

A METHOD OF INDICATING MINI CELL SIZE

TECHNICAL FIELD OF THE INVENTION

5 This invention relates to telecommunication networks in general and to the transport network of a mobile telephone network. ATM cells are used for transmission of data. The payload of an ATM cell comprises mini cells.

10 DESCRIPTION OF RELATED ART

 To day the size of mini cells to be used for an individual connection is explicit indicated in each one of the mini cells of said individual connection. Typically 6 to 10 bits are used for size indication. The size information is of
15 static nature, i.e. it does not change, as long as the mini cells don't change size.

 It is sometimes necessary to modify the size of a mini cell during an established connection. For example, the speech rate should be changed from full to half, or should the
20 service be changed from speech to data or should a variable rate speech codec be used or should congestion of the traffic between two nodes be reduced by making the mini cells of all connections between said nodes smaller.

 The explicit method of indicating mini cell size means
25 that the bits used for size indication are overhead costs as regards the efficiency with which available band width is used as well as regards the efficiency with which mini cell size is indicated.

30 SUMMARY

 A main object of the invention is to provide a method for indicating the size of mini cells pertaining to an individual

connection only when needed. In accordance with the invention indication of the mini cell size is needed only at instants when the mini cell size is changed. At such instants the new mini cell size to be used for the following mini cells of the
5 connection are indicated.

Another object of the invention is to provide a method of dynamically changing the size of a mini cell during an ongoing connection.

10 Another object of the invention is to provide a method for changing the size of a mini cell belonging to an individual connection with the aid of a mini cell. A mini cell used for this purpose is called a control mini cell.

15 Another object of the invention is to provide a method by which the control mini cell is sent in a control channel different from the channel in which mini cells containing user data are transported.

20 Still another object of the invention is to provide a method by which the control mini cell is transported in the same channel as that in which mini cells containing user data are transported.

A further object of the invention is to provide a synchronization mechanism for effecting change of mini cell size of an individual, ongoing connection. In particular the synchronization mechanism is intended to be used when the
25 control mini cells are transported in a channel which is different from the channel in which mini cells containing user data mini cells are

30 In a mobile telephone system that uses ATM cells in the transport network reduced bandwidth or enhanced use of the available bandwidth will make it possible to add more channels into the system.

A large number of connections, which require a larger CID field, will increase the bandwidth gain if statistical multiplexing is used.

BRIEF DESCRIPTION OF THE DRAWINGS

The invention will be better understood and other characteristics thereof will emerge from the following description made with reference to the appended Figures wherein:

Figure 1 shows the format of an ATM cell transporting mini cells therein,

Figure 2 shows the header of a mini cell transported in the ATM cell in Figure 1,

Figure 3 shows an octet of the cell header of Figure 2, said octet comprising a fixed size length field for indicating the length of the mini cell,

Figure 4 shows an octet in the header of a mini cell, said octet comprising a linearly coded fixed size length field,

Figure 5 is a mapping table,

Figure 6 shows the fixed size length field and an extended fixed size length field created by the extension bit method,

Figure 7 is a mapping table,

Figure 8 shows a fixed size length field and an extended fixed size length field created with the extension code method,

Figure 9 shows the basic format of a mini cell the header of which is provided with a short fixed length field and a length extension qualifier field LEQ comprising different extension codes,

Figure 10 is a table,

Figure 11 shows the extended format of a mini cell,

Figure 12 shows the mini cell of Figure 9 in its extended

format when predefined extension codes are present in the length extension qualifier field,

Figure 13 is a table,

Figure 14 shows an operation and maintenance cell,

5 Figure 15 is a block diagram showing a mini cell header analyzing unit used to extract, from the user data channel, the user data part of a mini cell in which the fixed size length field carries the non-linear coding,

10 Figure 16 shows a mini cell's header and user data as extracted from the user data channel,

Figure 17 is a block diagram of a mini cell header analyzing unit used to extract the user data part of a mini cell from a user data channel using the extension code method,
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Figure 18 is a block diagram of a modified mini cell header analyzing unit used to extract the user data part of a mini cell from a user data channel using either the extension code method or the extension bit method,

20 Figure 19 is a block diagram showing a mini cell header analyzing unit used to extract the user data part of a mini cell from a user data channel using the bit extension method,

Figure 20 shows a mini cell's header wherein the circuit identifier CID is used to indirectly indicate the mini cell size,
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Figure 21 is a mapping table used together with the indirect method for indicating the cell size,

Figure 22 shows different tables which together span up an address space used on the links of the transport network in a mobile telephone system,
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Figure 23 shows a mini cell's header provided with a

synchronization bit used for synchronization of a message that orders change of the size of a mini cell,

5 Figure 24 shows a system global specific mini cell used for changing the size of mini cells belonging to an individual connection,

Figure 25 shows an OAM mini cell used for changing the size of mini cells belonging to the connection to which the OAM cell is related,

10 Figure 26 shows a specific mini cell used for changing the size of mini cells of a connection, said specific mini cell belonging to the connection the mini cells of which are to be modified,

15 Figure 27 shows a combined user data and control mini cell provided with an extension bit in its header and an optional extension field in its payload, said optional extension field containing the new cell size to be used for the mini cells in the connection,

20 Figure 28 shows a combined user data and control mini cell similar to that of Figure 27

Figure 29 illustrates units involved in a control plane signaling method used for changing the size of a mini cell, in accordance with the invention,

25 Figure 30 illustrates in more detail the control plane signaling method described in Figure 23,

Figure 31 is a signaling diagram for the control plane signaling method,

Figure 32 is a block diagram illustrating a second method for changing the size of a mini cell,

30 Figure 33 is a signaling diagram for the second method,

Figure 34 is a sequence of incoming mini cells to a mini cell packetizing device, said Figure illustrating a variant of the second method,

Figure 35 is a signaling diagram related to Figure 34,

5 Figure 36 is a block diagram illustrating delay caused by interpretation of an cell size change order by the control system,

Figure 37 is block diagram illustrating a fourth method for changing the size of a mini cell,

10 Figure 38 is a signaling diagram for the fourth method,

Figure 39 is a block diagram illustrating a fifth method of changing the size of mini cells,

15 Figure 40 is an user data mini cell provided with an optional field for indicating that the mini cell size shall be changed to the new size indicated in the optional field

Figure 41 is a signaling diagram relating to the fifth method and

20 Figure 42 is a block diagram of mobile telephone system provided with the cell header analyzing units.

DETAILED DESCRIPTION OF EMBODIMENTS

In Figure 1 an ATM cell 1 is shown which comprises a header 2 and a payload 3. Conventionally the payload comprises
25 user data relating to an individual connection. In the aforesaid PCT/SE95/00575 patent document an ATM cell is disclosed which in its payload carries one or more mini cells. In the example shown in Figure 1 three mini cells 4, 5 and 6 of different sizes are shown. The ATM header 2 comprises 5
30 octets (1 octet = 8 bits = 1 byte) and its payload 3 comprises 48 octets. Each mini cell 4, 5, 6 comprises a header 7 and user data.

