiment of the invention;

Fig. 15 shows a schematic representation of a selection of a part of the time-frequency-domain representation as a candidate message;

Fig. 16 shows a block-schematic diagram of an analysis module;

Fig. 17a shows a graphical representation of an output of a synchronization correlator;

Fig. 17b shows a graphical representation of decoded messages;

Fig. 17c shows a graphical representation of a synchronization position, which is extracted from a watermarked signal;

Fig. 18a shows a graphical representation of a payload, a payload with a Viterbi termination sequence, a Viterbi-encoded payload and a repetition-coded version of the Viterbi-coded payload;

Fig. 18b shows a graphical representation of subcarriers used for embedding a watermarked signal;

Fig. 19 shows a graphical representation of an uncoded message, a coded message, a synchronization message and a watermark signal, in which the synchronization sequence is applied to the messages;

Fig. 20 shows a schematic representation of a first step of a so-called "ABC synchronization" concept;

Fig. 21 shows a graphical representation of a second step of the so-called "ABC synchronization" concept;

Fig. 22 shows a graphical representation of a third step of the so-called "ABC synchronization" concept;

Fig. 23 shows a graphical representation of a message comprising a payload and a CRC portion;

Fig. 24 shows a block-schematic diagram of a watermark generator, according to an embodiment of the invention;

Fig. 25 shows a block-schematic diagram of a watermark decoder, according to an embodiment of the invention;

Fig. 26 shows a flowchart of a method for providing a watermark signal in dependence on binary message data; and

Fig. 27 shows a flowchart of a method for providing binary message data in dependence on a watermarked signal.

Detailed Description of the Embodiments

1. Watermark generation

1.1 Watermark generator according to Fig. 24

In the following, a watermark generator 2400 will be described taking reference to Fig. 24, which shows a block-schematic diagram of such a watermark generator.

The watermark generator 2400 is configured to receive binary message data 2410 and to provide, on the basis thereof, a watermark signal 2420. The watermark generator comprises an information spreader 2430, which is configured to receive an information unit, for example a bit of the binary message data 2410, and to spread the information unit to a plurality of time-frequency-domain values, to obtain a spread information 2432. The watermark generator 2400 also comprises a synchronization inserter 2440, which is configured to receive the spread information 2432 and a synchronization sequence 2442 and to multiplicatively combine the spread information representation 2432 with the synchronization sequence 2442, to obtain a combined information-synchronization representation 2444. The watermark generator 2400 also comprises a watermark signal provider 2450 configured to provide the watermark signal 2420 on the basis of the combined information-synchronization information 2444.

The watermark generator 2400 may be supplemented by any of the features and functionalities which are discussed in more detail in section 3 below.

1.2. Method for providing a watermark signal in dependence on binary message data according to Fig. 26

In the following, a method for providing a watermark signal in dependence on binary message data will be explained taking reference to Fig. 26. Fig. 26 shows a flowchart of such a method.

The method 2600 of Fig. 26 comprises a step 2610 of spreading an information unit (for example, a bit) to a plurality of time-frequency-domain values, to obtain a spread information representation. The method 2600 also comprises a step 2620 of multiplicatively combining the spread information representation with a synchronization sequence to obtain a combined information-synchronization representation. The method 2600 also comprises a step 2630 of providing the watermark signal on the basis of the combined information-synchronization representation.

Naturally, the method 2600 can be supplemented by any of the features and functionalities discussed herein, also with respect to the inventive apparatus.

2. Watermark decoding

2.1 Watermark decoder according to Fig. 25

In the following, a watermark decoder 2500 will be described taking reference to Fig. 25, which shows a block-schematic diagram of such a watermark decoder. The watermark decoder 2500 is configured to provide binary message data 2520 in dependence on a watermarked signal 2510.

The watermark decoder 2500 comprises a time-frequency-domain representation provider, which is configured to provide a time-frequency-domain representation 2532 of the watermarked signal 2510. The time-frequency-domain representation provider 2530 may, for example, comprise a time-domain to time-frequency-domain transformer or a filterbank.

The watermark decoder 2500 further comprises a synchronization determinator 2540, which is configured to find a temporal alignment of the watermark information within the time-frequency-domain representation 2532 of the watermarked signal 2510. The synchronization determinator 2540 is configured to compute a first sum of result values of an element-wise multiplication between a first subset of values of the time-frequency-domain representation 2532, values of a first synchronization despreading sequence and values of a bit despreading sequence, and to compute a second sum of result values of an element-wise multiplication between a second subset of values of the time-frequency-domain representation, values of a second synchronization despreading sequence and values of the bit despreading sequence. The synchronization determinator 2540 is also configured to sum absolute values of the first sum of result values and of the second sum of result values, to obtain a likelihood measure representing a likelihood that the first and second subsets of values of the time-frequency-domain representation carry a synchronization information, which is in temporal alignment with the synchronization spreading sequence. The synchronization determinator 2540 is also configured to determine the synchronization in dependence on the likelihood measure, and to provide an appropriate synchronization information 2542.

The watermark decoder 2500 may be supplemented by any of the means and functionalities discussed herein with respect to the watermark decoding.

2.2 Method for providing binary message data in dependence on a watermarked signal according to Fig. 27

In the following, a method 2700 for providing binary message data in dependence on a watermarked signal will be described taking reference to Fig. 27, which shows a flow chart of such a method.

The method 2700 comprises a step 2710 of providing a time-frequency-domain representation of the watermarked signal. The method 2700 also comprises step 2720 of finding a temporal alignment of watermark information within the time-frequency-domain representation of the watermarked signal. The step 2720 of finding a temporal alignment of watermark information within the time-frequency-domain representation of the watermarked signal comprises a substep 2722 of computing a first sum of result values of an element-wise multiplication between a first subset of values of the time-frequency-domain representation, values of a first synchronization despreading sequence and values of a bit despreading sequence. The step 2720 also comprises substep 2724 of computing a second sum of result values of an element-wise multiplication between a second subset of

values of the time-frequency-domain representation, values of a second synchronization despreading sequence and values of the bit despreading sequence. The step 2720 also comprises a substep 2726 of summing absolute values of the first sum of result values and of the second sum of result values, to obtain a likelihood measure representing a likelihood that the first and second subsets of values of the time-frequency-domain representation carry a synchronization information, which is in temporal alignment with the synchronization spreading sequence. The step 2720 also comprises a substep 2728 of determining the synchronization dependent on the likelihood measure.

The method 2700 can be supplemented by any of the features and functionalities described herein with the watermark decoding.

3. System Description

In the following, a system for a watermark transmission will be described, which comprises a watermark inserter and a watermark decoder. Naturally, the watermark inserter and the watermark decoder can be used independent from each other.

For the description of the system a top-down approach is chosen here. First, it is distinguished between encoder and decoder. Then, in sections 3.1 to 3.5 each processing block is described in detail.

The basic structure of the system can be seen in Figures 1 and 2, which depict the encoder and decoder side, respectively. Fig 1 shows a block schematic diagram of a watermark inserter 100. At the encoder side, the watermark signal 101b is generated in the processing block 101 (also designated as watermark generator) from binary data 101a and on the basis of information 104, 105 exchanged with the psychoacoustical processing module 102. The information provided from block 102 typically guarantees that the watermark is inaudible. The watermark generated by the watermark generator101 is then added to the audio signal 106. The watermarked signal 107 can then be transmitted, stored, or further processed. In case of a multimedia file, e.g., an audio-video file, a proper delay needs to be added to the video stream not to lose audio-video synchronicity. In case of a multichannel audio signal, each channel is processed separately as explained in this document. The processing blocks 101 (watermark generator) and 102 (psychoacoustical processing module) are explained in detail in Sections 3.1 and 3.2, respectively.

The decoder side is depicted in Figure 2, which shows a block schematic diagram of a watermark detector 200. A watermarked audio signal 200a, e.g., recorded by a microphone, is made available to the system 200. A first block 203, which is also designated as an analysis module, demodulates and transforms the data (e.g., the watermarked audio signal) in time/frequency domain (thereby obtaining a time-frequency-domain representation 204 of the watermarked audio signal 200a) passing it to the synchronization module 201, which analyzes the input signal 204 and carries out a temporal synchronization, namely, determines the temporal alignment of the encoded data (e.g. of the encoded watermark data relative to the time-frequency-domain representation). This information (e.g., the resulting synchronization information 205) is given to the watermark extractor 202, which decodes the data (and consequently provides the binary data 202a, which represent the data content of the watermarked audio signal 200a).

3.1 The Watermark Generator 101

The watermark generator 101 is depicted detail in Figure 3. Binary data (expressed as ± 1) to be hidden in the audio signal 106 is given to the watermark generator 101. The block 301 organizes the data 101a in packets of equal length M_p. Overhead bits are added (e.g. appended) for signaling purposes to each packet. Let M_s denote their number. Their use will be explained in detail in Section 3.5. Note that in the following each packet of payload bits together with the signaling overhead bits is denoted message.

Each message 301a, of length $N_m = M_s + M_p$, is handed over to the processing block 302, the channel encoder, which is responsible of coding the bits for protection against errors. A possible embodiment of this module consists of a convolutional encoder together with an interleaver. The ratio of the convolutional encoder influences greatly the overall degree of protection against errors of the watermarking system. The interleaver, on the other hand, brings protection against noise bursts. The range of operation of the interleaver can be limited to one message but it could also be extended to more messages. Let R_c denote the code ratio, e.g., 1/4. The number of coded bits for each message is N_m/R_c. The channel encoder provides, for example, an encoded binary message 302a.

The next processing block, 303, carries out a spreading in frequency domain. In order to achieve sufficient signal to noise ratio, the information (e.g. the information of the binary message 302a) is spread and transmitted in N_f carefully chosen subbands. Their exact position in frequency is decided a priori and is known to both the encoder and the decoder. Details on the choice of this important system parameter is given in Section 3.2.2. The

spreading in frequency is determined by the spreading sequence c_f of size $N_f \times 1$. The output 303a of the block 303 consists of N_f bit streams, one for each subband. The i-th bit stream is obtained by multiplying the input bit with the i-th component of spreading sequence c_f. The simplest spreading consists of copying the bit stream to each output stream, namely use a spreading sequence of all ones.

Block 304, which is also designated as a synchronization scheme inserter, adds a synchronization signal to the bit stream. A robust synchronization is important as the decoder does not know the temporal alignment of neither bits nor the data structure, i.e., when each message starts. The synchronization signal consists of N_s sequences of N_f bits each. The sequences are multiplied element wise and periodically to the bit stream (or bit streams 303a). For instance, let **a**, **b**, and **c**, be the Ns = 3 synchronization sequences (also designated as synchronization spreading sequences). Block 304 multiplies **a** to the first spread bit, **b** to the second spread bit, and **c** to the third spread bit. For the following bits the process is periodically iterated, namely, **a** to the fourth bit, **b** for the fifth bit and so on. Accordingly, a combined information-synchronization information 304a is obtained. The synchronization sequences (also designated as synchronization spread sequences) are carefully chosen to minimize the risk of a false synchronization. More details are given in Section 3.4. Also, it should be noted that a sequence **a**, **b**, **c,...** may be considered as a sequence of synchronization spread sequences.

Block 305 carries out a spreading in time domain. Each spread bit at the input, namely a vector of length N_f, is repeated in time domain N_t times. Similarly to the spreading in frequency, we define a spreading sequence c_t of size $N_t \times 1$. The i-th temporal repetition is multiplied with the i-th component of c_t.

The operations of blocks 302 to 305 can be put in mathematical terms as follows. Let **m** of size $1 \times N_m = R_c$ be a coded message, output of 302. The output 303a (which may be considered as a spread information representation **R**) of block 303 is

$$c_f \cdot m \quad \text{of size} \quad N_f \times N_m / R_c \tag{1}$$

the output 304a of block 304, which may be considered as a combined information-synchronization representation **C**, is

$$S \circ (c_f \cdot m) \quad \text{of size} \quad N_f \times N_m / R_c \tag{2}$$

where ∘ denotes the Schur element-wise product and

$$S = [\ \ldots \quad a \quad b \quad c \quad \ldots \quad a \quad b \quad \ldots \ ] \quad \text{of size} \quad N_{\mathrm{f}} \times N_{\mathrm{m}}/R_{\mathrm{c}}. \tag{3}$$

The output 305a of 305 is

$$(S \circ (c_{\mathrm{f}} \cdot m)) \diamond c_{\mathrm{t}}^{\mathrm{T}} \quad \text{of size} \quad N_{\mathrm{f}} \times N_{\mathrm{t}} \cdot N_{\mathrm{m}}/R_{\mathrm{c}} \tag{4}$$

where ◇ and ^{T} denote the Kronecker product and transpose, respectively. Please recall that binary data is expressed as ±1.

Block **306** performs a differential encoding of the bits. This step gives the system additional robustness against phase shifts due to movement or local oscillator mismatches. More details on this matter are given in Section 3.3. If b(i; j) is the bit for the i-th frequency band and j-th time block at the input of block 306, the output bit b_{diff} (i; j) is

$$b_{\mathrm{diff}}(i,j) = b_{\mathrm{diff}}(i,j-1) \cdot b(i,j). \tag{5}$$

At the beginning of the stream, that is for j = 0, b_{diff} (i,j - 1) is set to 1.

Block 307 carries out the actual modulation, i.e., the generation of the watermark signal waveform depending on the binary information 306a given at its input. A more detailed schematics is given in Figure 4. N_f parallel inputs, 401 to $40N_f$ contain the bit streams for the different subbands. Each bit of each subband stream is processed by a bit shaping block (411 to $41N_f$). The output of the bit shaping blocks are waveforms in time domain. The waveform generated for the j-th time block and i-th subband, denoted by $s_{i,j}(t)$, on the basis of the input bit b_{diff}(i, j)
is computed as follows

$$s_{i,j}(t) = b_{\mathrm{diff}}(i,j)\gamma(i,j) \cdot g_i(t - j \cdot T_{\mathrm{b}}), \tag{6}$$

where $\gamma(i; j)$ is a weighting factor provided by the psychoacoustical processing unit 102, T_b is the bit time interval, and $g_i(t)$ is the bit forming function for the i-th subband. The bit forming function is obtained from a baseband function $g_i^T(t)$ modulated in frequency with a cosine

$$g_i(t) = g_i^T(t) \cdot \cos(2\pi f_i t) \tag{7}$$

where f_i is the center frequency of the i-th subband and the superscript T stands for transmitter. The baseband functions can be different for each subband. If chosen identical, a more efficient implementation at the decoder is possible. See Section 3.3 for more details.

The bit shaping for each bit is repeated in an iterative process controlled by the psychoacoustical processing module (102). Iterations are necessary to fine tune the weights $\gamma(i, j)$ to assign as much energy as possible to the watermark while keeping it inaudible. More details are given in Section 3.2.

The complete waveform at the output of the i-th bit shaping fillter 41i is

$$s_i(t) = \sum_j s_{i,j}(t). \tag{8}$$

The bit forming baseband function $g_i^T(t)$ is normally non zero for a time interval much larger than T_b, although the main energy is concentrated within the bit interval. An example can be seen if Figure 12a where the same bit forming baseband function is plotted for two adjacent bits. In the figure we have T_b = 40 ms. The choice of T_b as well as the shape of the function affect the system considerably. In fact, longer symbols provide narrower frequency responses. This is particularly beneficial in reverberant environments. In fact, in such scenarios the watermarked signal reaches the microphone via several propagation paths, each characterized by a different propagation time. The resulting channel exhibits strong frequency selectivity. Interpreted in time domain, longer symbols are beneficial as echoes with a delay comparable to the bit interval yield constructive interference, meaning that they increase the received signal energy. Notwithstanding, longer symbols bring also a few drawbacks; larger overlaps might lead to intersymbol interference (ISI) and are for sure more difficult to hide in the audio signal, so that the psychoacoustical processing module would allow less energy than for shorter symbols.