In Figure 2 an example of a mini cell header 7 is shown to comprise 2 octets 8, 9. Other mini cell header sizes are also conceivable depending on the ATM system design. A mini cell header size of 3 octets or more are also conceivable. The mini cell header 7 comprises a circuit identifier CID, which identifies the established connection/circuit, a payload type selector PTS which identifies different payload types such as user data, control data, maintenance data, a length indicator LEN, and a header integrity check field/bit HIC, which supervises the header integrity. The length indicator LEN defines the size of the payload of the individual mini cell.

There is a need for distinguishing between different types of mini cells. The following is required to indicate with the PTS field:

User information of fixed length: The length indicator LEN is not necessary in the header and the user information length is instead configured into the system and into the service. For "GSM full rate", the user information length is 35 octets, for PDC full rate it is 20 octets and for "D-AMPS full rate" it is 23 octets.

User information of different sizes, i.e. user information with variable length: This is the preferred embodiment and will be described below. To use the PTS field in order to indicate user information with variable length is a future proof solution.

User information of different sizes of extended lengths. OAM information per circuit/connection.

Synchronization information: The use of the PTS field for this purpose is optional.

In Figure 3 the cell header 7 is shown to comprise a fixed size length field 10, referred to as LEN field, which is used to indicate the size of the user data of the mini cell to which the header belongs. The size of the mini cell is indicated in this field 10 using linear coding. Linear coding

means that the code corresponds to the actual size of the mini cell. For example, if the cell length is 5 octets a binary 5 (000101) is written into the LEN field. For short mini cell sizes the fixed length field 10 will occupy much band width but all of the occupied band width is not used for transmission of useful information as exemplified by the leading zeros in the two examples given. It should be noted that the LEN field 10 is carried by each mini cell of an individual connection. A further drawback with this fixed size LEN field 10 is that the range of cell sizes which can be expressed with linear coding is restricted. With a fixed size LEN field 10 comprising 6 bits cell sizes from 1 to 64 octets can be indicated. Should larger cell sizes be used for an individual connection, then the length of the fixed size length field 10 must be enlarged which in turn leads to even more waste of band width.

In Figure 4 a fixed size length field 11 is shown. Non-linear coding is used to indicate a wide range of different cell sizes. In the example given 3 bits are used in an octet, for example octet 9, of a mini cell's header. The rest of the bits of the same octet are free and can be used for any of the above listed purposes. This contributes to reduce the overall size of the header which in its turn increases the efficiency with which the band width is used.

In a mobile telephony system mini cells are generated by voice coders. Today the current IS 95 voice coders uses 2, 5, 10 or 22 octets. Using the fixed size length field 10 in accordance with said ANSI document 7 bits would be required in the header of the mini cell in order to indicate a cell size of 22 octets. With the non-linear coding in accordance with Figure 4 the fixed size length field 11 is 3 bits. This gives a band width saving of 10% for an IS 95 voice coder that operates at 2 kbps (5 octets per 20 ms).

In Figure 5 a mapping table 12 is shown which is to be used together with a fixed size length field 11. As appears

from the table the code values do not correspond to the mini cells sizes but instead predefined cell sizes are allocated to a respective code value only three code bits are used.

5 Examples of mini cells sizes are given in the size column of the mapping list. The sizes vary from 4 to 60 octets. Of course the range can be increased, but the maximum number of sizes is given by the number of code bits used.

10 To expand the number of sizes that can be used together with the non-linear coding it is possible to extend the fixed LEN field 11 on demand. Two methods will be described. Either an extension bit in the fixed size LEN field 11 is used as a qualifier for extension of the LEN field 11 and the method is referred to as the extension bit method, or is one of the length field codes used as qualifier for extension of the LEN
15 field 11 in which case the method is referred to as the extension code method.

In Figure 6 a bit 13, also labeled E, following the LEN field 11 is reserved as an extension bit 13. When the extension bit 13 is set to 1 this will indicate that the
20 header of the mini cell comprises an extended LEN field 14 of the same size as the fixed size LEN field. When the extension bit is zero, the cell header comprises the fixed LEN field 11 only.

25 The extended length field 14 comprises 3 bits in the illustrated example.

When the extension bit 13 is set the number of bits available for the mapping table 12 will increase from 3 to 6 bits leaving a mapping table 15 shown in Figure 7. Since the extension bit 13 is reserved for this purpose it cannot be
30 used for code size mapping purposes.

A variation of the extension bit method is to append an extension bit 11B to the extended LEN field 14. The appended extension field is used to indicate if there is a further extended LEN field in the header in the mini cell or not. If

the appended extension bit 11B is set to 1 this indicates that a second extended length field 14A should be added to the header, thus increasing the number of code bits in table 15 from 6 to 9. If the appended extension field comprises a bit which is set to 0 no such second field is used.

In Figure 8 the extension code method is illustrated. In accordance with this method a code in the fixed length field 11 of Figure 4 is reserved and is used as extension code. Suppose, as an example, that binary code 111 in mapping table 12 is used as an extension code. When this code 111 is present in the fixed length field 11 it means that an extended length field 14 should be included in the header of the mini cell. Thus another 3 bits are available for size mapping. This has been illustrated in Figure 8. This method will reduce the number of sizes in mapping table 12 with 1 and will add another seven cell sizes that can be mapped on the additional 8 code values of the extended length field 14.

From band width efficiency view the extension code method is better than the extension bit method since it requires 3 bits, while the extension bit method requires 4 bits. Looking on the value range the extension bit method is better than the extension code method since it provides 16 different cell sizes compared to 14 as provided by the extension code method.

In Figure 9 the extension bit method has been combined with the extension code method in a manner that allows for high efficiency use of the bits available in a cell header while at the same time a broad range of cell sizes are covered and the band width is used efficiently.

The basic format of the mini cell using this combined coding method is shown in Figure 9. The mini cell comprises a header 21 of 2 octets and a payload part 22 which may comprise from 1 to 48 octets. The four least significant bits of the length of the mini cell is indicated in a small fixed size length field 23, LEN field, in the header. The LEN field 23 comprises 4 bits. The header also comprises a CID field 24

which occupies 8 bits and which identifies the circuit to which the mini cell belongs. Also in the header there is a length extension qualifier field 25, LEQ field, and a header integrity field 26, HIC field, both 2 bits long.

5 The length extension qualifier LEQ 25 is defined as a length extension for the payload and as a header extension. When LEQ takes the binary codes of 00, 01 and 10 the mini cell has the basic format shown in Figure 9 and the code bits of LEQ constitute bits to be appended to the LEN field 23. In
10 this case the LEQ field will thus serve as an extension of the LEN field 23.

 In particular, 2^4 different values in LEN field 23 is associated to the binary 00 code existing in the LEQ field 25, 2^4 different values in LEN field 23 is associated with the
15 binary 01 code existing in LEQ field 25 and 2^4 different values in LEN 23 is associated with the binary code 10 existing in LEQ field 25. This is illustrated in Figure 10. This gives a total of 48 different length values in accordance with the following general expression:

20 $[2^{\text{length of LEQ in bits}-m}] \times [2^{\text{length of LEN in bits}}],$

where m is the number of codes used to indicate the extended format of the mini cell.

 Accordingly the payload size can be chosen from forty-eight length values. In the example given the length values
25 are coded as 1 to 48.

 When the LEQ field 25 takes the binary code 11 this signifies that the basic cell format should be extended. The extended format is shown in Figure 11. The LEQ field 25 has a double meaning. The double meaning of LEQ is (i) it is used as
30 the two most significant bits of length indication, i.e. LEQ x 2^4 +LEN as shown in Figure 9 and (ii) it is used as indication of extended header format as shown in Figures 11 and 12, i.e. the LEN field 23 is interpreted as an extension qualifier field 27, EXQ field 27. The EXQ field 27 comprises 4 bits.

Of the four bits of the EXQ field 27 the binary values of 0000 and 0001 are reserved for use together with a further length field 29, LENE field, in the manner shown in Figures 12 and 13. In particular the least significant bit in EXQ field 27 should be appended to the seven bits in the further LENE field 29 in a manner shown in the dashed rectangle 31 in Figure 13. This is similar to what shown in Figure 10. For the EXQ binary value of 0 this will give 128 different length values and for the EXQ binary value of 1 this will give another 128 different length values.

The number of different length values that can be used with this method is given by the following general expression:

$$[2^{\text{number of EXQ bits used}}] \times [2^{\text{number of bits in LENE 29}}]$$

In a preferred embodiment of the invention an EXQ value of 0 is used to indicate mini cell lengths varying from 1 to 128 octets and an EXQ value of 1 is used to indicate mini cell lengths varying from 129 to 256 octets.

It should be noted that the length of the mini cell shown in Figures 9 and 12 is indicated by using a linear coding.

An EXQ value of 2 (binary 0010) is used to signify that the mini cell is an operation and maintenance cell, OAM cell, that comprises a header 32, and an OAM information field 33 as shown in Figure 14. The header 32 is similar to the header 21 in Figure 12. In the LEQ field 25 the binary code 11 is present and in the EXQ field 27 the binary code 0010 is present.

The EXQ code 3 (binary 0011) is used to indicate a fixed length mini cell, for example for the DAMPS system standard. Other EXQ values can be used for other systems standards or services.

EXQ code values 1xxx are used as synchronization cells; wherein xxx is timing information.

In the preferred embodiment a main requirement is that the header of the mini cell at the maximum has a length of 2 octets. Given this restriction the available bits are used in an efficient way to cover all ranges of values.

5 In Figures 9, 11, 12, 14 preferred sizes are indicated under the respective fields. The indicated sizes are just examples and many other sizes of the different fields can be used. Other LEQ and EXQ codes than the indicated can be used as bits that are appended to the LEN field 23 and LENE field
10 29.

In Figure 15 a block schema of a cell header reading device is shown. It comprises a shift register 19, a first counter 20, a latch register 30, a ROM memory 40, a second counter 50 and a multiplexor 60. A bit stream comprising the
15 user data of the mini cells is shifted into shift register 19 at one input thereof. A clock signal controls the frequency at which the data bits are shifted into the shift register 19. The clock signals are counted by the first counter 20 which is used to extract the fixed size length field 11 of a mini cell
20 and write its data into the register 30. The fixed length field or rather the information therein is used as address to the ROM memory 40 which has been configured with the mapping table shown in Figure 5. Accordingly, an individual code, in the following referred to as length code, will correspond to a
25 specific length of the user data. From the ROM memory 40 the size of the user data (mini cell size minus the size of the header) is read and is sent to the second counter 50 which controls the multiplexor 60 such that at the output 61 thereof the user data will appear. Suppose the first counter 20 reads
30 the binary code 011 from the user data channel. This code is used as address to the ROM memory and at this address the cell size 20 is stored. Accordingly the length of the user data should be 20 octets. Next the second counter 50 counts the following 20 octets bit by bit by counting a corresponding
35 number of clock pulses. The multiplexor 60 is shown to have an

arm 62 which is movable between the indicated two positions. Initially counter 50 sets the arm 62 to the lower position shown with dashed lines and no output data will appear at output 61. When the second counter 50 receives the cell size
5 from the ROM memory 40 it moves arm 62 into the upper position. In the upper position arm 62 connects to a line 63 which in its turn is connected to the input user data channel. When the second counter 50 has counted 20 octets it moves arm 62 back to its initial position and the correct number of
10 octets has now been produced at output 61.

In Figure 16 the extraction of the fixed size length field 11 from the user data channel at time t_0 is indicated. At time t_0 counter 20 starts to count 20 octets bit by bit and at time t_1 counter 20 has counted 20 octets. Accordingly arm 62
15 will be in the upper position in Figure 15 between times t_0 and t_1 .

In the cell header reading device shown in Figure 15 a predefined number of length codes and cell sizes are stored in ROM 40. In the cell header reading device shown in Figure 17 a
20 RAM memory 70 is used to which length codes and cell sizes are written from a control system 80. In this manner it is possible to configure different specific mini cell sizes for individual mobile telephone systems.

The mini cell sizes stored in ROM 40 are global in the
25 sense that an individual length code, for example 101, relate to all connections which use mini cells with this length code.

It is, however possible to have a specific mini cell size for a specific connection or for a specific physical link by using the control system 80 and the RAM memory 70 as will be
30 described in connection with Figures 18-27.

Figure 18 is a block diagram of a cell header reading device used for implementing the extension code method. In Figure 18 blocks with the same functions as corresponding blocks in Figures 15 and 17 have the same reference

designations. The circuit differs from that shown in Figures 15 and 17 in that there is a comparator 90 which is used to detect the extension code. If there is a match, the comparator triggers a subtractor 100 which counts down the first counter 5 20 by 3 counts. When this has been done the extended length field, or specifically the data therein, is again written into the register 30. The various sizes associated with the extended field 14 must be added to the RAM memory 70. This implies that the number of cell sizes in the RAM memory will be doubled. In practice this means that a new memory bank will 10 be used in the RAM memory 70. Unit 110 is a D-latch which latches the output value of the comparator 90 and uses it to address the new memory bank, in the RAM memory 70.

The comparator 90 and the subtractor 100 are the units 15 that will handle the extended length field 14 so that the position in the header will be moved when the extension code is detected. Three extra bits will be added to the length field 11 and it is these extra bits that will be used to indicate the cell length. Accordingly the fixed size length 20 field 11 is replaced with the extended length field 14 which is inserted into the data stream.

Compared with the operation of the circuit in Figure 15 or 17 where a field is written into the memory, in Figure 18 another field is written into the memory 70.

25 The cell header reading device shown in Figure 18 can also be used in order to implement the extension bit method. This is indicated in Figure 19. From the register 30 that contains the fixed size length field 11 the extension bit 13 is extracted and is used to increase the address range. The 30 extension bit will count down the first counter 20 with three bits, indicated by the subtractor 100. This implies that three new bits will be written into register 30 and these new three

bits plus the old three bits, i.e. altogether six bits, are used to address the RAM memory 70 as symbolized by the six arrows. In this manner the number of cell sizes has been increased.

5 The ROM memory 40 may have several different mapping tables of the kind shown in Figure 5. It is possible to change from one mapping table to another in response to a predefined length code provided in the header of a mini cell. In this
10 manner it will be possible to switch from a first set of mini cell lengths, for example 4, 8, 16, 20 to a second set of lengths, for example 3, 6, 9, 12. Instead of using a ROM memory 40 configured with the mapping table shown in Figure 5 a RAM memory can be used for the same purpose. This will enable the control system 80 to write in new a new set of mini
15 cell lengths in the RAM memory. The whole table can also be transferred in the control message.