The watermark signal is obtained by summing all outputs of the bit shaping filters

$$\sum_i s_i(t). \tag{9}$$

3.2 The Psychoacoustical Processing Module 102

As depicted in Figure 5, the psychoacoustical processing module 102 consists of 3 parts. The first step is an analysis module **501** which transforms the time audio signal into the time/frequency domain. This analysis module may carry out parallel analyses in different time/frequency resolutions. After the analysis module, the time/frequency data is transferred to the psychoacoustic model (PAM) **502**, in which masking thresholds for the watermark signal are calculated according to psychoacoustical considerations (see E. Zwicker H.Fastl, "Psychoacoustics Facts and models"). The masking thresholds indicate the amount of energy which can be hidden in the audio signal for each subband and time block. The last block in the psychoacoustical processing module 102 depicts the amplitude calculation module **503**. This module determines the amplitude gains to be used in the generation of the watermark signal so that the masking thresholds are satisfied, i.e., the embedded energy is less or equal to the energy defined by the masking thresholds.

3.2.1 The Time/Frequency Analysis 501

Block **501** carries out the time/frequency transformation of the audio signal by means of a lapped transform. The best audio quality can be achieved when multiple time/frequency resolutions are performed. One efficient embodiment of a lapped transform is the short time Fourier transform (STFT), which is based on fast Fourier transforms (FFT) of windowed time blocks. The length of the window determines the time/frequency resolution, so that longer windows yield lower time and higher frequency resolutions, while shorter windows vice versa. The shape of the window, on the other hand, among other things, determines the frequency leakage.

For the proposed system, we achieve an inaudible watermark by analyzing the data with two different resolutions. A first filter bank is characterized by a hop size of T_b, i.e., the bit length. The hop size is the time interval between two adjacent time blocks. The window length is approximately T_b. Please note that the window shape does not have to be the

same as the one used for the bit shaping, and in general should model the human hearing system. Numerous publications study this problem.

The second filter bank applies a shorter window. The higher temporal resolution achieved is particularly important when embedding a watermark in speech, as its temporal structure is in general finer than T_b.

The sampling rate of the input audio signal is not important, as long as it is large enough to describe the watermark signal without aliasing. For instance, if the largest frequency component contained in the watermark signal is 6 kHz, then the sampling rate of the time signals must be at least 12 kHz.

3.2.2 The Psychoacoustical Model 502

The psychoacoustical model 502 has the task to determine the masking thresholds, i.e., the amount of energy which can be hidden in the audio signal for each subband and time block keeping the watermarked audio signal indistinguishable from the original.

The i-th subband is defined between two limits, namely $f_i^{(\min)}$ and $f_i^{(\max)}$. The subbands are determined by defining N_f center frequencies f_i and letting $f_{i-1}^{(\max)} = f_i^{(\min)}$ i for i = 2, 3, ... , N_f . An appropriate choice for the center frequencies is given by the Bark scale proposed by Zwicker in 1961. The subbands become larger for higher center frequencies. A possible implementation of the system uses 9 subbands ranging from 1.5 to 6 kHz arranged in an appropriate way.

The following processing steps are carried out separately for each time/frequency resolution for each subband and each time block. The processing step **801** carries out a spectral smoothing. In fact, tonal elements, as well as notches in the power spectrum need to be smoothed. This can be carried out in several ways. A tonality measure may be computed and then used to drive an adaptive smoothing filter. Alternatively, in a simpler implementation of this block, a median-like filter can be used. The median filter considers a vector of values and outputs their median value. In a median-like filter the value corresponding to a different quantile than 50% can be chosen. The filter width is defined in Hz and is applied as a non-linear moving average which starts at the lower frequencies and ends up at the highest possible frequency. The operation of **801** is illustrated in Figure 7. The red curve is the output of the smoothing.

Once the smoothing has been carried out, the thresholds are computed by block **802** considering only frequency masking. Also in this case there are different possibilities. One way is to use the minimum for each subband to compute the masking energy E_i. This is the equivalent energy of the signal which effectively operates a masking. From this value we can simply multiply a certain scaling factor to obtain the masked energy J_i. These factors are different for each subband and time/frequency resolution and are obtained via empirical psychoacoustical experiments. These steps are illustrated in Figure 8.

In block **805**, temporal masking is considered. In this case, different time blocks for the same subband are analyzed. The masked energies J_i are modified according to an empirically derived postmasking profile. Let us consider two adjacent time blocks, namely k-1 and k. The corresponding masked energies are $J_i(k-1)$ and $J_i(k)$. The postmasking profile defines that, e.g., the masking energy E_i can mask an energy J_i at time k and $\alpha \cdot J_i$ at time k+1. In this case, block **805** compares $J_i(k)$ (the energy masked by the current time block) and $\alpha \cdot J_i(k+1)$ (the energy masked by the previous time block) and chooses the maximum. Postmasking profiles are available in the literature and have been obtained via empirical psychoacoustical experiments. Note that for large T_b, i.e., > 20 ms, postmasking is applied only to the time/frequency resolution with shorter time windows.

Summarizing, at the output of block **805** we have the masking thresholds per each subband and time block obtained for two different time/frequency resolutions. The thresholds have been obtained by considering both frequency and time masking phenomena. In block **806**, the thresholds for the different time/frequency resolutions are merged. For instance, a possible implementation is that **806** considers all thresholds corresponding to the time and frequency intervals in which a bit is allocated, and chooses the minimum.

3.2.3 The Amplitude Calculation Block 503

Please refer to Figure 9. The input of 503 are the thresholds 505 from the psychoacoustical model 502 where all psychoacoustics motivated calculations are carried out. In the amplitude calculator 503 additional computations with the thresholds are performed. First, an amplitude mapping **901** takes place. This block merely converts the masking thresholds (normally expressed as energies) into amplitudes which can be used to scale the bit shaping function defined in Section 3.1. Afterwards, the amplitude adaptation block **902** is run. This block iteratively adapts the amplitudes $\gamma(i, j)$ which are used to multiply the bit shaping functions in the watermark generator 101 so that the masking thresholds are indeed fulfilled. In fact, as already discussed, the bit shaping function normally extends for a time interval larger than T_b. Therefore, multiplying the correct amplitude $\gamma(i, j)$ which

fulfills the masking threshold at point i, j does not necessarily fulfill the requirements at point i, j-1. This is particularly crucial at strong onsets, as a preecho becomes audible. Another situation which needs to be avoided is the unfortunate superposition of the tails of different bits which might lead to an audible watermark. Therefore, block 902 analyzes the signal generated by the watermark generator to check whether the thresholds have been fulfilled. If not, it modifies the amplitudes $\gamma(i, j)$ accordingly.

This concludes the encoder side. The following sections deal with the processing steps carried out at the receiver (also designated as watermark decoder).

3.3 The Analysis Module 203

The analysis module 203 is the first step (or block) of the watermark extraction process. Its purpose is to transform the watermarked audio signal 200a back into N_f bit streams $\hat{b}_i(j)$ (also designated with 204), one for each spectral subband i. These are further processed by the synchronization module 201 and the watermark extractor 202, as discussed in Sections 3.4 and 3.5, respectively. Note that the $\hat{b}_i(j)$ are soft bit streams, i.e., they can take, for example, any real value and no hard decision on the bit is made yet.

The analysis module consists of three parts which are depicted in Figure 16: The analysis filter bank 1600, the amplitude normalization block 1604 and the differential decoding 1608.

3.3.1 Analysis filter bank 1600

The watermarked audio signal is transformed into the time-frequency domain by the analysis filter bank 1600 which is shown in detail in Figure 10a. The input of the filter bank is the received watermarked audio signal r(t). Its output are the complex coefficients $b_i^{AFB}(j)$ for the i-th branch or subband at time instant j. These values contain information about the amplitude and the phase of the signal at center frequency f_i and time $j \cdot Tb$.