 Instead of providing each cell with a fixed size length field which is used to indicate the mini cell size it is possible to use an implicit method of indicating the mini cell
20 size which does not use any length field in the mini cell header. According to the implicit method of indicating mini cell sizes, information relating to the sizes is resident within the system network. Instead of using a dedicated field to indicate the cell size an existing field in the mini cell
25 header is used. In the preferred embodiment of the invention mini cell sizes are mapped on the identities of established connections. Accordingly sizes are not global but connection oriented.

 The identity of a connection is given by the CID field of
30 a connection. In Figure 20 the mini cell header 7 is shown to comprise a CID field 71. The actual size of the CID field 71 depends on the system but generally two octets should be sufficient. By using the same mapping method as described in connection with Figures 6 and 7 a mapping table 72 results.

Accordingly the fixed length field 11 has been discarded. This will increase the band width efficiency. The CID value is used as address to the RAM memory 70 in Figure 17 and is provided by the control system 80. So, instead of latching the length field 11 in the register 30, the CID value is latched in register 30 and is used as address to the RAM memory 70. In this manner there will be a relation between the identity of the established connection and the length of the mini cells used in the connection. Accordingly no additional memory places are needed for storing the relation between a CID and a size of the mini cell associated with said CID.

At set up of a connection the control system 80 will receive a message which requests (a) that a connection should be set up between to identified end points and (b) that this connection shall use mini cells having a size of X octets. X is supposed to be an integer selected among the available cell sizes. Next the control circuit selects a free CID among logical addresses provided by the ATM network. For the sake of the example CID=7 is selected. The control system 80 will now use 7 as an address to the RAM memory 70 and will write at this address the mini cell size X. The cell header reading device shown in Figure 17 will then operate in the same manner as described. It should be noted that the mapping takes place at connection set-up.

It should be noted that one and the same CID may relate to several different mini cell sizes depending on the fact that cells having the same CID can be transported on different virtual connections VC:s. This is illustrated in Figure 22 wherein a typical address structure used in an ATM network is shown. To each physical link, referred to as physical route, in the ATM network, there is a physical link table 140 having a number of entries, for example the indicated entries 0-23. To each physical link is associated a respective VPI/VCI (virtual path/virtual identifier) table 150. As an example there are 256 virtual paths VP 0-255 in each physical link. In

each VC connection, identified with an VCI-/VPI value, there is as an example 256 mini cell connections each having its individual CID.

5 Certain applications require that the mini cell size is changed on a millisecond basis. Other size change methods fulfilling this requirement use cell size change mini cells transported in the user data channel. This will be described in connection with Figures 24-27. The size change methods used require no processing by the control system 80 and require no
10 synchronization mechanism.

In particular a specific mini cell is used to indicate the new size in accordance with the methods described in connection with Figures 24-27. The new mini cell size is given in the payload 94. Four different types are used:

15 1) a specific EXQ-value defines a size indicator mini cell, as shown in Figure 26,

2) a defined EXQ-value of 2, that is an OAM cell, is used as shown in Figure 25,

20 3) the size change indicating mini cell is indicated by a specific CID-value, for example CID=0, and the connection is identified by the CID field 93 in the payload, as shown in Figure 24,

4) the type described in the fifth method below.

25 The new mini cell size to be used for the following mini cells in a connection is indicated in the length field 94. All mini cells following the mini cells 91, 95, 97 in the data stream and having the same CID will have the new cell size and will have their size indicator set to zero thus indicating that the mini cell is used for user data.

30 Depending on the implementation of the telecommunication system the mini cell that comprises information indicative of mini cell size change is, or is not, in sequence with, i.e. linked with, the first mini cell that has the new length in

the user data stream. This may or may not give rise to a synchronization problem. A mini cell comprising mini cell size change information will in the following be referred to a control mini cell. It should be noted that a control mini cell
5 may comprise, further to mini cell size change information, other information such as user data, control data, OAM data etc.

Five methods for changing the size of a mini cell of an on-going, that is a connection in operation, will be
10 described.

Method 1 - General overview

If the mini cell size should be changed not to frequently, i.e. less frequent than each second, it is
15 proposed in accordance with the present invention to change the size with a control message which is sent over the access protocol between a base station and a controlling node such as for example a mobile switching center MSC. The controlling node will be handling and controlling all equipment involved
20 in the establishment of the mini cell connection, in particular the control system 80 in Figures 17, 18 and 19. The control message is sent over a channel which is different from that in which mini cells are transported. There will thus be necessary to provide synchronization between the sending end
25 of the mini cells and the receiving end of the same mini cells. In accordance with the invention such synchronization is provided by setting a flag in a bit of the header of a mini cell as indicated in Figure 23 wherein the flag bit is labeled 82.

30 Figure 29 shows an embodiment of the cell size change method, in the following referred to as control plane signaling. A cell header reading device 83 identical to that shown in Figure 15 receives the user data bit stream 84 transmitted from a symbolically shown sending device 85. When

it is desired to change the mini cell size the sending device 85 sends a control message 86 which is transported on a control channel and which indicates that the connection having the CID=N shall change its cell size from length L1 to a new length L2, wherein L is the number of octets making up the mini cell.

The control message 86 is sent in a bit stream 87 in a control channel. The bit stream 87 is not synchronized with the bit stream 84. A protocol handler for signaling messages 88 receives the control message and delivers it to the control system 80. The control system 80 will now write the new cell length L1 into the cell header reading device 83 at the address of the identified connection CID=N.

After a sufficient time has lapsed for the control system 80 of the ATM network to process the control message the sending device 85 changes the cell size from L1 to L2 by setting a flag 82 in the first mini cell 89 using the new size L2. This will signal to the receiving side that this cell and the following cells have the new size L2.

Finally, when the first mini cell 89 which carries the flag 82 is received by the cell header reading device 83 and the CID of the mini cell 89 is received by register 30 the new length L2 will be read from the mapping table associated with this CID. The second counter 50 will thus control multiplexor 60 in such manner that the new cell size will be applied to the mini cell 89 in the shift register 10 and all further cells belonging to this connection. In this manner no information will be lost when the cell size is changed.

Control plane signaling can trigger cell size changes on a second basis. This is so because the control system 80 has to process control signals which typically takes about ½-second. Accordingly control plane signaling is slow and require synchronization.

It should be noted that Figure 29 is somewhat simplified in order to clearly teach the synchronizing method. In reality bit stream 87 is interleaved in bit stream 84 on an irregular time basis.

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Method 1 - Detailed description

The general method described in connection with Figure 29 will now be described in detail with reference to a GSM system shown in Figures 30. In Figure 30 a controlling node 100 is a MSC (Mobile Services Switching Center) in a GSM network. The controlling node 100 comprises a control system 80 which controls a transmission equipment 101. The transmission equipment in its turn comprises a mini cell packetizing device 102 that comprises a cell header reading device 103 of the kind shown in Figures 15, 17, 18 and 19. A base station 104 in the GSM network has a similar transmission equipment 105 with a cell depacketizing device 106 that comprises a cell head reading device 107. The transmission equipment 101 has a non shown mini cell depacketizing device and the transmission equipment 105 has a non-shown mini cell packetizing device. The transmission equipments 101 and 105 are exchanging packets over a link 108. Several connections may exist simultaneously but for the sake of the example only one specific connection 109 is considered. The mini cells, symbolically shown at 110 and 111, used in the packets for connection 109 are supposed to have a length of 15 octets each. Traffic is supposed to take place continuously on an on-demand basis. At some instant, depending on an external event, such as for example switching from one service to another, switching from speech service to data service or from half rate speed to full rate speed, the control system 80 initiates a change of the size of the mini cells by sending a respective control message 112, 113 to each one of the transmission equipments 101 and 105. Each control message indicates that for connection 109 shall the mini cell size shall be changed to a new size of 23

octets. Upon reception of the control message no immediate action is taken other than that each equipment now knows that the size is about to be changed to 23 octets. Not until the transmission equipment 105 has information to send it will
5 act. The sequence of operations which then will take place will be described with reference to Figure 31.

Figure 31 is a signaling diagram showing packets being signaled between transmission equipments 101 and transmission equipment 105. The packets are represented by arrows and above
10 the arrows the size of the mini cells used in the packets are shown. Following the arrows in order of time from the top to the bottom the following happens: Initially packets, represented by the two uppermost signal arrows, are exchanged between the transmission equipments 101, 105. The cell length
15 is initially 15 octets. Next the control message 112 is sent to transmission equipment 105. It may also happen that one or more further packets 114 are sent from transmission equipment 101 to transmission equipment 105 before the control message 113 is sent to the transmission equipment 101. As indicated
20 above the receiving transmission equipment 105 will not react on the reception of the control message 112 until it has something to send. The next time the transmission equipment 105 has information to send, it will send it in a packet 115 in which cells having the new cell length of 23 octets are
25 used. In the first cell of the new length a first flag is set. This first flag is flag 82 in Figure 29. The first flag indicates that this mini cell 115 and the following ones are of the new length. In the transmission equipment 101, now acting as a receiving unit, the flag is detected by the cell
30 header reading device of the non shown depacketizing device and is used to trig a mechanism that alters the length of the mini cells to be sent in the future from this transmission equipment 101 to the new length. This mechanism is resident in the mini cell packetizing device 102. Next, when transmission
35 equipment 101 has information to send, it will send it in packets with mini cells of the new length. In the first one,

represented by signal arrow 116, of these a second flag is set. Packets which thereafter are exchanged between transmission equipments 101 and 105 will all have the new length as indicated by packets 117 and 118 and will all have no second flag set.

From the above it is obvious that the first flag in the first mini cell 115 acts as a synchronizing flag. The second flag in mini cell 116 acts as an acknowledgment flag that confirms to the transmission equipment 101 that the transmission equipment 105 has received the synchronizing flag. After exchange of the two flags connection 109 is in a synchronized state in which both transmission equipments 101 and 105 are sending and receiving packets with the new length. In this manner the length of the mini cells used for a particular connection is changed while the connection remains established.

The new length is transmitted in control messages 112, 113. A control message is typically a separate cell, such as an OAM cell. OAM mini cells are sent over a separate connection or in connection 109. The use of control messages has no influence of the bandwidth available to the connection 109. When the mini cell length is to be changed, this will cost only one binary digit, namely the flag bit, of a mini cell. In other words, only one bit needs to be used in the protocol for exchanging information between any two users in the mobile radio system. From bandwidth utilization point of view the control plane signaling method is effective.

It should be recalled that the available mini cell lengths are configured into the mobile radio system.

In accordance with a modification of the above method the first flag is sent in a mini cell that has the old length of 15 octets. This will give the transmission equipment 101 ample time to instruct its mini cell header reading device to change its settings for reception of cells with the new length.

The way the control messages 112, 113 are signaled to the sending and receiving side of the particular connection is not specified by method 1. To summarize method 1 a separate control mini cell is required in order to change the size of the mini cells of the connection and a synchronization mechanism is required in order to make the new mini cell size effective at the right instant.

Method 2.

This method is an example of method 1 and illustrates how the control message is signaled to the transmission equipments 101, 105. In this method a control mini cell is of type 2) above is used. The control mini cell is of the type shown in Figure 24 and contains a field 94 containing the new mini cell size. The CID value of the control mini cell is different from that of connection 109 the cell size of which is to be changed. Accordingly the control message is sent on a connection which is different from that over which user data is sent.

Refer to Figure 32. There is a synchronization problem because the control mini cell is not linked to sequence of the mini cells the size of which is to be modified. Thus there are two mutual independent connections; one, 125, for the control mini cell and another one, 109, for the user data mini cells. For the sake of the example the control mini cell, labeled 127, has a CID value of 0 and the user data mini cell has a CID=7. Connection 109 is active sending packets, which all are filled with mini cells 128 of a length of 15 octets, to transmission equipment 105. At some instant the control system wants to change the mini cell size of the packets in the connection 109 from 15 to 23 octets. The control system orders the sending of a control mini cell 127. The control mini cell is sent in connection 125. The control mini cell has a CID=0 and contains in its payload: (a) the CID field 93 and the length field 94. CID-field 93 refers to the connection the mini cell size of which is to be changed; in this case CID=7

is indicated. In the length field 94 the new length 23 is indicated.

5 The sending transmission equipment 101 and the receiving transmission equipment 105 will both receive the control mini cell 127, compare the signaling diagram shown in Figure 33, and they will both learn the new length of 23 octets. Any of the two units can now start to use the new length at the next synchronization instant. To generate a synchronization instant the receiving transmission equipment 105 sets a flag 129 in
10 the first mini cell 130 with the new length it sends in connection 109. The flag comprises one binary digit and indicates to the receiving unit that the following mini cells, starting with the mini cell in which the flag bit is set, will all have the new length $L=23$. All further mini cells sent from
15 the transmission equipment 105 will then have the new length. When the transmission equipment 101 receives mini cell 130 the cell flag will indicate that the mini cell is formatted with the new length. The transmission equipment 101 will therefore depacketize cell 130 and all following mini cells using the
20 new length of 23 octets. When transmission equipment 101 has anything to send to the transmission equipment 105 it will use the new length as exemplified by arrow 131.

25 The next synchronization instant referred to above may occur when a new service is invoked by the control system or when the control system for other reasons want to change the cell size of the particular connection.

Accordingly it is possible to send the flag bit as soon as the sending unit as well as the receiving unit of the connection 109 has received the control mini cell 127. The
30 flag bit acts as a means of synchronization for switching from the old mini cell size to the new mini cell size. The synchronization is performed by the sending and receiving units themselves with no help from the control system. The transmission equipment first to send sets the synchronization
35 flag when it changes the length of the mini cells it

transmits. The receiving unit, upon reception of the flag, starts to use the new length format.

In the illustrated example the transmission equipment first to send after reception of the control cell 127 is
5 equipment 105. It could as well be transmission equipment 101.

This method operates quicker than method 1 since the control mini cell contains the new size and therefore the transmission equipment 101, 105 need not await to be contacted by the control system in order to have the new cell size.
10 Method 2 has an attractive bandwidth utilization since the overhead in the payload is only one bit each time the size is changed.