The filter bank 1600 consists of N_f branches, one for each spectral subband i. Each branch splits up into an upper subbranch for the in-phase component and a lower subbranch for the quadrature component of the subband i. Although the modulation at the watermark generator and thus the watermarked audio signal are purely real-valued, the complex-

valued analysis of the signal at the receiver is needed because rotations of the modulation constellation introduced by the channel and by synchronization misalignments are not known at the receiver. In the following we consider the i-th branch of the filter bank. By combining the in-phase and the quadrature subbranch, we can define the complex-valued baseband signal $b_i^{\mathrm{AFB}}(t)$ as

$$b_i^{\mathrm{AFB}}(t) = r(t) \cdot e^{-j2\pi f_i t} * g_i^{\mathrm{R}}(t) \tag{10}$$

where $*$ indicates convolution and $g_i^{\mathrm{R}}(t)$ is the impulse response of the receiver lowpass filter of subband i. Usually $g_i^{\mathrm{R}}(t)$i (t) is equal to the baseband bit forming function $g_i^{\mathrm{T}}(t)$of subband i in the modulator 307 in order to fulfill the matched filter condition, but other impulse responses are possible as well.

In order to obtain the coefficients $b_i^{\mathrm{AFB}}(j)$ with rate 1=T_b, the continuous output $b_i^{\mathrm{AFB}}(t)$ must be sampled. If the correct timing of the bits was known by the receiver, sampling with rate 1=T_b would be sufficient. However, as the bit synchronization is not known yet, sampling is carried out with rate N_{os}/T_b where N_{os} is the analysis filter bank oversampling factor. By choosing N_{os} sufficiently large (e.g. N_{os} = 4), we can assure that at least one sampling cycle is close enough to the ideal bit synchronization. The decision on the best oversampling layer is made during the synchronization process, so all the oversampled data is kept until then. This process is described in detail in Section 3.4.

At the output of the i-th branch we have the coefficients $b_i^{\mathrm{AFB}}(j,k)$, where j indicates the bit number or time instant and k indicates the oversampling position within this single bit, where k = 1; 2; …., N_{os}.

Figure 10b gives an exemplary overview of the location of the coefficients on the time-frequency plane. The oversampling factor is N_{os} = 2. The height and the width of the rectangles indicate respectively the bandwidth and the time interval of the part of the signal that is represented by the corresponding coefficient $b_i^{\mathrm{AFB}}(j,k)$.

If the subband frequencies f_i are chosen as multiples of a certain interval Δf the analysis filter bank can be efficiently implemented using the Fast Fourier Transform (FFT).

3.3.2 Amplitude normalization 1604

Without loss of generality and to simplify the description, we assume that the bit synchronization is known and that $N_{os} = 1$ in the following. That is, we have complex coeffcients $b_i^{\mathrm{AFB}}(j)$at the input of the normalization block 1604. As no channel state information is available at the receiver (i.e., the propagation channel in unknown), an equal gain combining (EGC) scheme is used. Due to the time and frequency dispersive channel, the energy of the sent bit $b_i(j)$ is not only found around the center frequency f_i and time instant j, but also at adjacent frequencies and time instants. Therefore, for a more precise weighting, additional coefficients at frequencies $f_i \pm n\,\Delta f$ are calculated and used for normalization of coefficient $b_i^{\mathrm{AFB}}(j)$. If $n = 1$ we have, for example,

$$b_i^{\mathrm{norm}}(j) = \frac{b_i^{\mathrm{AFB}}(j)}{\sqrt{1/3 \cdot (|b_i^{\mathrm{AFB}}(j)|^2 + |b_{i-\Delta f}^{\mathrm{AFB}}(j)|^2 + |b_{i+\Delta f}^{\mathrm{AFB}}(j)|^2)}} \tag{11}$$

The normalization for $n > 1$ is a straightforward extension of the formula above. In the same fashion we can also choose to normalize the soft bits by considering more than one time instant. The normalization is carried out for each subband i and each time instant j. The actual combining of the EGC is done at later steps of the extraction process.

3.3.3 Differential decoding 1608

At the input of the differential decoding block 1608 we have amplitude normalized complex coefficients $b_i^{\mathrm{norm}}(j)$which contain information about the phase of the signal components at frequency f_i and time instant j. As the bits are differentially encoded at the transmitter, the inverse operation must be performed here. The soft bits $\widehat{b}_i(j)$are obtained by first calculating the difference in phase of two consecutive coefficients and then taking the real part:

$$\widehat{b}_i(j) = \mathrm{Re}\{b_i^{\mathrm{norm}}(j) \cdot b_i^{\mathrm{norm}^*}(j-1)\} \tag{12}$$

$$= \mathrm{Re}\{|b_i^{\mathrm{norm}}(j)| \cdot |b_i^{\mathrm{norm}}(j-1)| \cdot e^{j(\varphi_j - \varphi_{j-1})}\} \tag{13}$$

This has to be carried out separately for each subband because the channel normally introduces different phase rotations in each subband.

3.4 The Synchronization Module 201

The synchronization module's task is to find the temporal alignment of the watermark. The problem of synchronizing the decoder to the encoded data is twofold. In a first step, the analysis filterbank must be aligned with the encoded data, namely the bit shaping functions $g_i^{\mathrm{T}}(t)$ used in the synthesis in the modulator must be aligned with the filters $g_i^{\mathrm{R}}(t)$ used for the analysis. This problem is illustrated in Figure 12a, where the analysis filters are identical to the synthesis ones. At the top, three bits are visible. For simplicity, the waveforms for all three bits are not scaled. The temporal offset between different bits is T_b. The bottom part illustrates the synchronization issue at the decoder: the filter can be applied at different time instants, however, only the position marked in red (curve 1299a) is correct and allows to extract the first bit with the best signal to noise ratio SNR and signal to interference ratio SIR. In fact, an incorrect alignment would lead to a degradation of both SNR and SIR. We refer to this first alignment issue as "bit synchronization". Once the bit synchronization has been achieved, bits can be extracted optimally. However, to correctly decode a message, it is necessary to know at which bit a new message starts. This issue is illustrated in Figure 12b and is referred to as message synchronization. In the stream of decoded bits only the starting position marked in red (position 1299b) is correct and allows to decode the k-th message.

We first address the message synchronization only. The synchronization signature, as explained in Section 3.1, is composed of Ns sequences in a predetermined order which are embedded continuously and periodically in the watermark. The synchronization module is capable of retrieving the temporal alignment of the synchronization sequences. Depending on the size N_s we can distinguish between two modes of operation, which are depicted in Figure 12c and 12d, respectively.

In the full message synchronization mode (Fig. 12c) we have $N_s = N_m/R_c$. For simplicity in the figure we assume $N_s = N_m/R_c = 6$ and no time spreading, i.e., $N_t = 1$. The synchronization signature used, for illustration purposes, is shown beneath the messages. In reality, they are modulated depending on the coded bits and frequency spreading sequences, as explained in Section 3.1. In this mode, the periodicity of the synchronization signature is identical to the one of the messages. The synchronization module therefore can identify the beginning of each message by finding the temporal alignment of the

synchronization signature. We refer to the temporal positions at which a new synchronization signature starts as synchronization hits. The synchronization hits are then passed to the watermark extractor 202.

The second possible mode, the partial message synchronization mode (Fig. 12d), is depicted in Figure 12d. In this case we have $N_s < N_m = R_c$. In the figure we have taken $N_s = 3$, so that the three synchronization sequences are repeated twice for each message. Please note that the periodicity of the messages does not have to be multiple of the periodicity of the synchronization signature. In this mode of operation, not all synchronization hits correspond to the beginning of a message. The synchronization module has no means of distinguishing between hits and this task is given to the watermark extractor 202.

The processing blocks of the synchronization module are depicted in Figures 11a and 11b. The synchronization module carries out the bit synchronization and the message synchronization (either full or partial) at once by analyzing the output of the synchronization signature correlator 1201. The data in time/frequency domain 204 is provided by the analysis module. As the bit synchronization is not yet available, block 203 oversamples the data with factor N_{os}, as described in Section 3.3. An illustration of the input data is given in Figure 12e. For this example we have taken $N_{os} = 4$, $N_t = 2$, and $N_s = 3$. In other words, the synchronization signature consists of 3 sequences (denoted with a, b, and c). The time spreading, in this case with spreading sequence $c_t = [1\ 1]^T$, simply repeats each bit twice in time domain. The exact synchronization hits are denoted with arrows and correspond to the beginning of each synchronization signature. The period of the synchronization signature is $N_t \cdot N_{os} \cdot N_s = N_{sbl}$ which is $2 \cdot 4 \cdot 3 = 24$, for example. Due to the periodicity of the synchronization signature, the synchronization signature correlator (1201) arbitrarily divides the time axis in blocks, called search blocks, of size N_{sbl}, whose subscript stands for search block length. Every search block must contain (or typically contains) one synchronization hit as depicted in Figure 12f. Each of the N_{sbl} bits is a candidate synchronization hit. Block 1201's task is to compute a likelihood measure for each of candidate bit of each block. This information is then passed to block 1204 which computes the synchronization hits.