OAM mini cells are handled by the operation and maintenance system of the communication system. In accordance
15 with a modification of method 2 the OAM mini cells indicating a size change of connection 109 are handled by the terminating transmission equipment 105.

Refer to Figure 34 and 35. The control mini cell is of the type shown in Figure 24. There is a problem linked to cell
20 size change since many other things happen in the system before the cell size of connection 109 is changed. Figure 34 illustrates a sequence of packets which belong to three different connections having the respective ID:s of CID=1, CID=2 and CID=7. All of the three connections are sending
25 information. Right in the middle of the packet sequence it is desired to change the cell size in the connection the CID value of which is CID=7 and which uses a cell size of 15 octets.

First in the packet sequence comes a packet 140 belonging
30 to the connection that has CID=1 and a length L=10. Next comes a packet 141 that belongs to the connection that has CID=2 and a length L=8. Then comes a packet 142 that belongs to the connection that has CID=7 and L=23 octets. Next comes a packet 143 belonging to CID=1, next a packet 144 belonging to CID=2.

For some reason the control system 80 has decided to change the size of the cells in the connection that has a CID=7 from current L=15 to L=23. The control system 80 therefore inserts a control mini cell 145 into the outgoing data stream. A CID value of CID=0 signifies an OAM mini cell. In the payload of the control mini cell 145 there is a reference to the connection CID=7, the connection the cell size of which shall be changed, and an indication of the new cell length L=23. When the control mini cell is received by the transmission equipment 105 at the receiving side the depacketizing device 106 stores the information given in the control mini cell 145, i.e. it stores the following: in the connection which has CID=7 the cell size shall be changed to 23 octets. The information is stored until the next time a mini cell arrives on this connection. During this time, i.e. during the time from the reception of mini cell 145 and the arrival of the next mini cell which has a CID=7, many other mini cells belonging to the other two connections arrive to the packetizing device 106. This is illustrated by mini cells 146 and 147 belonging to connections CID=1 and CID=2 respectively. When the next cell in connection CID=7 finally arrives, i.e. when mini cell 148 arrives, then the transmission equipment 105 reads it and its cell depacketizing device formats it into segments that all are 23 octets long. All of the above method steps are shown in Figure 35.

The transmission equipment 105 is simple provided the control system 80 allows for cell modification of one connection at a time. If several connections shall change cells size simultaneously the implementation of the transmission equipment 105 will be more complex.

Method 3.

From system point of view mini cell size change in accordance with method 3 is handled by the operation and

maintenance system and the control message for size change is transported in the traffic flow, that is the flow in which user data mini cells are transported.

In method 3 the control mini cell is of type 3) above.

- 5 The control mini cell is an OAM mini cell having an EXQ-value of 2 (binary 11). The OAM cell, shown in Figure 25, has a CID-value equal to that of the connection the cell size of which is to be changed. In other words the OAM cell is transported in the same connection the cell size of which is to be
- 10 changed. This will ensure that the control mini cell is in the right place in the mini cell flow of the connection the size of which shall be changed. With right place is meant that the control mini cell lies between two mini cells of different sizes which both belong to the connection the mini cell size
- 15 of which shall be changed. In principle no synchronization mechanism will therefore be required. However, since an OAM mini cell is not terminated in the same manner as a traffic mini cell synchronization problems may arise. The payload type selector PTS=OAM of the control mini cell indicates that the
- 20 mini cell is an OAM cell. The transmission equipment handle user data mini cells at the traffic plane and they do not handle OAM mini cells. OAM mini cells are handled by the operation and maintenance system at the control plane.

- Similar hardware devices as described in Figures 30 and
- 25 32 are used in method 3 and will therefor not be described again. Following the examples of Figures 30 and 32 it is also supposed that connection 109 shall change size of its mini cells. Connection 109 has a CID=7. The mini cell size of this connection 109 shall be altered from 15 octets to 23 octets.
- 30 Now, refer to Figure 36. On this connection an OAM control mini cell 134 is inserted. The OAM control mini cell comprises the new length of 23 octets in its length field 94A (Figure 25).



The control system 80 shown in Figure 36 comprises an OAM mini cell handler 133 and operates on the cell packetizing device 102 with the cell header reading device 103.

5 Incoming packets to the cell packetizing device 102 arrive from the left in Figure 36 and outgoing packets leave to the right. As long as incoming mini cells have an EXQ value different from 2 (binary 11) the non shown cell packetizing unit of transmission equipment 105 will packetize them in packets which are sent to their destination along arrow 136.

10 The transmission equipment 105 that terminates connection 109 will identify the control mini cell as an OAM mini cell since its PTS=OAM. The transmission equipment 105 will lift the OAM control mini cell out of the data stream and send it, as indicated by arrow 137, to the control system 80 in which
15 it is treated by the OAM cell handler 133. Logic residing in the OAM mini cell handler will interpret the OAM mini cell. The logic will in this case find that the OAM cell refers to CID=7 and that a change of mini cell length shall take place. In response the interpretation the OAM handler returns to the
20 cell packetizing unit a size indication message SI indicating the mini cells on connection 109 shall have a new length SI=23. This message is illustrated by arrow 138. Upon receipt of this message the packetizing device 102 starts reading the incoming mini cells applying the new length L=23. To do so the
25 packetizing device 102 orders a new length setting of the its RAM memory 70.

OAM cells may be of many kinds. An OAM cell contains a message that indicates to the OAM handler 133 what kind of action to be taken by the control system in response to a
30 received OAM cell. For example it can be a message ordering the control system 80 to measure the error bit rate. Another OAM cell may contain a message reporting a hardware malfunction to the control system 80. Still another OAM mini cell message is to order the control to test a multitude of
35 mini cells in some respect, for example with regard to a check



sum. The packetizing device 102 will send all mini cells having PTS=OAM to the OAM handler in the control system 80

Although the OAM mini cell 134 is linked to the instant of cell size change and although it belongs to the connection 109 the mini cell size of which is to be changed it appears as if no synchronization is required. This is however not always true. The system design may also affect the behavior of method 3. Therefore it might be necessary to provide some synchronization mechanism. Why this is so will be described next.

It takes the control system 80 some time to interpret the OAM mini cell 134. Also the return transmission of SI to the packetizing device 102 takes time. During this time a new cell may have arrived in the connection 109. During this time the cell depacketizing unit will not know which mini cell size to apply on the incoming mini cells. This is why some synchronization mechanism is needed.

In accordance with a modification of method 3 it is therefore suggested that the transmission equipment 105 that terminates the connection 109 itself reads the OAM mini cell in order to learn its type. If the OAM mini cell is a cell size modification mini cell the receiving transmission equipment itself will handle the OAM cell and will start to receive and to send mini cells with the new length. This will eliminate the time delay referred to above in connection with interpretation and SI return transmission interval. Still method 3 will ensure that the instant at which the mini cell size is changed will be in correct sequence. In this manner the transmission equipment will ensure that the mini cell size modification cell will be linked to the first mini cell that has the new cell size. According to this modification of method 3 the OAM cell is handled at the traffic plane.

From bandwidth utility point of view method 3 has no loss of bandwidth provided that the frequency at which the mini cell size is changed is moderate.

Method 4.

In accordance with this method the control mini cell is transported in the traffic flow and is handled autonomously by the packetizing and depacketizing devices 102, 106. The control mini cell is terminated in the traffic plane. The loop 137 shown in Figure 36 is eliminated.

In method 4 the control mini cell is of type 1) above the cell is of the type shown in Figure 26 wherein the CID value of the control cell is the same as that of the connection. The method will be described with reference to Figures 37 and 38. In Figure 37 all units are the same as those described in connection with Figures 30 and 32. Suppose transmission equipment 101 is sending packets to transmission equipment 105 and that the packets are filled with mini cells 110 having the length of 15 octets.

At a certain instant the control system 80 orders a change of the length of the cells. The new length shall be 23 octets. A cell size change order is sent to the transmission equipment 101, for example in an OAM cell. The transmission equipment 101 responds to this order by sending a control mini cell 119 of the type shown in Figure 26 to transmission equipment 105. All of the control mini cell 119 is used for carrying cell size change information. The receiving transmission equipment 105 must therefore change the packet size to 23 octets. The question is when.

Provided the transmission equipment 101 sends the control cell 119 after the last one of the cells 110 of length 15 octets, then all further cells which are sent from the transmission equipment 101 may all have the new length of 23 octets. No additional synchronization will be needed.

This is so, because in an ATM network ATM cells are guaranteed to arrive to their destination in the correct order of time. In other words the order of time in which the ATM

cells are sent will not be reversed at the receiving side. So therefore the control system 80 may at any time order the sending transmission equipment 101 to change to the new mini cell size.

5 A signaling diagram shown in Figure 38 illustrates the method. Figure 38 is similar to Figure 31 and will therefore not be described in detail. In Figure 38 the order to change mini cell size is indicated by arrow 121. The control mini cell is indicated by arrow 122 and has a length of 15 octets.
10 It cannot have the new length. The next mini cell 123 sent by the transmission equipment 101 and all mini cells following this are sent with the new size. Upon receipt of the control mini cell 122 at the transmission equipment 105 a mechanism, resident in the packetizing device 107 and the non-shown
15 depacketizing device, will set the length of the mini cells that follow mini cell 122 to the new length of 23 octets. So when mini cell 123, as sent from the transmission equipment 101 arrives to the transmission equipment 105 it will be decoded with the new length. Similarly, when transmission equipment
20 105 sends its next mini cell 124, it will send it with the new length.

 The bandwidth loss adhering method 4 is proportional to the how often, i.e. the rate or frequency by which, the cell size is modified. It is only in connection with mini cell size
25 change that overhead appears in the form of a control mini cell. Mini cells 110 and 120 contain no fields for cell size indication. Therefore they have no overhead relating to cell size indication. This is in contrast to the mini cells described in connection with Figures 4, 6, 9, 11, 12 which all
30 contain a mini cell size indication.

 Method 4 will result in a complex implementation of the transmission equipment 101, 105 since they must be able to handle a great number of connections simultaneously and in a very short time.

It should be remembered that ATM is a connection oriented technique that defines point-to-point connections. This is in contrast to a packet switched network which is of a connection-less nature. In a packet switched network packets that have the same destination can take different routes through the network and may therefore arrive in reversed time order.

Method 5.

Method 5 is an improvement of method 4. Instead of using a complete mini cell for changing the size of the mini cells of an ongoing connection an optional field is inserted into a mini cell carrying user data. When the optional field is present it indicates the new size to be used for the mini cells of the connection. In method 5 the information to change size of the mini cells are transported in the traffic flow. According to method 5 a somewhat different message format is used. In principle an explicit, that is separate, length field is used. The method will be described with reference to Figures 39-41.

In method 5 a cell size modification mini cell 170 of the type shown in Figure 27 is used. An optional field 171 is used to indicate the new length to be used for the mini cells following this cell 170 and belonging to the same connection. The connection is indicated by the CID in the header of the mini cell. In the header there is also an extension bit 13 (Figure 27) which when set indicates that the cell contains the optional length field 171. If the extension bit is set to 0 no field 171 is present in the cell 170.

Method 5 will be described in connection with Figures 39-41. Figure 39 is similar to Figures 30 and the same devices are shown. When the cell size shall be changed the control system 80 will send a cell size modification order 149 to the sending transmission equipment 101 only.

Upon receipt of the order to change mini cell size the sending transmission equipment 101 sets flag 150, adds the optional length field 171 to the mini cell 170, states the new cell size in the added length field, 23 in the example, sends the mini cell in the new length format of 23 octets and continues to send all further mini cells with the new length. Upon reception of the control mini cell 170 transmission equipment 105 detects the flag. In response to the detection of the flag the transmission equipment 105 changes the mini cell size from the current size to the new size indicated in the added optional length field 171. The receiving transmission equipment 105 starts to use the new size of 23 octets, beginning with the control cell and continuing with all of the following cells on this connection.

In Figure 41 a self explanatory graphic illustration of method 5 is shown.

Method 5 will save bandwidth since the length field is not present all the time, as is the case with method 4, but is present only when cell size modification is to take place.

A modification of method 5 is to use the new length format beginning with the mini cell following the mini cell that contains the optional length field 171. In this case the old length format is applied to the mini cell that contains the optional length field.

Comparison of methods 2-5.

Table 1 below is a comparison of some of the characteristic features of methods 2-4. All figures given in the table are estimated.

TABLE 1

Method No.#	Change rate in relation to packet rate	Implementation complexity	Robustness in relation to the explicit length indication method
2	< 1/100	factor 1	0.1
3	1/20	factor 3	0.2
4	1/10	factor 7	0.2
5	1	factor 8	0.1

Methods 2 and 3 are more robust compared to methods 4 and 5 in that they require synchronization. If the control mini cell for some reason is lost, no synchronization will be achieved and no length modification will take place.

Information will be transmitted and received with the old cell length and no information will be lost. Synchronization is achieved by sending two mini cells, one in each direction of the connection, comprising synchronizing information. The robustness will be increased but the change rate will be reduced by a factor of 2. In methods 4 and 5 no synchronization is required. If the control cell is lost it will not be received at the receiving side. The transmitting side will change to the new cell size and will start to send the information in cells of the new length. The receiving side will continue to receive cells which it thinks still have the old size. The received information will therefore be corrupted.

With method 5 it is in principle possible to change the mini cell length of each one of the successive mini cell (each mini cell has a new length; change rate = cell rate) of an

individual connection. If this happens method 5 degenerates and becomes the same method as the explicit length method with the addition of a flag 150.

5 If the size is changed from mini cell to mini cell using method 5, then method 5 is slightly inferior to the explicit length method. If, however, the change rate is each second cell, then method 5 is better than the explicit length method. If the change rate is one per each hundred cells, then method 5 is superior to the explicit length method. Method 5
10 is preferred when the change rate is in the order of one length modification at each ten cells.