3.4.1 The synchronization signature correlator 1201

For each of the N_{sbl} candidate synchronization positions the synchronization signature correlator computes a likelihood measure, the latter is larger the more probable it is that the

temporal alignment (both bit and partial or full message synchronization) has been found. The processing steps are depicted in Figure 12g.

Accordingly, a sequence 1201aof likelihood values, associated with different positional choices, may be obtained.

Block 1301 carries out the temporal despreading, i.e., multiplies every N_t bits with the temporal spreading sequence c_t and then sums them. This is carried out for each of the N_f frequency subbands. Figure 13a shows an example. We take the same parameters as described in the previous section, namely $N_{os} = 4$, $N_t = 2$, and $N_s = 3$. The candidate synchronization position is marked. From that bit, with N_{os} offset, $N_t \cdot N_s$ are taken by block 1301 and time despread with sequence c_t, so that Ns bits are left.

In block 1302 the bits are multiplied element-wise with the N_s spreading sequences (see Figure 13b).

In block 1303 the frequency despreading is carried out, namely, each bit is multiplied with the spreading sequence c_f and then summed along frequency.

At this point, if the synchronization position were correct, we would have N_s decoded bits. As the bits are not known to the receiver, block 1304 computes the likelihood measure by taking the absolute values of the N_s values and sums.

The output of block 1304 is in principle a non coherent correlator which looks for the synchronization signature. In fact, when choosing a small N_s, namely the partial message synchronization mode, it is possible to use synchronization sequences (e.g. **a**, **b**, **c**) which are mutually orthogonal. In doing so, when the correlator is not correctly aligned with the signature, its output will be very small, ideally zero. When using the full message synchronization mode it is advised to use as many orthogonal synchronization sequences as possible, and then create a signature by carefully choosing the order in which they are used. In this case, the same theory can be applied as when looking for spreading sequences with good auto correlation functions. When the correlator is only slightly misaligned, then the output of the correlator will not be zero even in the ideal case, but anyway will be smaller compared to the perfect alignment, as the analysis filters cannot capture the signal energy optimally.

3.4.2 Synchronization hits computation 1204

This block analyzes the output of the synchronization signature correlator to decide where the synchronization positions are. Since the system is fairly robust against misalignments of up to $T_b/4$ and the T_b is normally taken around 40 ms, it is possible to integrate the output of 1201 over time to achieve a more stable synchronization. A possible implementation of this is given by an IIR filter applied along time with a exponentially decaying impulse response. Alternatively, a traditional FIR moving average filter can be applied. Once the averaging has been carried out, a second correlation along different $N_f \cdot N_s$ is carried out ("different positional choice"). In fact, we want to exploit the information that the autocorrelation function of the synchronization function is known. This corresponds to a Maximum Likelihood estimator. The idea is shown in Figure 13c. The curve shows the output of block 1201 after temporal integration. One possibility to determine the synchronization hit is simply to find the maximum of this function. In Figure 13d we see the same function (in black) filtered with the autocorrelation function of the synchronization signature. The resulting function is plotted in red. In this case the maximum is more pronounced and gives us the position of the synchronization hit. The two methods are fairly similar for high SNR but the second method performs much better in lower SNR regimes. Once the synchronization hits have been found, they are passed to the watermark extractor 202 which decodes the data.

In some embodiments, in order to obtain a robust synchronization signal, synchronization is performed in partial message synchronization mode with short synchronization signatures. For this reason many decodings have to be done, increasing the risk of false positive message detections. To prevent this, in some embodiments signaling sequences may be inserted into the messages with a lower bit rate as a consequence.

This approach is a solution to the problem arising from a sync signature shorter than the message, which is already addressed in the above discussion of the enhanced synchronization. In this case, the decoder doesn't know where a new message starts and attempts to decode at several synchronization points. To distinguish between legitimate messages and false positives, in some embodiments a signaling word is used (i.e. payload is sacrified to embed a known control sequence). In some embodiments, a plausibility check is used (alternatively or in addition) to distinguish between legitimate messages and false positives.

3.5 The Watermark Extractor 202

The parts constituting the watermark extractor 202 are depicted in Figure 14. This has two inputs, namely 204 and 205 from blocks 203 and 201, respectively. The synchronization module 201 (see Section 3.4) provides synchronization timestamps, i.e., the positions in time domain at which a candidate message starts. More details on this matter are given in Section 3.4. The analysis filterbank block 203, on the other hand, provides the data in time/frequency domain ready to be decoded.

The first processing step, the data selection block 1501, selects from the input 204 the part identified as a candidate message to be decoded. Figure 15 shows this procedure graphically. The input 204 consists of N_f streams of real values. Since the time alignment is not known to the decoder a priori, the analysis block 203 carries out a frequency analysis with a rate higher than $1/T_b$ Hz (oversampling). In Figure 15 we have used an oversampling factor of 4, namely, 4 vectors of size $N_f \times 1$ are output every T_b seconds. When the synchronization block 201 identifies a candidate message, it delivers a timestamp 205 indicating the starting point of a candidate message. The selection block 1501 selects the information required for the decoding, namely a matrix of size $N_f \times N_m/R_c$. This matrix 1501a is given to block 1502 for further processing.

Blocks 1502, 1503, and 1504 carry out the same operations of blocks 1301, 1302, and 1303 explained in Section 3.4.

An alternative embodiment of the invention consists in avoiding the computations done in 1502-1504 by letting the synchronization module deliver also the data to be decoded. Conceptually it is a detail. From the implementation point of view, it is just a matter of how the buffers are realized. In general, redoing the computations allows us to have smaller buffers.

The channel decoder 1505 carries out the inverse operation of block 302. If channel encoder, in a possible embodiment of this module, consisted of a convolutional encoder together with an interleaver, then the channel decoder would perform the deinterleaving and the convolutional decoding, e.g., with the well known Viterbi algorithm. At the output of this block we have N_m bits, i.e., a candidate message.

Block 1506, the signaling and plausibility block, decides whether the input candidate message is indeed a message or not. To do so, different strategies are possible.

The basic idea is to use a signaling word (like a CRC sequence) to distinguish between true and false messages. This however reduces the number of bits available as payload.

Alternatively we can use plausibility checks. If the messages for instance contain a timestamp, consecutive messages must have consecutive timestamps. If a decoded message possesses a timestamp which is not the correct order, we can discard it.

When a message has been correctly detected the system may choose to apply the look ahead and/or look back mechanisms. We assume that both bit and message synchronization have been achieved. Assuming that the user is not zapping, the system "looks back" in time and attempts to decode the past messages (if not decoded already) using the same synchronization point (look back approach). This is particularly useful when the system starts. Moreover, in bad conditions, it might take 2 messages to achieve synchronization. In this case, the first message has no chance. With the look back option we can save "good" messages which have not been received only due to back synchronization. The look ahead is the same but works in the future. If we have a message now we know where the next message should be, and we can attempt to decode it anyhow.

3.6. Synchronization Details

For the encoding of a payload, for example, a Viterbi algorithm may be used. Fig. 18a shows a graphical representation of a payload 1810, a Viterbi termination sequence 1820, a Viterbi encoded payload 1830 and a repetition-coded version 1840 of the Viterbi-coded payload. For example, the payload length may be 34 bits and the Viterbi termination sequence may comprise 6 bits. If, for example a Viterbi code rate of 1/7 may be used the Viterbi-coded payload may comprise (34+6)*7=280 bits. Further, by using a repetition coding of 1/2, the repetition coded version 1840 of the Viterbi-encoded payload 1830 may comprise 280*2=560 bits. In this example, considering a bit time interval of 42.66 ms, the message length would be 23.9 s. The signal may be embedded with, for example, 9 subcarriers (e.g. placed according to the critical bands) from 1.5 to 6 kHz as indicated by the frequency spectrum shown in Fig. 18b. Alternatively, also another number of subcarriers (e.g. 4, 6, 12, 15 or a number between 2 and 20) within a frequency range between 0 and 20 kHz maybe used.

Fig. 19 shows a schematic illustration of the basic concept 1900 for the synchronization, also called ABC synch. It shows a schematic illustration of an uncoded messages 1910, a coded message 1920 and a synchronization sequence (synch sequence) 1930 as well as the application of the synch to several messages 1920 following each other.

The synchronization sequence or synch sequence mentioned in connection with the explanation of this synchronization concept (shown in Fig. 19 – 23) may be equal to the synchronization signature mentioned before.

Further, Fig. 20 shows a schematic illustration of the synchronization found by correlating with the synch sequence. If the synchronization sequence 1930 is shorter than the message, more than one synchronization point 1940 (or alignment time block) may be found within a single message. In the example shown in Fig. 20, 4 synchronization points are found within each message. Therefore, for each synchronization found, a Viterbi decoder (a Viterbi decoding sequence) may be started. In this way, for each synchronization point 1940 a message 2110 may be obtained, as indicated in Fig. 21.

Based on these messages the true messages 2210 may be identified by means of a CRC sequence (cyclic redundancy check sequence) and/or a plausibility check, as shown in Fig. 22.

The CRC detection (cyclic redundancy check detection) may use a known sequence to identify true messages from false positive. Fig. 23 shows an example for a CRC sequence added to the end of a payload.

The probability of false positive (a message generated based on a wrong synchronization point) may depend on the length of the CRC sequence and the number of Viterbi decoders (number of synchronization points within a single message) started. To increase the length of the payload without increasing the probability of false positive a plausibility may be exploited (plausibility test) or the length of the synchronization sequence (synchronization signature) may be increased.

4. Concepts and Advantages

In the following, some aspects of the above discussed system will be described, which are considered as being innovative. Also, the relation of those aspects to the state-of-the-art technologies will be discussed.

4.1. Continuous synchronization

Some embodiments allow for a continuous synchronization. The synchronization signal, which we denote as synchronization signature, is embedded continuously and parallel to the data via multiplication with sequences (also designated as synchronization spread sequences) known to both transmit and receive side.

Some conventional systems use special symbols (other than the ones used for the data), while some embodiments according to the invention do not use such special symbols. Other classical methods consist of embedding a known sequence of bits (preamble) time-multiplexed with the data, or embedding a signal frequency-multiplexed with the data.

However, it has been found that using dedicated sub-bands for synchronization is undesired, as the channel might have notches at those frequencies, making the synchronization unreliable. Compared to the other methods, in which a preamble or a special symbol is time-multiplexed with the data, the method described herein is more advantageous as the method described herein allows to track changes in the synchronization (due e.g. to movement) continuously.

Furthermore, the energy of the watermark signal is unchanged (e.g. by the multiplicative introduction of the watermark into the spread information representation), and the synchronization can be designed independent from the psychoacoustical model and data rate. The length in time of the synchronization signature, which determines the robustness of the synchronization, can be designed at will completely independent of the data rate.

Another classical method consists of embedding a synchronization sequence code-multiplexed with the data. When compared to this classical method, the advantage of the method described herein is that the energy of the data does not represent an interfering factor in the computation of the correlation, bringing more robustness. Furthermore, when using code-multiplexing, the number of orthogonal sequences available for the synchronization is reduced as some are necessary for the data.

To summarize, the continuous synchronization approach described herein brings along a large number of advantages over the conventional concepts.

However, in some embodiments according to the invention, a different synchronization concept may be applied.

4.2. 2D spreading

Some embodiments of the proposed system carry out spreading in both time and frequency domain, i.e. a 2-dimensional spreading (briefly designated as 2D-spreading). It has been found that this is advantageous with respect to 1D systems as the bit error rate can be further reduced by adding redundance in e.g. time domain.

However, in some embodiments according to the invention, a different spreading concept may be applied.

4.3. Differential encoding and Differential decoding

In some embodiments according to the invention, an increased robustness against movement and frequency mismatch of the local oscillators (when compared to conventional systems) is brought by the differential modulation. It has been found that in fact, the Doppler effect (movement) and frequency mismatches lead to a rotation of the BPSK constellation (in other words, a rotation on the complex plane of the bits). In some embodiments, the detrimental effects of such a rotation of the BPSK constellation (or any other appropriate modulation constellation) are avoided by using a differential encoding or differential decoding.

However, in some embodiments according to the invention, a different encoding concept or decoding concept may be applied. Also, in some cases, the differential encoding may be omitted.

4.4. Bit shaping

In some embodiments according to the invention, bit shaping brings along a significant improvement of the system performance, because the reliability of the detection can be increased using a filter adapted to the bit shaping.

In accordance with some embodiments, the usage of bit shaping with respect to watermarking brings along improved reliability of the watermarking process. It has been found that particularly good results can be obtained if the bit shaping function is longer than the bit interval.

However, in some embodiments according to the invention, a different bit shaping concept may be applied. Also, in some cases, the bit shaping may be omitted.

4.5. Interactive between Psychoacoustic Model (PAM) and Filter Bank (FB) synthesis

In some embodiments, the psychoacoustical model interacts with the modulator to fine tune the amplitudes which multiply the bits.

However, in some other embodiments, this interaction may be omitted.

4.6. Look ahead and look back features

In some embodiments, so called "Look back" and "look ahead" approaches are applied.

In the following, these concepts will be briefly summarized. When a message is correctly decoded, it is assumed that synchronization has been achieved. Assuming that the user is not zapping, in some embodiments a look back in time is performed and it is tried to decode the past messages (if not decoded already) using the same synchronization point (look back approach). This is particularly useful when the system starts.

In bad conditions, it might take 2 messages to achieve synchronization. In this case, the first message has no chance in conventional systems. With the look back option, which is used in some embodiments of the invention, it is possible to save (or decode) "good" messages which have not been received only due to back synchronization.

The look ahead is the same but works in the future. If I have a message now I know where my next message should be, and I can try to decode it anyhow. Accordingly, overlapping messages can be decoded.

However, in some embodiments according to the invention, the look ahead feature and/or the look back feature may be omitted.

4.7. Increased synchronization robustness

In some embodiments, in order to obtain a robust synchronization signal, synchronization is performed in partial message synchronization mode with short synchronization

signatures. For this reason many decodings have to be done, increasing the risk of false positive message detections. To prevent this, in some embodiments signaling sequences may be inserted into the messages with a lower bit rate as a consequence.

However, in some embodiments according to the invention, a different concept for improving the synchronization robustness may be applied. Also, in some cases, the usage of any concepts for increasing the synchronization robustness may be omitted.

4.8. Other enhancements

In the following, some other general enhancements of the above described system with respect to background art will be put forward and discussed:

1. lower computational complexity

2. better audio quality due to the better psychoacoustical model

3. more robustness in reverberant environments due to the narrowband multicarrier signals

4. an SNR estimation is avoided in some embodiments. This allows for better robustness, especially in low SNR regimes.

Some embodiments according to the invention are better than conventional systems, which use very narrow bandwidths of, for example, 8Hz for the following reasons:

1. 8 Hz bandwidths (or a similar very narrow bandwidth) requires very long time symbols because the psychoacoustical model allows very little energy to make it inaudible;

2. 8 Hz (or a similar very narrow bandwidth) makes it sensitive against time varying Doppler spectra. Accordingly, such a narrow band system is typically not good enough if implemented, e.g., in a watch.

Some embodiments according to the invention are better than other technologies for the following reasons:

1. Techniques which input an echo fail completely in reverberant rooms. In contrast, in some embodiments of the invention, the introduction of an echo is avoided.

2. Techniques which use only time spreading have longer message duration in comparison embodiments of the above described system in which a two-dimensional spreading, for example both in time and in frequency, is used.