Instead of using an extension bit 13 it is also possible to use a specific payload type selector code, PTS-code, similar to extension code method described in connection with
15 Figure 28 to indicate that the mini cell is used for size change of the following mini cells in the same connection. This is shown in Figure 28. In stead of using a separate extension bit 13 in the mini cell header in order to indicate that the mini cell is a user data cell that also comprises
20 information indicating that the mini cell size shall be changed, a code point in the payload type selector PTS is used for this purpose. In particular PTS has a particular code PTSI indicating this.

In Figure 42 there is shown a mobile telephone system
25 comprising an ATM network 200 with a sending unit 201 and a receiving unit 202 are connected via respective link 205 and 206. User data sources 203 are connected to the sending unit over a respective connection as shown symbolically by the lines 209. User data sinks 204 are connected to the receiving
30 unit 202 over a respective connection 210. Connections 209, formed by mini cells, are multiplexed together in the sending unit 201 with a non shown multiplexor. Likewise there is a non shown demultiplexor in the receiving unit 202 that demulti-
plexes mini cells belonging to connections which are termin-
35 ated by the user data sinks 204. In the sending unit 201 there

is a mini cell header reading device 207 of the kind shown in Figure 12 and in the receiving unit there is a similar mini cell header reading device 208 of the kind shown in Figure 12.

In Figure 38 there is shown a mobile telephone system comprising an ATM network 200 with a sending unit 201 and a receiving unit 202 are connected via respective link 205 and 206. User data sources 203 are connected to the sending unit over a respective connection as shown symbolically by the lines 209. User data sinks 204 are connected to the receiving unit 202 over a respective connection 210. Connections 209, formed by mini cells, are multiplexed together in the sending unit 201 with a non shown multiplexor. Likewise there is a non shown demultiplexor in the receiving unit 202 that demultiplexes mini cells belonging to connections which are terminated by the user data sinks 204. In the sending unit 201 there is a mini cell header reading device 207 of the kind shown in Figure 11 and in the receiving unit there is a similar mini cell header reading device 208 of the kind shown in Figure 11.



CLAIMS

- 5 1. A method of indicating the size of mini cell
pertaining to an individual connection, characterized in that
the size is indicated only at instants when the size of the
mini cells of said connection is changed to a new size, and
that after each one of said instants the mini cells of said
10 connection are sent with the new size.
2. A method in accordance with claim 1 for dynamically
changing the size of a mini cell from a first size to a second
size during an established connection, characterized in that
information signalling that the size of minicells used for an
15 individual connection shall be changed is sent in a mini cell
referred to as a control mini cell.
3. A method in accordance with claim 1 for dynamically
changing the size of a mini cell from a first size to a second
size during an established connection, characterized in that
20 the control mini cell is sent in a channel different from the
channel in which mini cells containing user data of said
individual connection are sent.
4. A method in accordance with claim 2 for dynamically
changing the size of a mini cell from a first size to a second
25 size during an established connection, characterized in that
said change of mini cell size is effected following a
synchronizing procedure.
5. A method in accordance with claim 3 for dynamically
changing the size of a mini cell from a first size to a second
30 size during an established connection, characterized in that
said synchronization procedure comprises the following method
steps: (a) any one of said terminating transmission equipment
and said originating transmission equipment upon receipt of
said control mini cell will send its next mini cell pertaining





to said individual connection with said second length and with a flag set in the header of said next mini cell, (b) the other one of said terminating transmission equipment and said originating transmission equipment upon receipt of said next
5 mini cell and detection of said flag starting to send its next mini cell, and all of the following ones pertaining to said individual connection, starting to send mini cells with said second length.

6. A method in accordance with claim 5 for dynamically
10 changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said control mini cell is sent to a transmission equipment situated at the originating side of said individual connection as well as to a transmission equipment situated at the
15 terminating side of said individual connection.

7. A method in accordance with claim 5 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said control mini cell is an operation and maintenance control
20 mini cell containing a field wherein : (a) the identity of said individual connection is indicated and (b) the second length is indicated.

8. A method in accordance with claim 6 for dynamically changing the size of a mini cell from a first size to a second
25 size during an established connection, characterized in that said control mini cell is an operation and maintenance control mini cell which when received at said terminating transmission equipment is sent to a control system for interpretation, said control system upon receipt of said operation and maintenance
30 cell and interpretation thereof as a mini cell that signals the change of size of mini cells pertaining to said connection sending the second length to said terminating transmission equipment, said terminating transmission equipment upon receipt of said second length applying said second length on
35 all further received mini cells pertaining to said connection.



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9. A method in accordance with claim 7 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said control mini cell is an operation and maintenance control mini cell which when received at said terminating transmission equipment is interpreted thereof as a mini cell that signals the change of size of mini cells pertaining to said connection, said terminating transmission equipment in response to said interpretation, from which it learns said second length, applying said second length on all further received mini cells pertaining to said connection.

10. A method in accordance with claim 1 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that the control mini cell is sent in the same channel in which mini cells pertaining to said individual connection are sent.

11. A method in accordance with claim 2 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said control mini cell in its header contains an indication (PTS=SI) indicating that said control cell in its payload comprises said second length.

12. A method in accordance with claim 10 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said originating transmission equipment first sends said control mini cell and thereafter sends all further mini cells belonging to said individual connection with said second length.

13. A method in accordance with claim 11 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that transmission equipment at the terminating end of said

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individual connection upon receipt of said control mini cell starts to read all further mini cells pertaining to said individual connection with said second length.

14. A method in accordance with claim 12 for dynamically
5 changing the size of a mini cell from a first size to a second size during an established connection, characterized in that said control cell is a user data mini cell which in its header comprises an extension bit, which when set indicates that user data mini cell comprises an optional field wherein said second
10 length is indicated.

15. A method in accordance with claim 13 for dynamically changing the size of a mini cell from a first size to a second size during an established connection, characterized in that a control message indicating cell size modification is sent and
15 that following this a synchronization signal is sent.

16. A method of changing the size of a mini cell from a first size to a second size, during an established connection, wherein mini cells belonging to the same connection are transported in a user data channel in accordance with claim 14
20 characterized in that the control message comprises the second size and the identity of the connection, that the control message is sent from a base station to a control system over a control channel using an access protocol, that said control system associates said second cell size with mini cells
25 belonging to said connection, that a cell length modifying indicator, used as said synchronization signal, is set in the first mini cell having the new, second size, that said control system upon detection of the cell length modifying indicator in said first mini cell changes the size of the mini cells
30 belonging to said connection from said first to said second size.

17. A method of changing the size of a mini cell from a first size to a second size during an established connection, wherein an operation and maintenance connection is common to

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all connections and has a predefined first connection identity (CID=0), said operation and maintenance connection being carried by operation and maintenance cells, OAM cells, having said predefined first connection identity characterized by

- 5 providing an OAM cell with an identification (CID) of the connection the band width of which is to be changed, as well as with information on said second size.

- 10 18. A method of changing the size of a mini cell from a first size to a second size during an established connection, wherein a resource management connection is common to all connections and has a predefined first connection identity (CID=0), said resource management connection being carried by resource management cells, having said predefined first connection identity characterized by providing a resource
- 15 management cell with an identification (CID) of the connection the band width of which is to be changed, as well as with information on said second size.

- 20 19. A method of changing the size of a mini cell from a first size to a second size during an established connection, characterized in that a size modifying mini cell is sent in the user data channel in synchronism to the mini cells therein.