Some embodiments according to the invention are better than the system described in DE 196 40 814, because one of more of the following disadvantages of the system according to said document are overcome:

- the complexity in the decoder according to DE 196 40 814 is very high, a filter of length 2N with N = 128 is used
- the system according to DE 196 40 814 comprises a long message duration
- in the system according to DE 196 40 814 spreading only in time domain with relatively high spreading gain (e.g. 128)
- in the system according to DE 196 40 814 the signal is generated in time domain, transformed to spectral domain, weighted, transformed back to time domain, and superposed to audio, which makes the system very complex

5. Applications

The invention comprises a method to modify an audio signal in order to hide digital data and a corresponding decoder capable of retrieving this information while the perceived quality of the modified audio signal remains indistinguishable to the one of the original.

Examples of possible applications of the invention are given in the following:

1. Broadcast monitoring: a watermark containing information on e.g. the station and time is hidden in the audio signal of radio or television programs. Decoders, incorporated in small devices worn by test subjects, are capable to retrieve the watermark, and thus collect valuable information for advertisements agencies, namely who watched which program and when.

2. Auditing: a watermark can be hidden in, e.g., advertisements. By automatically monitoring the transmissions of a certain station it is then possible to know when exactly the ad was broadcast. In a similar fashion it is possible to retrieve statistical information

about the programming schedules of different radios, for instance, how often a certain music piece is played, etc.

3. Metadata embedding: the proposed method can be used to hide digital information about the music piece or program, for instance the name and author of the piece or the duration of the program etc.

6. Implementation Alternatives

Although some aspects have been described in the context of an apparatus, it is clear that these aspects also represent a description of the corresponding method, where a block or device corresponds to a method step or a feature of a method step. Analogously, aspects described in the context of a method step also represent a description of a corresponding block or item or feature of a corresponding apparatus. Some or all of the method steps may be executed by (or using) a hardware apparatus, like for example, a microprocessor, a programmable computer or an electronic circuit. In some embodiments, some one or more of the most important method steps may be executed by such an apparatus.

The inventive encoded watermark signal, or an audio signal into which the watermark signal is embedded, can be stored on a digital storage medium or can be transmitted on a transmission medium such as a wireless transmission medium or a wired transmission medium such as the Internet.

Depending on certain implementation requirements, embodiments of the invention can be implemented in hardware or in software. The implementation can be performed using a digital storage medium, for example a floppy disk, a DVD, a Blue-Ray, a CD, a ROM, a PROM, an EPROM, an EEPROM or a FLASH memory, having electronically readable control signals stored thereon, which cooperate (or are capable of cooperating) with a programmable computer system such that the respective method is performed. Therefore, the digital storage medium may be computer readable.

Some embodiments according to the invention comprise a data carrier having electronically readable control signals, which are capable of cooperating with a programmable computer system, such that one of the methods described herein is performed.

Generally, embodiments of the present invention can be implemented as a computer program product with a program code, the program code being operative for performing one of the methods when the computer program product runs on a computer. The program code may for example be stored on a machine readable carrier.

Other embodiments comprise the computer program for performing one of the methods described herein, stored on a machine readable carrier.

In other words, an embodiment of the inventive method is, therefore, a computer program having a program code for performing one of the methods described herein, when the computer program runs on a computer.

A further embodiment of the inventive methods is, therefore, a data carrier (or a digital storage medium, or a computer-readable medium) comprising, recorded thereon, the computer program for performing one of the methods described herein.

A further embodiment of the inventive method is, therefore, a data stream or a sequence of signals representing the computer program for performing one of the methods described herein. The data stream or the sequence of signals may for example be configured to be transferred via a data communication connection, for example via the Internet.

A further embodiment comprises a processing means, for example a computer, or a programmable logic device, configured to or adapted to perform one of the methods described herein.

A further embodiment comprises a computer having installed thereon the computer program for performing one of the methods described herein.

In some embodiments, a programmable logic device (for example a field programmable gate array) may be used to perform some or all of the functionalities of the methods described herein. In some embodiments, a field programmable gate array may cooperate with a microprocessor in order to perform one of the methods described herein. Generally, the methods are preferably performed by any hardware apparatus.

The above described embodiments are merely illustrative for the principles of the present invention. It is understood that modifications and variations of the arrangements and the details described herein will be apparent to others skilled in the art. It is the intent, therefore, to be limited only by the scope of the impending patent claims and not by the

specific details presented by way of description and explanation of the embodiments herein.

Claims

1. A watermark generator (101;2400) for providing a watermark signal (101b, $s_i(t)$; 2420) in dependence on binary message data (101a, m; 2410), the watermark generator comprising:

an information spreader (303; 2430) configured to spread an information unit to a plurality of time-frequency-domain values, to obtain a spread information representation (303a, $\mathbf{c_f} \cdot \mathbf{m}$; 2432);

a synchronization inserter configured to multiplicatively combine the spread information representation (303a, $\mathbf{c_f} \cdot \mathbf{m}$; 2432) with a synchronization sequence (**S**; **a**, **b**, **c**; 2442) to obtain a combined information-synchronization representation (304a, $S \circ \mathbf{c_f} \cdot \mathbf{m}$; 2444); and

a watermark signal provider (305,306,307; 2450) configured to provide the watermark signal (101b, $s_i(t)$; 2420) on the basis of the combined information-synchronization representation (304a, $S \circ \mathbf{c_f} \cdot \mathbf{m}$; 2444).

2. The watermark generator according to claim 1, wherein the information spreader is configured to frequency-spread a bit of the binary message data in accordance with a bit spread sequence (**a**, **b**, **c**), such that, in the spread information representation, a set of time-frequency-domain values associated with a plurality of different frequencies is associated to the bit; and

wherein the synchronization sequence comprises a plurality of synchronization spread sequences (**a**, **b**, **c**), the synchronization spread sequences defining a plurality of synchronization patterns spread in frequency.

3. The watermark generator according to claim 2, wherein the synchronization inserter is configured to multiplicatively combine different time portions of the spread information representation, which different time portions represent different bits of the binary message data, with different synchronization spread sequences (**a**, **b**, **c**), such that the different synchronization spread sequences are combined with different bit spread sequences or with differently scaled versions of a given bit spread sequence.

4. The watermark generator according to claim 2 or claim 3, wherein the synchronization inserter is configured to multiplicatively combine the spread information representation with a periodic sequence (**S**) of synchronization spread sequences (**a**, **b**, **c**, **a**, **b**, **c**); and

wherein a periodicity of the sequence of synchronization spread sequences is equal to a number of bits of a binary message, or

wherein a periodicity of the sequence of synchronization spread sequences is shorter than the number of bits of the binary message.

5. The watermark generator according to one of claims 2 to 4, wherein the synchronization spread sequences (**a**, **b**, **c**) are chosen to be orthogonal with respect to each other.

6. The watermark generator according to one of claims 1 to 5, wherein the synchronization inserter is configured to multiplicatively combine each time portion of the spread information representation with at least one synchronization spread sequence (**a**, **b**, **c**), such that each time portion of the combined information-synchronization representation comprises a plurality of values which are dependent on a bit of the binary message data and which also represent a synchronization information.

7. The watermark generator according to one of claims 1 to 6, wherein the information spreader is configured to selectively spread a given bit onto a first bit representation, which is a positive multiple of a bit spread sequence ($\mathbf{c}_f$) , or onto a second bit representation, which is a negative multiple of the bit spread sequence ($\mathbf{c}_f$), in dependence on the value of the given bit.

8. The watermark generator according to one of claims 1 to 7, wherein the information spreader is configured to obtain the spread information representation **R** according to

$$\mathbf{R}=\mathbf{c}_f \cdot \mathbf{m},$$

wherein $\mathbf{c}_f$ is a vector of size $N_f \times 1$ representing a bit spread sequence of frequency spreading width N_f,

wherein **m** is a vector of size $1 \times N_{mc}$ representing N_{mc} bits of the binary message data, wherein binary values of the bits are represented by entries of the vector **m** having a different sign;

wherein the synchronization inserter is configured to combine the spread information with a synchronization sequence **S** according to

$$\mathbf{C}=\mathbf{S} \circ \mathbf{R},$$

to obtain the combined information-synchronization representation **C**,

wherein **S** designates a matrix of size $N_f \times N_{mc}$ representing a sequence of N_{mc} synchronization spread sequences (**a**, **b**, **c**); and

wherein $\circ$ designates a Schur element-wise product.

9. The watermark generator according to one of claims 1 to 8, wherein the information spreader is configured to spread the information unit of the binary message data in a first spreading direction, and wherein the watermark generator is configured to spread information units of the combined information-synchronization representation in a second spreading direction.

10. The watermark generator according to one of claims 1 to 9, wherein the synchronization inserter is configured to multiplicatively combine the spread information representation with the synchronization spread sequence in an element-wise manner, to obtain the combined information-synchronization representation.

11. A watermark decoder (200; 2500) for providing binary message data (202a, **m**; 2520) in dependence on a watermarked signal (200a, r(t); 2510), the watermark decoder comprising:

a time-frequency-domain representation provider (203; 2530) configured to provide a time-frequency-domain representation (204; 2532) of the watermarked signal; and

a synchronization determinator (201; 2540) configured to find a temporal alignment of watermark information within the time-frequency-domain representation of the watermarked signal,

wherein the synchronization determinator is configured to compute a first sum of result values of an element-wise multiplication between a first subset (1301aa) of values of the time-frequency-domain representation, values of a first synchronization despreading sequence (**a**) and values of a bit despreading sequence ($\mathbf{c}_f$), and to compute a second sum of result values of an element-wise multiplication between a second subset (1301ab) of values of the time-frequency-domain representation, values of a second synchronization despreading sequence (**b**) and values of the bit despreading sequence ($\mathbf{c}_f$), and

wherein the synchronization determinator is configured to sum absolute values of the first sum (1303aa) of result values and of the second sum (1303ab) of result values, to obtain a likelihood measure representing a likelihood that the first and second subsets of values of the time-frequency-domain representation carry a synchronization information, which is in temporal alignment with the synchronization spreading sequences; and

wherein the synchronization determinator is configured to determine the synchronization in dependence on the likelihood measure.

12. The watermark decoder according to claim 11, wherein the first subset of values of the time-frequency-domain representation comprises values of the time-frequency-domain representation associated with a first time portion of the time-frequency-domain representation, and

wherein the second subset of values of the time-frequency-domain representation comprises values of the time-frequency-domain representation associated with a second time portion of the time-frequency-domain representation, and

wherein the first synchronization despreading sequence (**a**) and the second synchronization despreading sequence (**b**) are orthogonal.

13. The watermark decoder according to claim 11 or claim 12, wherein the synchronization determinator is configured to receive, as the time-frequency-domain representation, are hard bit information or a soft bit information ($\hat{b}_i(j)$); and

wherein the synchronization determinator is configured to apply the first synchronization despreading sequence (**a**) to the first subset of values,

to reverse a first synchronization embedding operation, to apply the second synchronization despreading sequence (**b**) to the second subset of values, to reverse a synchronization embedding operation,

and to apply the bit despreading sequence ($\mathbf{c}_f$) to reverse a bit spreading operation.

14. The watermark decoder according to one of claims 11 to 13, wherein the synchronization determinator is configured to compute likelihood measures for a plurality of positional choices of the first and second subset with reference to the time-frequency-domain representation,

to obtain a sequence of likelihood measures associated with the different positioned choices of the first and second subset; and

wherein the synchronization determinator is configured to correlate the sequence of likelihood measures with a reference sequence, in order to determine a synchronization hit, or

to filter the sequence of likelihood measures with a matched filter, adapted to an expected sequence of likelihood values, in order to determine a synchronization hit.

15. The watermark decoder according to one of claims 11 to 14, wherein the synchronization determinator is configured to perform a temporal despreading, to obtain the time-frequency-domain representation.

16. The watermark decoder according to one of claims 11 to 15, wherein the synchronization determinator is configured to decode a message using a plurality of candidate synchronization points, and wherein the synchronization determinator is configured to evaluate whether an expected signaling word is present in decoded data in order to select a correct synchronization point out of the plurality of candidate synchronization points.

17. The watermark decoder according to one of claims 11 to 15, wherein the synchronization determinator is configured to decode a message using a plurality of candidate synchronization points, and wherein the synchronization determinator is

configured perform a plausibility check in order to select a correct synchronization point out of the plurality of candidate synchronization points.

18. A method for providing a watermark signal in dependence on binary message data, the method comprising:

spreading an information unit to a plurality of time-frequency-domain values, to obtain a spread information representation;

multiplicatively combining the spread information representation with a synchronization sequence (**a**, **b**, **c**) to obtain a combined information-synchronization representation; and

providing the watermark signal on the basis of the combined information-synchronization representation.

19. A method for providing binary message data in dependence on a watermarked signal, the method comprising:

providing a time-frequency-domain representation of the watermarked signal; and

finding a temporal alignment of watermark information within the time-frequency-domain representation of the watermarked signal,

wherein finding the temporal alignment of the watermark information comprises:

computing a first sum of result values of an element-wise multiplication between a first subset of values of the time-frequency-domain representation, values of a first synchronization despreading sequence (**a**) and values of a bit despreading sequence ($\mathbf{c}_f$);

computing a second sum of result values of an element-wise multiplication between a second subset of values of the time-frequency-domain representation, values of a second synchronization despreading sequence (**b**) and values of the bit despreading sequence ($\mathbf{c}_f$);

summing absolute values of the first sum of result values and of the second sum of result values, to obtain a likelihood measure representing a likelihood that the first and second subsets of values of the time-frequency-domain representation carry a synchronization information, which is in temporal alignment with the synchronization spreading sequence (**a**, **b**, **c**); and

determining the synchronization dependent on the likelihood measure.

20. A computer program for performing the method according to claim 18 or claim 19, when the computer program runs on a computer.

21. A watermark generator (101;2400) for providing a watermark signal (101b, $s_i(t)$; 2420) in dependence on binary message data (101a, m; 2410), the watermark generator comprising:

an information spreader (303; 2430) configured to spread an information unit of the binary message data to a plurality of time-frequency-domain values, to obtain a spread information representation (303a, $\mathbf{c}_f \cdot \mathbf{m}$; 2432);

a synchronization inserter configured to multiplicatively combine the spread information representation (303a, $\mathbf{c}_f \cdot \mathbf{m}$; 2432) with a synchronization spread sequence (**S**; **a**, **b**, **c**; 2442) to obtain a combined information-synchronization representation (304a, $S\circ \mathbf{c}_f \cdot \mathbf{m}$; 2444); and

a watermark signal provider (305,306,307; 2450) configured to provide the watermark signal (101b, $s_i(t)$; 2420) on the basis of the combined information-synchronization representation (304a, $S\circ \mathbf{c}_f \cdot \mathbf{m}$; 2444);

wherein the information spreader is configured to selectively spread a given bit onto a first bit representation, which is a positive multiple of a bit spread sequence ($\mathbf{c}_f$) , or onto a second bit representation, which is a negative multiple of the bit spread sequence ($\mathbf{c}_f$), in dependence on the value of the given bit;

wherein the information spreader is configured to obtain the spread information representation **R** according to

$$\mathbf{R}=\mathbf{c}_f \cdot \mathbf{m},$$

wherein $\mathbf{c}_f$ is a vector of size $N_f \times 1$ representing a bit spread sequence of frequency spreading width N_f,

wherein $\mathbf{m}$ is a vector of size $1 \times N_{mc}$ representing N_{mc} bits of the binary message data, wherein binary values of the bits are represented by entries of the vector $\mathbf{m}$ having a different sign;

wherein the synchronization inserter is configured to combine the spread information with a synchronization sequence $\mathbf{S}$ according to

$$\mathbf{C}=\mathbf{S} \circ \mathbf{R},$$

to obtain the combined information-synchronization representation $\mathbf{C}$,

wherein $\mathbf{S}$ designates a matrix of size $N_f \times N_{mc}$ representing a sequence of N_{mc} synchronization spread sequences ($\mathbf{a}$, $\mathbf{b}$, $\mathbf{c}$); and

wherein $\circ$ designates a Schur element-wise product.